## Features

## General Description

- Optimized for $\mathrm{FLEX}^{\circledR}$ and $\mathrm{MAX}^{\circledR}$ architectures
- Offers eight levels of individually maskable interrupts
- Expandable to 64 interrupts
- Offers a flexible priority resolution scheme
- Provides programmable interrupt modes and vectors
- Uses approximately 399 logic elements (LEs) in FLEX devices
- Functionally based on the Intel 8259 device, except as noted in the "Variations \& Clarifications" section on page 79

The Altera ${ }^{\circledR}$ a8259 MegaCore ${ }^{\text {TM }}$ function is a programmable interrupt controller. The a8259 can be initialized by the microprocessor through eight data bus lines (din [7..0] and dout [7..0]), and the ncs, nrd, nwr, int, and ninta control signals. Figure 1 shows the symbol for the a8259.

Figure 1. a8259 Symbol


Table 1 describes the input and output ports of the a8259.

## Table 1. a8259 Ports

| Name | Type | Polarity | Description |
| :---: | :---: | :---: | :---: |
| nmrst | Input | Low | Master reset. When nmrst is asserted, all internal registers assume their default state. The a8259 is idle, awaiting initialization. |
| clk | Input | - | Clock. All registers are clocked on the positive edge of the clock. |
| ncs | Input | Low | Chip select. When low, this signal enables the nwr and nrd signals and register access to and from the a8259. |
| nwr | Input | Low | Write control. When this signal is low (and ncs signal is also low), it enables write transactions to the a8259. |
| nrd | Input | Low | Read control. When this signal is low (and nes signal is also low), it enables read transactions from the a8259. |
| a0 | Input | High | Address. This signal serves as a register selector when writing to and reading from internal a8259 registers. |
| ninta | Input | Low | Interrupt acknowledge. This signal serves as the primary handshake between the a8259 and microprocessor during an interrupt service cycle. |
| nsp | Input | Low | Slave processor. This signal indicates that the a 8259 should be configured as a slave. However, this signal is ignored when the a8259 is configured as a single device. This signal should also be ignored in buffered mode. |
| casin [2..0] | Input | High | Cascade data bus. These bus signals act as a cascade mode control to a slave a8259. If the a8259 is configured as a master, the bus should be driven low. |
| ir [7..0] | Inputs | High (1) | Interrupt request. These are eight maskable, prioritized interrupt service request signals. |
| din [7..0] | Input | - | Data bus. This bus inputs data when writing to internal a8259 registers. |
| int | Output | High | Interrupt. This signal indicates that the a8259 has made an unmasked service request. |
| casout [2..0] | Output | High | Cascade data bus. These bus signals act as cascade mode control, and should be connected to the casin [2. . 0] bus of a slave a8259. When the a8259 is configured as a master, the casout [2..0] bus is ignored. |
| cas_en | Output | High | Cascade directional bus enable. This signal is intended as a tri-state enable signal to external bidirectional I/O buffers on the cascade control bus. |
| dout [7..0] | Output | - | Data bus. The output data when reading from internal 28259 registers. |
| nen | Output | Low | Data enable. This signal indicates that a read cycle is being performed on an internal a8259 register, and it is intended as a tri-state enable to external bidirectional I/O buffers. |

## Note:

(1) The interrupt request signals can be set as active high or positive-edge-triggered via bit 3 of Initialization Command Word (ICW) 1 (see "ICW 1" on page 62 for more information).

## Functional Description

Figure 2 shows the a 8259 block diagram.

Figure 2. a8259 Block Diagram


The int and ninta signals provide the handshaking mechanism for the a8259 to signal the microprocessor. The a8259 requests service via the int signal and receives an acknowledgment of acceptance from the microprocessor via the ninta signal. The int signal is applied directly to the microprocessor's interrupt input. Whenever the a8259 receives a valid interrupt request on an ir pin (ir1 through ir7), the int signal goes high.

The ninta input is connected to the microprocessor's interrupt acknowledgment signal. The microprocessor pulses the ninta signal twice during the interrupt acknowledgment cycle, which tells the a8259 that the interrupt request has been acknowledged. Then, the a 8259 sends the highest priority active interrupt type number onto the din [7..0] bus for the microprocessor to acknowledge.

The ir inputs are used by external devices to request service, and they can be configured for level-sensitive or edge-sensitive operation.

## Programming \& Initialization

The casin [2..0] and casout [2..0] buses, and nsp and cas_en pins are used to implement the cascade interface. These pins are used when more than one a 8259 functions are interconnected in a master/slave configuration, expanding the number of interrupts from 8 up to 64 .

The a8259 operation depends on initial programming. Two types of command words are used for programming the a8259: initialization command words (ICWs) and operation command words (OCWs). ICWs are used to load the a8259 internal control registers, while the OCWs permit the microprocessor to initiate variations in the basic operating modes defined by the ICW registers. Table 2 summarizes how to access the ICW and OCW registers for programming and initialization (for more information on ICW and OCW registers, see "Register Descriptions" on page 62 ).

| Table 2. ICW \& OCW Register Access for Programming \& Initialization |  |  |  | Note(1) |  |  |
| :--- | :---: | :---: | :---: | :--- | :--- | :--- |
| Register | Mnemonics |  |  | Description |  | Access Method |
|  | A0 | D4 | D3 |  |  |  |
| ICW 1 | 0 | 1 | Don't Care | A write with A0 low and D4 high is <br> interpreted as the beginning of an <br> initialization sequence. | Sequential access <br> which starts with ICW 1 <br> and timed by the pulsing |  |
| ICW 2 | 1 | Don't Care | Don't Care | This register always follows ICW 1. | nwr signal. |  |

## Note:

(1) "Don't Care" indicates that the bit has no address significance for this register access method. However, the bit will usually have data significance.

To begin an initialization sequence, the a0 pin must be low, and bit 4 of the din [7. . 0] bus must be high during a valid write cycle. Figure 3 shows the a8259 initialization sequence flow diagram.

Figure 3. a8259 Initialization Sequence Flow Diagram


## Note:

(1) For more information on SINGLE and IC4, see Table 3 on page 62.

Figures 4 and 5 show typical write and read cycles, respectively. The ncs, nwr, and nrd signals enable data to be written to and read from the a8259. This data is clocked by the rising edge of clk. The ncs and nwr signals must be held low for an entire clock cycle in order to read or write valid data.

Figure 4. Typical Write Cycle
Xindicates "don't care." DV indicates "data valid."


Figure 5. Typical Read Cycle
Xindicates "don't care." DV indicates "data valid."


## Register Descriptions

The a8259 contains three type of registers:

- Initialization command word (ICW) registers
- Operation command word (OCW) registers
- Interrupt registers


## Initialization Command Word Registers

There are four ICW command registers: ICW 1, ICW 2, ICW 3, and ICW 4.
/CW1
Input data for ICW 1 is sent via the din [7. . 0] bus (a0 must be low and bit 4 of din [7. . 0] must be held high). ICW 1 is deselected with the rising edge of the nwr signal. Table 3 describes the ICW 1 register format.

Table 3. ICW 1 Register Format (Part 1 of 2)

| Bit | Mnemonic | Description |
| :---: | :---: | :--- |
| 0 | IC4 | When low, this bit causes ICW 4 to be reset (i.e., non- <br> buffered mode, no automatic EOI, and a 3-byte interrupt <br> sequence), and the initialization cycle to skip ICW 4. <br> When high, ICW 4 is accessed normally. |
| 1 | SINGLE | Single mode. When high, this bit indicates that the <br> a8259 is not cascaded with other a8259 functions. <br> When low, this bit causes the a8259 to operate in <br> cascade mode. |
| 2 | ADI | Address interval. When using a 3-byte interrupt <br> sequence, this bit selects the address interval. When <br> low, the address interval is set to eight; otherwise, it is <br> set to four. |


| Table 3. ICW 1 Register Format (Part 2 of 2) |  |  |
| :---: | :---: | :--- |
| Bit | Mnemonic | Description |
| 3 | LTIM | Level-sensitive or edge-triggered input mode. When <br> high, the ir [ 7. . 0] pins are level-sensitive inputs; <br> otherwise, they are positive-edge-triggered. |
| 4 | 1 | This bit is used in conjunction with the a0 signal to select <br> other command registers (see "Interrupt Registers" on <br> page 69). |
| 5 | A5 | These bits set the interrupt vector address (bits 5 <br> through 7) in a 3-byte interrupt sequence (see "Interrupt <br> Sequencing" on page 70). |
| 6 | A6 |  |
| 7 | A7 |  |

## ICW2

ICW 2 is selected after the a0 signal has been high. Input data for ICW 2 is sent via the din [7. . 0] bus, and data is clocked by the rising edge of clk. ICW 2 is deselected with the next falling edge of the nwr signal. Table 4 describes the ICW 2 register format.

Table 4. ICW 2 Register Format

| Bit | Mnemonic | Description |
| :---: | :--- | :--- |
| 0 | A 8 | These bits set the interrupt vector address. For bits 8 <br> through 15, the interrupt vector address is set in a single- <br> byte interrupt sequence mode. For bits 3 through 7, the |
| 1 | A 9 | inter <br> interrupt vector address is set in the same mode. See |
| 2 | A 10 | "Operating Modes \& Sequence of Events" on page 76 |
| 3 | $\mathrm{~A} 11 / \mathrm{T} 3$ |  |

If SINGLE (bit 1 of ICW 1) is low, ICW 3 is the next register selected (see "ICW 3" on page 64). If SINGLE is high, ICW 3 is skipped.

The next register considered is ICW 4 (see Figure 3 on page 61). If bit 0 of ICW 1 is high, then ICW 4 is the next register selected; if it is low, ICW 4 is skipped. When a write transaction is completed for ICW 4-or if it is skipped-the initialization sequence is finished, and the a8259 is now ready to accept interrupts.

## ICW3

If SINGLE is low, ICW 3 must be initialized. Input data for ICW 3 is sent via the $\operatorname{din}[7 . .0]$ bus, and data is clocked by the rising edge of clk . ICW 3 is deselected with the next falling edge of the nwr signal.

The meaning of the ICW 3 contents depends on whether the a8259 is configured as a master or slave. Table 5 describes the ICW 3 register format for the a8259 configured as a master.

Table 5. ICW 3 Register Format (a8259 Master Configuration)

| Bit | Mnemonic | Description |
| :---: | :---: | :--- |
| 0 | S0 | These bits are slave inputs. When high, each bit |
| indicates that the corresponding interrupt request line is |  |  |
| 1 | S1 | a cascaded slave input. For instance, if S2 is high, the |
| 2 | S2 | ir2 pin is treated as a slave input and receives data |
| 3 | S3 | from the int signal of another a8259. |
| 4 | S4 |  |
| 5 | S5 |  |
| 6 | S6 |  |
| 7 | S7 |  |

Table 6 describes the ICW 3 register format when the a 8259 is configured as a slave.

Table 6. ICW 3 Register Format (a8259 Slave Configuration)

| Bit | Mnemonic | Description |  |
| :---: | :---: | :--- | :---: |
| 0 | ID0 | Slave identification. These bits set the slave ID for the <br> a8259. |  |
| 1 | ID1 |  |  |
| 2 | ID2 |  |  |
| 3 | 0 | These bits are not used when the a8259 is configured |  |
| 4 | 0 | as a slave, and they should be low. |  |
| 5 | 0 |  |  |
| 6 | 0 |  |  |
| 7 | 0 |  |  |

At this point in the initialization process, the next register selected depends on whether bit 0 of ICW 1 is high. If bit 0 of ICW 1 is high, ICW 4 is selected (see "ICW 4" on page 65). If bit 0 is low, ICW 4 is skipped and the a 8259 is ready to accept interrupts.

## ICW4

ICW 4 is initialized when bit 0 of ICW 1 is high. Input data for ICW 4 is sent via the din [7. . 0] bus, and the data is clocked by the rising edge of clk. ICW 4 is deselected with the next falling edge of the nwr signal. When a write transaction for ICW 4 is finished-or if ICW 4 is skippedthe initialization sequence is complete, and the a8259 is ready to accept interrupts.

Table 7 describes the ICW 4 register formats.

Table 7. ICW 4 Register Format

| Bit | Mnemonic | Description |
| :---: | :---: | :--- |
| 0 | $\mu$ PM | Microprocessor mode. When this bit is low, the a8259 <br> operates in a 3-byte interrupt sequence mode. If the bit is <br> high, it operates in a single-byte interrupt sequence <br> mode. |
| 1 | AEOI | Automatic end of interrupt. When this bit is high, the <br> AEOI is enabled; otherwise, the AEOI is disabled. |
| 2 | M/S | Master/slave. When this bit is high in buffered mode, the <br> a8259 is configured as a slave, and when it is low, the <br> a8259 is configured as a master. When the device is not <br> in buffered mode, this bit is in a "don't care" condition. |
| 3 | BUF | Buffered mode. When this bit is high, the a8259 is in <br> buffered mode. See "Operating Modes \& Sequence of <br> Events" on page 76 for more information. |
| 4 | SFNM | Special fully nested mode. When this bit is high, the <br> a8259 is in special fully nested mode. |
| 5 | 0 | These bits are unused and should be set low. <br> 6 |
| 7 | 0 | 0 |

## Operation Command Word Registers

Once the appropriate OCW registers have been issued to the a8259, they will be ready for operation.

There are three OCW registers: OCW 1, OCW 2, and OCW 3. These command registers control the operation of the a8259, and permit the interrupt interface operation to be further modified—after the a8259 has been initialized. Unlike the initialization sequence, which requires the outputs of an ICW to be in a special sequence, the OCWs can be issued under program control whenever needed and in any order.

## OCW1

OCW 1 is selected by setting the a 0 pin high. Input data for OCW 1 is sent via the din [7..0] bus, and the data is clocked by the rising edge of clk.

Table 8 describes the OCW 1 register format.

Table 8. OCW 1 Register Format

| Bit | Mnemonic | Description |
| :---: | :---: | :--- |
| 0 | M0 | When more than one of these bits is high, the <br> corresponding interrupt request inputs are masked; <br> otherwise, they are not masked. |
| 1 | M 1 |  |
| 2 | M 2 |  |
| 3 | M 3 |  |
| 4 | M 4 |  |
| 5 | M5 |  |
| 6 | M6 |  |
| 7 | M7 |  |

## OCW2

OCW 2 is selected by setting the a0 pin and resetting bits 3 and 4 of the din [7. . 0] bus low. Input data for OCW 2 is sent via the din [7. . 0] bus, and the data is clocked by the rising edge of $c l k$.

Table 9 describes the OCW 2 register format.

Table 9. OCW 2 Register Format

| Bit | Mnemonic | Description |
| :---: | :---: | :--- |
| 0 | L0 | Interrupt level. These bits determine the interrupt level <br> that is acted upon when bit $6(\mathrm{SL})$ is asserted (see |
| 1 | L 1 | Table 10). |.

Table 10 describes the interrupt levels acted upon when SL (bit 6 of OCW 2) is asserted.

Table 10. Interrupt Levels for SL (Bit 6 of OCW 2)

| Interrupt Level | Mnemonic |  |  |
| :---: | :---: | :---: | :---: |
|  | L2 | L1 | L0 |
| 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 |
| 2 | 0 | 1 | 0 |
| 3 | 0 | 1 | 1 |
| 4 | 1 | 0 | 0 |
| 5 | 1 | 0 | 1 |
| 6 | 1 | 1 | 0 |
| 7 | 1 | 1 | 1 |

Table 11 describes the rotate and EOI commands controlled by bits 5 through 7 of the OCW 2 command register.

Table 11. Rotate \& EOI Commands Controlled by Bits 5 Through 7 of OCW 2

| $\mathbf{R}$ | SL | EOI | Command |
| :---: | :---: | :---: | :--- |
| 0 | 0 | 1 | Non-specific EOI command |
| 0 | 1 | 1 | Specific EOI command |
| 1 | 0 | 1 | Rotate on non-specific EOI command |
| 1 | 0 | 0 | Rotate on automatic EOI mode (set) |
| 0 | 0 | 0 | Rotate on automatic EOI mode (clear) |
| 1 | 1 | 1 | Rotate on specific EOI command (LO, <br> L1, and L2 are used) |
| 1 | 1 | 0 | Specific priority command (LO, L1, and <br> L2 are used) |
| 0 | 1 | 0 | No operation |

## OCW3

OCW 3 is selected by setting the a 0 pin, resetting bit 4 low, and bit 3 high. Input data for OCW 3 is sent via the din [7..0] bus, and the data is clocked by the rising edge of clk.

Table 12 describes the OCW 3 register format.

Table 12. OCW 3 Register Format

| Bit | Decode | Description |
| :---: | :--- | :--- |
| 0 | RIS | Read register command. These bits control which status <br> register will be accessed on the next read cycle (see <br> Table 13). |
| 1 | RR | When this bit is high and RR is high, the a8259 enters <br> the poll mode; the next nrd cycle ends the poll mode <br> (see "Poll Command" on page 75). |
| 2 | P | These bits are used as address decode. Bit 4 must be <br> low and bit 3 must be high. |
| 3 | 1 | 0 |
| 4 | SMM | Special mask mode. These bits are used to enable, set, <br> and clear the special mask mode function (see <br> Table 14). |
| 5 | This bit is unused and should be tied to GND. |  |
| 6 | 0 | ESMM (1) |

## Note:

(1) Enable special mask mode.

Table 13 describes the read register commands for bits 0 and 1 of the OCW 3 command register.

Table 13. Read Register Commands for Bits 0 \& 1 of OCW 3

| RR | RIS | Command |
| :---: | :---: | :--- |
| 0 | 0 | No action |
| 0 | 1 | No action |
| 1 | 0 | Read interrupt request register (IRR) on next read cycle |
| 1 | 1 | Read in-service register (ISR) on next read cycle |

Table 14 describes the special mask mode commands for bits 5 and 6 of the OCW 3 command register.

| Table 14. Read Register Commands for Bits 0 \& 1 of OCW 3 |  |  |
| :---: | :---: | :---: |
| Bit 6 ESMM | Bit 5 SMM | Command |
| 0 | 0 | No action |
| 0 | 1 | No action |
| 1 | 0 | Reset special mask |
| 1 | 1 | Set special mask |

## Interrupt Registers

The a8259 contains two interrupt registers:

- Interrupt request register (IRR)
- In-service register (ISR)


## Interrupt Request Register

The IRR stores all interrupts that are requesting service. In edge-triggered mode (when bit 3 of ICW 1 is low), each ir pin is synchronized to the clk signal. Positive-edge detection is performed, and the result is clocked into the IRR.

In level-triggered mode (when bit 3 of ICW 1 is high), each ir signal is clocked directly into the IRR. The falling edge of the first ninta signal from the microprocessor freezes the IRR so all interrupts can be evaluated. The level of an ir signal from the microprocessor must be maintained until after the falling edge of the ninta signal. Interrupt handshaking protocol must be completed before the next interrupt can be received. Table 15 shows the IRR format.

| Table 15. IRR Format |  |
| :---: | :---: |
| Bit | Decode |
| 0 | IR0 |
| 1 | IR1 |
| 2 | IR2 |
| 3 | IR3 |
| 4 | IR4 |
| 5 | IR5 |
| 6 | IR6 |
| 7 | IR7 |

## In-Service Register

The ISR stores the interrupt level currently being serviced. Data is enabled by the first ninta signal of the interrupt acknowledge sequence. In AEIO mode, data is reset upon the final rising edge of the ninta signal in the interrupt sequence. Otherwise, the microprocessor must issue an EOI command by writing the appropriate value to the OCW 2 command register. Table 16 shows the ISR format.

Table 16. ISR Format

| Bit | Decode |
| :---: | :---: |
| 0 | ISR0 |
| 1 | ISR1 |
| 2 | ISR2 |
| 3 | ISR3 |
| 4 | ISR4 |
| 5 | ISR5 |
| 6 | ISR6 |
| 7 | ISR7 |

## Interrupt Sequencing

The a8259 supports two interrupt sequencing modes:

- 3-byte interrupt sequence mode
- Single-byte interrupt sequence mode


## 3-Byte Interrupt Sequence Mode

The 3-byte interrupt sequence mode provides a 24-bit interrupt vector. The interrupt sequence for this mode is as follows:

1. One or more of the interrupt request signals (ir [7..0]) are high, which sets the corresponding bit in the IRR.
2. The a 8259 checks the priority and masks for the interrupt, and if appropriate, sets the int signal.
3. The microprocessor responds by asserting ninta to the a8259.
4. The a8259 latches the interrupt request signals on the falling edge of the ninta signal (when level-triggered). The a 8259 sets the corresponding bit in the IRR on the following rising edge of ninta. Simultaneously, the bit in the IRR is reset. The a 8259 places a fixed vector opcode of binary 11001101 on the dout [7. .0] bus while the ninta signal is low. The vector opcode indicates that the following two bytes will contain the interrupt vector.
5. The microprocessor responds to the vector opcode by sending two more ninta pulses. The falling edge of the first ninta pulse causes the lower eight interrupt vector address bits to be placed on the dout [7. . 0] bus. The contents of the lower eight bits depend on the value of the address interval bit (bit 2 of ICW 1). See Tables 17 and 18.

Table 17. Contents of the First Interrupt Vector Bytes Notes (1), (2)

| ir | dout7 | dout6 | dout5 | dout4 | dout3 | dout2 | dout1 | dout0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | A7 | A6 | A5 | 1 | 1 | 1 | 0 | 0 |
| 6 | A7 | A6 | A5 | 1 | 1 | 0 | 0 | 0 |
| 5 | A7 | A6 | A5 | 1 | 0 | 1 | 0 | 0 |
| 4 | A7 | A6 | A5 | 1 | 0 | 0 | 0 | 0 |
| 3 | A7 | A6 | A5 | 0 | 1 | 1 | 0 | 0 |
| 2 | A7 | A6 | A5 | 0 | 1 | 0 | 0 | 0 |
| 1 | A7 | A6 | A5 | 0 | 0 | 1 | 0 | 0 |
| 0 | A7 | A6 | A5 | 0 | 0 | 0 | 0 | 0 |

Table 18. Contents of the First Interrupt Vector Bytes Notes (2), (3)

| ir | dout7 | dout6 | dout5 | dout4 | dout3 | dout2 | dout1 | dout0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | A7 | A6 | 1 | 1 | 1 | 0 | 0 | 0 |
| 6 | A7 | A6 | 1 | 1 | 0 | 0 | 0 | 0 |
| 5 | A7 | A6 | 1 | 0 | 1 | 0 | 0 | 0 |
| 4 | A7 | A6 | 1 | 0 | 0 | 0 | 0 | 0 |
| 3 | A7 | A6 | 0 | 1 | 1 | 0 | 0 | 0 |
| 2 | A7 | A6 | 0 | 1 | 0 | 0 | 0 | 0 |
| 1 | A7 | A6 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | A7 | A6 | 0 | 0 | 0 | 0 | 0 | 0 |

Notes to tables:
(1) Interval $=4$ (bit 2 of ICW 1 is high).
(2) A7 through A5 are derived from the contents of bits 5, 6, and 7 of ICW 1 .
(3) Interval $=8$ (bit 2 of ICW 1 is low).
6. The upper eight interrupt vector address bytes are released on the falling edge of the second ninta pulse. The contents of the upper eight interrupt vector address bytes are always derived from the contents of ICW 2. See Table 19.

Table 19. Contents of the Second Interrupt Vector Bytes

| dout7 | dout6 | dout5 | dout4 | dout3 | dout2 | dout1 | dout0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A 15 | A 14 | A 13 | A 12 | A 11 | A 10 | A 9 | A 8 |

7. In AEOI mode, a bit in the ISR is reset on the rising edge of the last ninta pulse. When not in AEOI mode, an appropriate EOI command must be issued to end the interrupt sequence.

Figure 6 shows the timing waveforms for the 3-byte interrupt sequence mode.

Figure 6. 3-Byte Interrupt Sequence Mode Timing Waveforms
Xindicates "Don't Care."


## Single-Byte Interrupt Sequence Mode

The single-byte interrupt sequence mode provides an 8 -bit interrupt vector. The interrupt sequence for this mode is as follows:

1. One or more of the ir [7. . 0] signals are high, which sets the corresponding bit in the IRR.
2. The a 8259 checks the priority and masks for the interrupt, and if appropriate, sets the int signal.
3. The microprocessor responds by asserting ninta to the a8259.
4. The a8259 latches the ir signal on the falling edge of the ninta signal (when level-triggered). The a8259 sets the corresponding ISR bit on the following rising edge of the ninta signal. Simultaneously, the interrupt request bit is reset, and no data is driven onto the dout [7. . 0] bus for this cycle.
5. The microprocessor issues a second ninta pulse. An 8-bit interrupt vector is driven onto the dout [7. .0] bus. See Table 20.

Table 20. Contents of the Interrupt Vector Bytes Note(1)

| ir | dout7 | dout6 | dout5 | dout4 | dout3 | dout2 | dout1 | dout0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | T7 | T6 | T5 | T4 | T3 | 1 | 1 | 1 |
| 6 | T7 | T6 | T5 | T4 | T3 | 1 | 1 | 0 |
| 5 | T7 | T6 | T5 | T4 | T3 | 1 | 0 | 1 |
| 4 | T7 | T6 | T5 | T4 | T3 | 1 | 0 | 0 |
| 3 | T7 | T6 | T5 | T4 | T3 | 0 | 1 | 1 |
| 2 | T7 | T6 | T5 | T4 | T3 | 0 | 1 | 0 |
| 1 | T7 | T6 | T5 | T4 | T3 | 0 | 0 | 1 |
| 0 | T7 | T6 | T5 | T4 | T3 | 0 | 0 | 0 |

Note:
(1) T7 through T3 derive from the contents of ICW 2 (bits 4 through 7).
6. In AEOI mode, the ISR bit is reset on the rising edge of the last ninta pulse. When not in AEOI mode, an appropriate EOI command is issued to end the interrupt sequence.

Figure 7 shows timing waveforms for the single-byte interrupt sequence mode.

Figure 7. Single-Byte Interrupt Sequence Mode Timing Waveforms
Xindicates "Don't Care." VD indicates "Vector Data"


The a8259 supports several operational commands:

- Priority rotation
- EOI
- Special mask mode (SMM)
- Trigger modes
- Poll command


## Priority Rotation

The priority rotation command can adjust the interrupt request priority. The a8259 supports two types of rotation commands: automatic and specific rotation.

The automatic rotation command rotates an interrupt that has just been serviced to the lowest priority. For example, if the ir5 interrupt has just been serviced, it is assigned the lowest priority, and the ir6 interrupt is then given the highest priority. For a system with equal priority interrupts, this process ensures that an interrupt waits for no more than seven other devices to be serviced. Automatic priority can be configured to operate on a non-specific EOI, or an automatic EOI using the OCW 2 command register.

The specific rotation command is similar to the automatic rotation command, except the interrupt to be assigned the lowest priority is specified using bits 0 through 2 of the OCW 2 command register. Specific rotation can be accomplished by issuing a set priority command or a rotate-on-specific EOI command.

## End of Interrupt

The EOI command is used to clear the last interrupt request serviced in an ISR bit. There are two methods used to issue an EOI command: automatic or non-specific.

If an automatic EOI command is used (bit 1 of ICW 4 is high), a non-specific EOI command is issued at the rising edge of the last ninta pulse in the interrupt sequence. A non-specific EOI clears the ISR bit currently set at the highest priority. As long as a specific rotation is not used, a non-specific EOI will always clear last interrupt request serviced in the ISR bit.

When the interrupt priority scheme is disturbed (usually in specific rotation), a specific EOI command is issued to clear the ISR bit of the interrupt request specified in bits 0 through 2 of the OCW 2 command register.

## Special Mask Mode

In SMM, masking an interrupt does not inhibit the reception of lower priority interrupts. Only the interrupt being serviced is masked. With SMM, any interrupt may be selectively enabled using the mask register.

## Trigger Modes

Interrupt request lines may be configured in an edge- or level-triggered mode. In edge-triggered mode, an interrupt request is clocked on the rising edge of the clock. In level-triggered mode, an interrupt is generated merely by placing a high on an ir pin. This level must be maintained until after the falling edge of the first ninta pulse of the interrupt sequence. The trigger mode can be programmed in bit 3 of the ICW 1 command register.

## Poll Command

The poll command provides a way to expand a system allowing the microprocessor to service more than 64 interrupts. In poll mode, the int signal should be ignored. Each a8259 is "polled" individually to determine which interrupts are requesting service. After setting the poll (using bit 2 of OCW 3), the microprocessor simply reads from each a8259 (each interrupt read transaction must be preceded by a write transaction to the poll bit in order to reset the ISR). If an interrupt is pending, the corresponding ISR bit is set on the falling edge of the read cycle, and an interrupt ID byte (bit 7) set high is placed on the dout [7. . 0] bus.

Table 21 shows the interrupt ID word format for the poll command.

Table 21. Interrupt ID Word Format

| Bit | Decode | Description |
| :---: | :--- | :--- |
| 0 | ID0 | Interrupt ID. These bits identify the pending |
| interrupt. |  |  |

Figure 8 shows timing waveforms of the poll mode.

Figure 8. Poll Mode Timing Waveforms for Bit 2 of OCW 3
Xindicates "Don't Care." VD indicates "Vector Data"


## Operating Modes \& Sequence of Events

The a8259 can operate in four different modes:

- Fully nested mode
- Cascade mode
- Special fully nested mode
- Buffered mode


## Fully Nested Mode

The fully nested mode is the default mode, after the master clear. The a8259 will enter this mode when initialization is completed, unless another mode is specifically programmed. When the a8259 is in fully nested mode, the following sequence of events occurs:

1. The interrupt requests are prioritized from 0 (highest priority) to 7 (lowest priority).
2. When an interrupt request is acknowledged, the highest priority, unmasked request is determined, the corresponding bit in the ISR is set, and the IRR bit is reset. While the ISR bit is being set, further interrupts from lower priority sources are ignored. Interrupts from higher priority sources are clocked, causing the int signal to remain active until the interrupt is serviced.
3. During the appropriate handshaking sequence using the inta and ninta signals, the interrupt vector information is placed on the dout [7..0] bus.
4. In AEOI mode, the ISR bit is reset on the rising edge of the last ninta pulse. When not in AEOI mode, an appropriate EOI command is issued to end the interrupt sequence.

## Cascade Mode

The cascade mode provides easy expansion of the a8259. In this mode, a single a8259 is configured as a master, while other a 8259 functions (from 1 to 8 a8259 functions) are configured as slaves. The int signal of each slave is connected to an ir input on the master. The master's int signal serves as an interrupt to the microprocessor. The master's casout [2..0] bus is connected to a slave's casin [2..0] bus. Each a8259 has a unique ncs signal and all other inputs to the a8259 are connected in parallel.

When a slave receives an interrupt, the master asserts its int signal. The master enables the slave by placing the slave's address on the casout [2..0] bus at the rising edge of the first nint a pulse. The slave is then responsible for completing the int and ninta handshaking required by the interrupt sequence. The slave will place its interrupt vector information on the dout [7. . 0] bus as required by the interrupt sequence.

## 3-Byte Interrupt Sequence in Cascade Mode

For a 3-byte interrupt sequence in cascade mode, the handshaking between the int and ninta signals is as follows:

1. The master clocks the ISR bit that corresponds to the slave input on the falling edge of the first ninta pulse. The master also simultaneously resets the IRR bit and places a fixed vector opcode of binary 11001101 on the dout [7. . 0] bus. The vector opcode indicates that the two bytes that follow will contain the interrupt vector. The master enables the slave by placing the slave's address on the casout [7..0] bus at the rising edge of the first ninta pulse.
2. The microprocessor responds to the vector opcode by sending two more ninta pulses. The slave sets the appropriate ISR bit on the falling edge of the second ninta pulse. Simultaneously, the slave's IRR bit is reset. The falling edge of the second ninta pulse also causes the slave to place the lower eight interrupt vector address bits on the slave's dout [7..0] bus. The upper eight interrupt vector address bits are released on the falling edge of the third ninta pulse.
3. Two EOI commands must be issued to end the interrupt sequence: one to the master and the other to the slave.

## Single-Byte Interrupt Sequence in Cascade Mode

In a single-byte interrupt sequence in cascade mode, the handshaking between the int and ninta signals is as follows:

1. The master sets the ISR bit that corresponds to the slave input on the falling edge of the first ninta pulse. The master also simultaneously resets the IRR bit and no data is driven onto the dout [7..0] bus for this cycle. The master enables the slave by placing the slave's address on the casout [7. . 0] bus at the rising edge of the first ninta pulse.
2. The microprocessor issues a second ninta pulse. The slave sets the corresponding ISR bit on the falling edge of the second ninta pulse. Simultaneously, the slave's IRR bit is reset. The slave drives the eight interrupt vector address bits onto the dout [7..0] bus.
3. Two EOI commands must be issued to end the interrupt sequence: one to the master and one to the slave.

The slave's address will remain on the casout [2..0] bus until the rising edge of the last ninta pulse.

## Special Fully Nested Mode

This mode is used in conjunction with the cascade mode to preserve the priority structure within each slave. To operate in this mode, only bit 1 of ICW 4 of the master should be high. The slaves are configured in normal fully nested mode. When a slave is in service, it will not be locked out of the master's priority logic; the master can recognize interrupts from higher priority sources within that slave.

To complete the interrupt service and ensure that all interrupts from the slave have been serviced, the microprocessor sends a non-specific EOI command to the slave and reads the slave's interrupt request register for active low signals. If the interrupt request register is low, a non-specific EOI command is issued to the master. Otherwise, the master services the pending interrupt request.

## Buffered Mode

The buffered mode was originally intended to support board designs where tri-state buffers were needed to drive the data bus. The a8259 has separate nsp and nen signals, and the nen signal is always available. Instead of using the nsp signal, the buffered mode can determine the master/slave configuration by using bits 2 and 3 of the ICW 4 command register.

## Variations \& Clarifications

The following characteristics distinguish the Altera a8259 function from the Intel 8259A device:

- A master clear is provided with the a8259.
- A clock signal has been added, and synchronous design rules have been incorporated to improve operation and reliability. All input signals except nmrst and ir should be synchronous to the clock signal. All inputs must be asserted for one clock cycle to ensure reliable operation.
- Bidirectional I/O pins are split into separate inputs, outputs, and corresponding tri-state control lines. This features makes the a8259 compatible with the bus or multiplexer scheme used internally in a design.
- The dout [7..0], casout [2..0], int, cas_en, en, and nen outputs are driven by complex logic structures and are prone to glitches. If appropriate, these signals should be registered in a target application.
- Because the a8259 can be used in various Altera architectures, no timing information is included in this data sheet.
- Automatic EOI in slave mode is implemented within the a8259.

Copyright © 1995, 1996, 1997, 1998, 1999 Altera Corporation, 101 Innovation Drive, San Jose, CA 95134, USA, all rights reserved.

By accessing this information, you agree to be bound by the terms of Altera's Legal Notice.

