## Features

- AAL1 Segmentation and Reassembly device capable of simultaneously processing up to 2048 bidirectional VCs
- AAL1 cell format for "Structured DS1/E1 N x 64 kbps Service" as per ATM Forum AF-VTOA0078.000 "Circuit Emulation Services Interoperability Specifications" (Nx64 Basic Service, DS1 Nx64 Service with CAS, and E1 Nx64 Service with CAS)
- Two UTOPIA ports (Level 2, 16-bit, 50 MHz ) with loopback function for dual fibre ring applications
- Third UTOPIA port for connection to an external AAL5 SAR processor, or for chaining multiple MT90503 or other SAR or IMA devices
- Flexible aggregation capabilities ( $\mathrm{N} \times 64$ ) to allow any combination of 64 Kbps
- TDM bus provides 32 bidirectional serial TDM streams at 2.048, 4.096, or 8.192 Mbps for up to 4096 TDM 64 Kbps channels
ISSUE 2


## Ordering Information <br> MT90503AG <br> 503 Pin PBGA

- Compatible with H. 100 and H. 110 interfaces
- TDM to ATM transmission latency less than 250 $\mu \mathrm{s}$
- Support for clock recovery - Adaptive Clock Recovery, Synchronous Residual Time Stamp (SRTS) or external
- Support master and slave TDM bus clock operation
- 8- or 16-bit microprocessor port, configurable to Motorola or Intel timing
- Master clock rate up to 80 MHz
- Single power supply device (3.3V)
- IEEE 1149 (JTAG) interface


Figure 1 - Functional Block Diagram

## Description

The MT90503 is an AAL1 SAR, which offers a highly integrated solution for interfacing telecom bus-based systems with ATM networks. The device has the capability of simultaneously processing 2048 bidirectional channels of 64 kbps . The MT90503 can be connected directly to an H. 100 or H. 110 compatible bus. The device also offers the capability of using Channel Associated Signalling (CAS) to support Circuit Emulation Service (CES) for Structured Data Transfer (SDT).

The interface to the TDM port is provided by a TDM bus, which consists of 32 bidirectional serial TDM data streams at $2.048,4.096$, or 8.192 Mbps , therefore allowing for 2048 bidirectional TDM channels operating at 64 kbps . This TDM bus is compatible with the ECTF H. 100 and H. 110 specifications.

The interface to the ATM domain is provided by three UTOPIA ports (Ports A, B, and C). All three of the UTOPIA ports can operate in ATM (master) or PHY (slave) mode. Port A can also be configured as Level 2 M-PHY mode.

## Applications

- ATM Access and Multiplexing Equipment
- Switching Platforms that provide internetworking between TDM and ATM
- ATM Edge Switches
- ATM uplink for expansion of COs, PBXs, or open switching platforms using an adjunct ATM switch
- Integrated Digital Loop Carrier (IDLC)
- SONET or SDH Add and Drop Multiplexers (ADM)
- Next Generation Digital Loup Carrier (NGDIC)
- Digital Subscriber Line Access Multiplexer DSLAM with Gateway


## Switching Feature

Cells coming in from any of the UTOPIA ports can be switched to any other port. The user has the option to change the VPI and VCI fields.


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### 1.0 Introduction

### 1.1 Functional Overview of the MT90503

The MT90503 is an AAL1 SAR, which offers a highly integrated solution for interfacing telecom bus-based systems with ATM networks. The device has the capability of simultaneously processing 2048 bi-directional channels of 64 Kbps . The MT90503 can be connected directly to an H .100 or H .110 compatible bus. The device also offers the capability of using Channel Associated Signalling (CAS) to support Circuit Emulation Service (CES) for Structured Data Transfer (SDT).

The interface to the TDM port is provided by a TDM bus, which consists of 32 bi-directional serial TDM data streams at 2.048, 4.096, or 8.192 Mbps , therefore allowing for 2048 bi-directional TDM channels operating at 64 kbps . This TDM bus is compatible with the ECTF H. 100 and H. 110 specifications.

The interface to the ATM domain is provided by three UTOPIA ports (Ports A, B, and C). All three of the UTOPIA ports can operate in ATM (master) or PHY (slave) mode.

Port A is a UTOPIA Level 2 interface which can operate at up to 50 MHz using a 16- or an 8-bit data bus. This port is capable of operating in ATM-mode (single-PHY), in PHY-mode (slave-mode or level 1), or in slave MPHY-mode (Level 2).

Port B is a UTOPIA Level 2 interface, which can operate at up to 50 MHz using a 16- or an 8-bit data bus but does not support bus addressing. This port is capable of operating in ATM-mode (single-PHY), in PHY-mode (slave-mode).

Port C is a UTOPIA Level 1 interface which can operate at up to 50 MHz using an 8-bit data bus. This port is capable of operating in ATM-mode (master-mode), or in PHY-mode (slave-mode). The MT90503 is capable of performing a UTOPIA loopback from any incoming UTOPIA port to any outgoing UTOPIA port, including a loopback to the port of origin. The loopback capability could be used for dual fibre ring applications.

Figure 3 shows the data flow from the $\mathrm{H} .100 / \mathrm{H} .110$ bus to the TX UTOPIA interface.


Figure 3 - Transmit Data Flow - TDM to UTOPIA

Figure 4 shows the dataflow from the RX UTOPA interface to the H.100/H. 110 bus.


Figure 4 - Receive Data Flow - UTOPIA to TDM

### 2.0 Features Detailed Description

### 2.1 UTOPIA Interface

- Contains 3 UTOPIA ports with transmit and receive interfaces:
- Port A: 16-bit or 8-bit UTOPIA Level 2, ATM mode (single-PHY) or PHY mode (single or multi-PHY). Accepts data rates up to $622 \mathrm{Mb} / \mathrm{s}$.
- Port B: 16-bit or 8-bit UTOPIA Level 2 without bus addressing, ATM mode or PHY mode (restricted to 8 -bit when Port A is in multi-PHY mode). Accepts data rates up to $622 \mathrm{Mb} / \mathrm{s}$.
- Port C: 8-bit UTOPIA, ATM mode or PHY mode
- Supports cell switching through daisy chained SAR/PHY devices via the UTOPIA interface (AAL5 SARs, AAL1 SARs such as MT90503, and AAL2 SARs such as MT90502).
- Supports both UNI and NNI header formats.
- Supports any combination of VCI/VPI concatenation up to 16 bits in ATM Receive direction.
- Supports up to 65536 Virtual Circuits Per UTOPIA Port in ATM Receive direction.
- Rapid timing reference cell processing in Receive/Segmentation direction
- Can eliminate null cells (VPI $=0, \mathrm{VCI}=0$ ) received at UTOPIA A, B, and C port inputs.
- Filters received cells before accessing VCC Look Up Table (LUT).
- UTOPIA VCC loopback for bi-directional ring functionality (from RX A to TX B and from RX B to TX A).
- Per-VCC User cell and OAM cell destination control (for VCs that have a LUT entry).
- Per-UTOPIA-port User cell and OAM cell destination control (for VCs that do not have a LUT entry).


### 2.2 TDM Interface

- H.100/H. 110 compatible.
- Low latency TDM bus to TDM bus loopback of up to 2048 channels.
- Programmable value for the null-octet inserted during an underrun situation.
- Receive buffer replay capability or silent pattern insertion for underrun situations.
- Support of CAS and MFS for DS1(ESF) and E1.
- Automatic Detection of a change in the CAS value, for CAS received in ATM cells, and CAS received from TDM bus.


### 2.3 Clock Recovery

- SRTS clock recovery:
- Dual reference VCs (for redundancy).
- Broadcast SRTS VCs in Transmit /Segmentation direction.
- Adaptive clock recovery:
- Dual reference VCs (for redundancy).
- Limited jitter, precision enhanced, MCLK (chip clock) to 8 kHz dividers.
- Direct 8 kHz clock recovery:
- Can generate an 8 kHz reference using one of 8 multipurpose timing reference pins.
- Supports all n * 8 kHz input reference, ( $1.544 \mathrm{MHz}, 2.048 \mathrm{MHz}, 19.44 \mathrm{MHz}$, etc.) up to 12500 * 8 kHz.
- Output high time and low time of the 8 kHz reference output can be modified relative to input signal.
- The eight multipurpose timing reference pins can be used to support many possible clock recovery configurations, including the following reference signals: SEC8K (MVIP), ATM8K (to/from PHY25 or from PHY155), FNXI (SRTS), CT_NETREF (for CT-Bus).
- Can generate a 20 MHz clock for an external PLL, e.g. MT9042 or MT9044 (output on one of the multipurpose timing reference pins).


### 2.4 ATM SAR

- Supports AAL1 (with pointer, or without pointer byte), CBR-AAL0, and CBR-AAL5 (AAL5-VTOA) Cell formats.
- Supports partially filled cells, with fills from 4 to 47 bytes.
- AAL1 cell format for "Structured DS1/E1 Nx64 Kbit/s Service" as per ATM Forum AF-VTOA-0078.000 "Circuit Emulation Services Interoperability Specification" (Nx64 Basic Service, DS1 Nx64 Service with CAS, and E1 Nx64 Service with CAS).
- VCs carrying 1 to 2048 TDM channels.
- TDM to ATM Transmission latency less than $250 \mu \mathrm{~s}$ (when minimum voice latency desired, and strict multiframe alignment of voice with CAS not required).
- TDM to ATM Transmission latency less than 3.25 ms (when strict multiframe alignment of voice with CAS required).
- ATM to TDM Reception latency less than CDV $+250 \mu \mathrm{~s}$ (when minimum voice latency desired, and strict multiframe alignment of voice with CAS not required).
- ATM to TDM Reception latency less than CDV +6.250 ms (when strict multiframe alignment of voice with CAS required).
- Per VCC monitoring (Receive/Reassembly direction):
- CDV Monitoring and Delay Correction Fields.
- Single cell loss, multiple cell loss, cell misinsertion, AAL1 parity, AAL1-CRC, P-byte Parity, P-Byte Out-of-Range error bits.
- Cell Arrival Counter.
- Underrun Slip Counter.
- Overrun Slip Counter.
- Per VCC monitoring (Transmit/Segmentation direction):
- Cell Transmission Counter.
- Single received cell loss compensated, replacing the payload with a programmable null-octet.
- Support segmentation and reassembly of 2048 full duplex TDM channels ( 2048 without CAS, subtract one TDM channel for each CAS channel carried).
- Support of up to $\pm 61 \mathrm{~ms}$ of CDV in non-multiframing mode, $\pm 45 \mathrm{~ms}$ in T 1 with strict multiframing, $\pm 29$ ms in E1 with strict multiframing.
- AALO Cell Generation / Reception for software implemented SAR function (cell buffer can contain up to 1024 cells).
- Percentage of bandwidth usage register (Transmit/Segmentation direction).


### 2.5 Required External Components

- $128 \mathrm{~K} / 256 \mathrm{~K} / 512 \mathrm{~K} \times 18$ Control Memory (can be used in up to 2 banks).
- Maximum addressable control memory: 1 MB .
- $128 \mathrm{~K} / 256 \mathrm{~K} / 512 \mathrm{~K} \times 18$ Data Memory (can be used in up to 4 banks).
- Maximum addressable data memory: 4 MB .
- 8 kHz to $16 / \mathrm{MHz}$ PLL when $\mathrm{H} .100 / \mathrm{H} .110$ interface used as master mode.
- Clock driver for mclk_src.


### 2.6 Particular Modes of Operation

- Test modes
- TX SAR to RX SAR internal loop-back of some VCs, while MT90503 is running.
- TDM Bus loop-back.


### 2.7 Miscellaneous

- Motorola/Intel CPU Interface (paged memory accesses).
- Programmable Maximum number of Interrupts per second.
- Multipurpose I/Os.
- LED pin generation for UTOPIA Interface.
- Parity bits on memory and UTOPIA interfaces to ensure clocking and memory access integrity.
- CPU-based OAM cell treatment.
- JTAG (IEEE 1149) Test Access Port.
- MCLK speed of 80 MHz .
- Global reset pin with I/O tri-state.
- Global power-down and tri-state.


### 2.8 Power

- 3.3 V core and I/O supply.
- All I/Os are 3.3 V with 5 V tolerance.
- TDM pins are PCI 5 V signalling tolerant (when PCI clamp rail tied to 5 V ).


### 3.0 Pin Designations and Descriptions

The following tables identify each pin of the MT90503 device's main functional areas. A description of each pin is also provided.

## Notes:

1 All outputs are $+3.3 \mathrm{~V}_{\mathrm{DC}}$.
2 All input and output pins that are designated ( $F$ ) can withstand $5 \mathrm{~V}_{\mathrm{DC}}$ being applied to them.
3 All input and output pins that are designated ( $F$ ) are tested with a 50 pF load unless otherwise specified.
4 Designations under the "rst" (reset condition) table column are: $\mathrm{X}=$ undefined; $\mathrm{Z}=$ high impedance; 1 = high (+3.3 $\mathrm{V}_{\mathrm{DC}}$ ).
5 I/O types include: Output (O), Input (I), Bidirectional (I/O), Power (PWR) and Ground (GND).
6 All buses have pins listed in order from MSB to LSB.

GND pins: A2, A4, A7, A8, A9, A12, A15, A18, A21, A22, A23, A25, A26, A28, B1, B29, D1, D26, D29, E1, E6, E29, G1, G29, H1, H29, J1, J29, L11, L13, L15, L17, L19, M1, M29, N11, N13, N15, N17, N19, R1, R11, R13, R17, R19, R29, T29, U11, U13, U15, U17, U19, V1, W11, W13, W15, W17, W19, AA1, AA29, AB1, AB29, AC1, AD5, AD25, AE1, AE13, AE29, AF1, AF29, AG29, AH1, AH2, AH28, AJ2, AJ4, AJ7, AJ8, AJ9, AJ12, AJ15, AJ18, AJ21, AJ22, AJ23, AJ25, AJ26, AJ28

VDD 3.3V pins: A3, A5, A10, A11, A14, A16, A19, A20, A27, A29, B2, C1, C29, K1, K29, L1, L29, P1, P29, T1, V29, W1, W29, Y1, Y29, AC29, AG1, AH3, AH29, AJ1, AJ3, AJ5, AJ10, AJ11, AJ14, AJ16, AJ19, AJ20, AJ27, AJ29

Pins not connected: A13, B3, B4, B7, B8, B9, B12, B13, B14, B15, B16, C3, C4, C7, C8, C9, C12, C13, C14, C15, D6, D7, D8, D9, D12, D13, D14, D15, D24, E5, E8, E9, E12, E13, E14, E15, E25, M5, M28, T25, W25, AB25, AE4, AE5, AE22, AE23, AE25, AF2, AF23

| Pin | rst | Name | I/O | Type | Description |
| :--- | :---: | :---: | :---: | :---: | :--- |
| AE18, AF18, AG18, <br> AH18, AE19, AF19, <br> AG19, AH19, AF20, <br> AG20, AH20, AE21, <br> AF21, AG21, AH21 | Z | inmo_a [14:0] | I | TTL (F) | Intel/Motorola interface address bus. Can <br> be used as a GPI. |
| AE14, AF14, AG14, <br> AH14, AG15, AH15, <br> AE16, AF16 | Z | inmo_d [7:0] | I/O | TTL, 4 mA (F) | Intel/Motorola interface data bus, low bits <br> AE12, AF12, AG12, <br> AH12, AF13, AG13, <br> AH13,AJ13 <br> AF11 |
| Z | inmo_d[15:8] | I/O | TTL, 4 mA (F) | Intel/Motorola interface data bus, high bits. <br> Can be used as a GPIO if an 8 bit CPU |  |
| Interface is used. |  |  |  |  |  |

Table 1-CPU Bus Interface Pins

| Pin | rst | Name | 1/0 | Type | Description |
| :---: | :---: | :---: | :---: | :---: | :---: |
| AG23, AH23, AE24, AF24, AG24, AH24, AJ24, AF25, AG25, AH25, AG26, AH26, AH27, AA26, AA27, AA28, W27, W28 | X | cmem_a [17:0] | 0 | 4 mA | Control memory address bus |
| W26 | 1 | 1. cmem a [18] <br> 2. cmem_cs [1] | $1.0$ | 4 mA | 1. Control memory address bus <br> 2. Control memory chip select 1 . |
| $\begin{gathered} \text { AG28, AF27, } \\ \text { AE26, AE28, } \\ \text { AD27, AC26, } \\ \text { AC28, AB27, AF28, } \\ \text { AE27, AD26, } \\ \text { AC25, AC27, } \\ \text { AB26, AB28, AA25 } \end{gathered}$ | Z | cmem_d[15:0] | I/O | TTL, 4 mA | Control memory data bus. |
| Y25 | 1 | cmem_cs[0] | 0 | 4 mA | Control memory chip select 0 |
| Y27 | X | cmem_ $\overline{\text { bws }}[0]$ | 0 | 4 mA | Control memory byte write select 0 . |
| Y26 | X | cmem_ ${ }^{\text {bws }}$ [1] | 0 | 4 mA | Control memory byte write select 1. |
| Y28 | X | cmem_r $\mathrm{r} / \mathrm{w}$ | 0 | 4 mA | Control Memory R/W. This signal is only used for late write memories. |
| AG27, AF26 | Z | cmem_par[1:0] | I/O | TTL, 4 mA | Control Memory Parity 1:0 |

Table 2 - Control Memory Bus Interface Pins

| Pin | rst | Name | 1/0 | Type | Description |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\begin{aligned} & \text { T28, K26, V28, } \\ & \text { U29, T27, T26, } \\ & \text { R28, K25, J27, } \\ & \text { J26, H28, H27, } \\ & \text { L25, H25, J28, } \\ & \text { G27, F25, F27, } \\ & \text { F29 } \end{aligned}$ | x | dmem_a[18:0] | O | 4 mA | Data memory address bus |
| P25, N28, N26, <br> M27, L28, L26, <br> K28, K27, P28, <br> P26, N29, N27, <br> N25, M26, L27, <br> M25 | z | dmem_d[15:0] | I/O | TTL, 4 mA | Data memory data bus |
| U25 | x | dmem_r/w | $\bigcirc$ | 4 mA | Data Memory R/W. This signal is only used for late write memories. |
| U26 | x | dmem_ $\overline{\text { bws }}[0]$ | $\bigcirc$ | 4 mA | Data memory byte write select 0 . |
| U27 | X | dmem_ $\overline{\text { bws }}[1]$ | $\bigcirc$ | 4 mA | Data memory byte write select 1 . |

Table 3 - Data Memory Bus Interface Pins

| U28 | 1 | dmem_Cs[0] | 0 | 4 mA | Data memory chip select 0 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| V27 | 1 | dmem_cs[1] | O | 4 mA | Data memory chip select 1 |
| V26 | 1 | dmem_cs[2] | 0 | 4 mA | Data memory chip select 2 |
| V25 | 1 | dmem_Cs[3] | 0 | 4 mA | Data memory chip select 3 |
| R27 | Z | dmem_par[0] | I/O | TTL, 4 mA | Data memory parity 0 |
| P27 | Z | dmem_par[1] | I/O | TTL, 4 mA | Data memory parity 1 |

Table 3 - Data Memory Bus Interface Pins

| Pin | rst | Name | I/O | Type | Description |
| :---: | :--- | :--- | :--- | :--- | :--- |
| AD29 |  | mem_clk_i | I | TTL | Data and Control memory clocks |
| AD28 | X | mem_clk_o | O | 4 mA | Data and Control memory clocks |
| R26 |  | mem_clk_positive_i | I | PECL | Data and Control memory clocks, PECL |
| R25 |  | mem_clk_negative_i | I | PECL | Data and Control memory clocks, PECL |
| AF15 | X | mem_clk_positive_o | O | PECL | Data and Control memory clocks, PECL |
| AE15 | X | mem_clk_negative_o | O | PECL | Data and Control memory clocks, PECL |

Table 4 - Data and Control Memory Clock Pins

| Pin | rst | Name | 1/0 | Type | Description |
| :---: | :---: | :---: | :---: | :---: | :---: |
| B17, A17, E22, D17, B19, B18, E16, C18, C16, C20, E18, E17, D16, E21, B22, B21, B20, B23, E19, C19, D23, D22, C22, D19, C24, B24, A24, E20, C21, B26, C23, D21 | Z | ct_d[31:0] | I/O | $\mathrm{PCI}(\mathrm{F})$ | H.100/H. 110 serial data bus |
| G25 | Z | ct_netref1 | I/O | Schmitt, 12 mA <br> (F) | H.100/H. 110 Network Reference 1. |
| G26 | Z | ct_netref2 | I/O | Schmitt, 12 mA (F) | H.100/H. 110 Network Reference 2. |
| C28 | Z | ct_c8_a | I/O | Schmitt, 12 mA (F) | H.100/H. 1108 MHz clock A |
| B28 | Z | ct_c8_b | I/O | Schmitt, 12 mA <br> (F) | H.100/H. 1108 MHz clock B |
| C27 | Z | ct_frame_a | I/O | Schmitt, 12 MA (F) | H.100/H. 110 Frame pulse A |
| C26 | Z | ct_frame_b | I/O | Schmitt, 12 MA (F) | H.100/H. 110 Frame pulse B |
| E26 | Z | ct_fr_comp | 0 | 12 mA (F) | H.100/H. 110 compatibility frame pulse |
| D27 | Z | ct_c2 | 0 | $12 \mathrm{~mA}(\mathrm{~F})$ | MVIP 90-bit clock |

Table 5-H.100/H. 110 Bus Interface Pins

| Pin | rst | Name | 1/0 | Type | Description |
| :---: | :---: | :---: | :---: | :---: | :---: |
| D28 | Z | ct_c4 | 0 | 12 mA (F) | MVIP 90-bit clock times two |
| E23 | Z | ct_c16- | 0 | 12 mA (F) | H-MVIP 16 MHz clock |
| D25 | Z | ct_c16+ | 0 | 12 mA (F) | H-MVIP 16 MHz clock |
| E27 | Z | ct_sclk | 0 | 12 mA (F) | SCBUS system clock |
| E28 | Z | ct_sclkx2 | 0 | 12 mA (F) | SCBUS system clock times two |
| F26 |  | ct_mc | I/O | $\begin{aligned} & \text { TTL, } \\ & 12 \mathrm{~mA}(F) \end{aligned}$ | H.100/H. 110 Message Channel. Open Collector output. |
| E24 | X | mc_clock | 0 | 4 mA (F) | H.100/H. 110 Message Channel extracted clock. 2 MHz . Nominal duty cycle: $62 \%$ high, $38 \%$ low. |
| B27 |  | mc_tx | I | TTL (F) | H.100/H. 110 Message channel transmit data. When this signal is ' 0 ', the MT90503 will drive ct_mc low. When ' 1 ', the MT90503 will not drive ct_mc. |
| C25 | X | mc_rx | 0 | 4 mA (F) | H.100/H. 110 Message channel receive data. The level of this pin directly reflects the value of ct_mc. |
| B25 |  | ct_vdd5_0 | I |  | 5 V power supply used in PCI Buffers of the ct_d[31:0] signals. Can also be connected to 3 V power supply. |
| D20 |  | ct_vdd5_1 | 1 |  | See ct_vdd5_0. |
| D18 |  | ct_vdd5_2 | I |  | See ct_vdd5_0. |
| C17 |  | ct_vdd5_3 | I |  | See ct_vdd5_0. |

Table 5-H.100/H. 110 Bus Interface Pins (continued)

| Pin | rst | Name | I/O | Type | Description |
| :--- | :--- | :--- | :--- | :--- | :--- |
| E10 | Z | recov_a | I/O | TTL, $4 \mathrm{~mA}(F)$ | Clock recovery general I/O A. This pin can be used as a GPIO. |
| D10 | Z | recov_b | I/O | TTL, $4 \mathrm{~mA}(F)$ | Clock recovery general I/O B. This pin can be used as a GPIO. |
| C10 | Z | recov_c | I/O | TTL, $4 \mathrm{~mA}(F)$ | Clock recovery general I/O C. This pin can be used as a GPIO. |
| B10 | Z | recov_d | I/O | TTL, $4 \mathrm{~mA}(F)$ | Clock recovery general I/O D. This pin can be used as a GPIO. |
| E11 | Z | recov_e | I/O | TTL, $4 \mathrm{~mA}(F)$ | Clock recovery general I/O E. This pin can be used as a GPIO. |
| D11 | Z | recov_f | I/O | TTL, $4 \mathrm{~mA}(F)$ | Clock recovery general I/O F. This pin can be used as a GPIO. |
| C11 | Z | recov_g | I/O | TTL, $4 \mathrm{~mA}(F)$ | Clock recovery general I/O G. This pin can be used as a GPIO. |
| B11 | Z | recov_h | I/O | TTL, $4 \mathrm{~mA}(F)$ | Clock recovery general I/O H. This pin can be used as a GPIO. |

Table 6 - Clock Recovery Pins

| Pin | Nom <br> Switch <br> (MHz) | rst | Name | I/O | Type |  |
| :--- | :---: | :--- | :--- | :--- | :--- | :--- |
| AE17 | 0 |  | global_tri_state | I | TTL (F) | (PU) Should be 1 for functional mode, 0 for tristate. |
| AH5 | 1 |  | tck | I | TTL (F) | JTAG Test Clock. Should be 1 when not in use |
| AG5 | 1 |  | tdi | I | TTL (F) | JTAG Test Data In. Should be 1 when not in use |
| AH4 | 1 | X | tdo | O | TTL, $4 \mathrm{~mA}(F)$ | JTAG Test Data Out |
| AG4 | 1 |  | tms | I | TTL (F) | JTAG Test Mode Select. Should be 1 when not in <br> use. |
| AE7 | 1 |  | trst | I | TTL (F) | JTAG Test Reset. Should be 0 when not in use. |

Table 7-Test Pins

| Pin | rst | Name | I/O | Type | Description |
| :---: | :---: | :---: | :---: | :---: | :---: |
| H3 | Z | txa_clk | I/O | TTL, 4 mA (F) | UTOPIA Port A TX Clock |
| C6 | Z | rxa_clk | I/O | TTL, 4 mA (F) | UTOPIA Port A RX Clock |
| C2 |  | phya_alm | 1 | TTL (F) | PHY alarm A <br> This pin can also act as a GPI. |
| D2 | Z | phya_rx_led | I/O | TTL, 12 mA (F) | LED signal. When the LED is on, this pin will be ' 0 '. When the LED is off, this pin will be tri-state. This pin can also act as a GPIO. |
| H5 | Z | phya_tx_led | I/O | TTL, 12 mA (F) | LED signal. When LED is on, this pin will be ' 0 '. When the LED is off, this pin will be tri-state. This pin can also act as a GIPO. |
| J5 |  | $\begin{aligned} & \text { 1. txa_clav } \\ & \text { 2. txa_enb } \end{aligned}$ | $\begin{aligned} & 1.1 \\ & 2.1 \end{aligned}$ | TTL (F) | 1. UTOPIA Port A TX Cell Available (in ATM) <br> 2. UTOPIA Port A TX Enable (in PHY) |
| H2 | $\begin{aligned} & 1 . Z \\ & 2 . Z \end{aligned}$ | 1. $\overline{\text { xxa_enb }}$ <br> 2. txa_clav | $\begin{aligned} & 1.0 \\ & 2.0 \end{aligned}$ | $4 \mathrm{~mA}(\mathrm{~F})$ | 1. UTOPIA Port A TX Enable (in ATM). This pin must be pulled-up externally. <br> 2. UTOPIA Port A TX Cell Available (in PHY). This pin must be pulled-down externally. |
| N2 | Z | txa_soc | 0 | 4 mA (F) | UTOPIA Port A TX Start of Cell |
| $\begin{aligned} & \text { L5, K2, K3, K4, } \\ & \text { K5, J2, J3, J4 } \end{aligned}$ | Z | txa_data[7:0] | 0 | 4 mA (F) | UTOPIA Port A TX Data bus |
| $\begin{aligned} & \text { N4, N5, M2, M3, } \\ & \text { M4, L2, L3, L4 } \end{aligned}$ | Z | txa_data[15:8] | I/O | TTL, 4 mA (F) | UTOPIA Port A TX Data bus Each of these pins can be used as a GPIO. |
| N3 | Z | txa_par | 0 | 4 mA (F) | UTOPIA Port A TX Parity |
| A6 |  | $\begin{aligned} & \text { 1. rxa_clav } \\ & \text { 2. rxa_enb } \end{aligned}$ | $\begin{aligned} & 1.1 \\ & 2.1 \end{aligned}$ | TTL (F) | 1. UTOPIA Port A RX Cell Available (in ATM) <br> 2. UTOPIA Port A RX Enable (in PHY) |
| B6 | $\begin{aligned} & \text { 1. Z } \\ & \text { 2. Z } \end{aligned}$ | 1. $\overline{\text { xa_enb }}$ <br> 2. rxa_clav | $\begin{aligned} & 1.0 \\ & 2.0 \end{aligned}$ | $4 \mathrm{~mA}(\mathrm{~F})$ | 1. UTOPIA Port A RX Enable (in ATM). This pin must be pulled-up externally. <br> 2. UTOPIA Port A RX Cell Available (in PHY). This pin must be pulled-down externally. |
| H4 |  | rxa_soc | I | TTL (F) | UTOPIA Port A RX Start of Cell |
| $\begin{aligned} & \text { E2, E3, E4, D3, } \\ & \text { D4, D5, C5, B5 } \end{aligned}$ |  | rxa_data[7:0] | I | TTL (F) | UTOPIA Port A RX Data bus |

Table 8 - UTOPIA Interface Pins

| Pin | rst | Name | I/O | Type | Description |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\begin{aligned} & \text { G3, G4, G5, F1, } \\ & \text { F2, F3, F4, F5 } \end{aligned}$ |  | rxa_data[15:8] | I | TTL (F) | UTOPIA Port A RX Data bus Each of these pins can be used as a GPI. |
| G2 |  | rxa_par | 1 | TTL (F) | UTOPIA Port A RX Parity |
| W4 | Z | txb_clk | I/O | TTL, 4 mA | UTOPIA Port B TX Clock |
| N1 | Z | rxb_clk | I/O | TTL, 4 mA (F) | UTOPIA Port B RX Clock |
| AB5 |  | phyb_alm | I | TTL (F) | PHY alarm B <br> This pin can also act as a GPI. |
| W5 | Z | phyb_rx_led | I/O | TTL, 12 mA (F) | LED signal. When LED is on, this pin will be ' 0 '. When the LED is off, this pin will be tri-state. This pin can also act as a GPIO. |
| T5 | Z | phyb_tx_led | I/O | TTL, $12 \mathrm{~mA} \mathrm{~F}^{\text {) }}$ | LED signal. When LED is on, this pin will be ' 0 '. When the LED is off, this pin will be tri-state. This pin can also act as a GPIO. |
| W2 |  | $\text { 1. } \frac{\mathrm{txb} \text { clav }}{\mathrm{txb} \mathrm{enb}}$ | $\begin{aligned} & 1.1 \\ & 2.1 \end{aligned}$ | TTL (F) | 1. UTOPIA Port B TX Cell Available (in ATM) <br> 2. UTOPIA Port B TX Enable (in PHY) |
| W3 | $\begin{aligned} & 1 Z \\ & 2 Z \end{aligned}$ | 1. $\overline{\mathrm{txb}}$ _enb <br> 2. txb_clav | $\begin{aligned} & 1.0 \\ & 2.0 \end{aligned}$ | 4 mA (F) | 1. UTOPIA Port B TX Enable (in ATM) This pin must be pulled-up externally. <br> 2. UTOPIA Port B TX Cell Available (in PHY). This pin must be pulled-down externally. |
| AD2 | Z | txb_soc | 0 | 4 mA (F) | UTOPIA Port B TX Start of Cell |
| AA2, AA3, AA4, AA5, Y2, Y3, Y4, Y5 | Z | txb_data[7:0] | 0 | 4 mA (F) | UTOPIA Port B TX Data bus |
| $\begin{aligned} & \text { AC3, AC4, AG2, } \\ & \text { AB2, AB3, AB4 } \end{aligned}$ | Z | txb_data[13:8] | I/O | TTL, 4 mA (F) | UTOPIA Port B TX Data bus Each of these pins can be used as a GPIO. |
| AC2 | Z | $\begin{aligned} & \text { 1. txb_data[14] } \\ & \text { 2. rxa_addr[4] } \end{aligned}$ | $\begin{aligned} & 1.0 \\ & 2.1 \end{aligned}$ | TTL, 4 mA (F) | 1. UTOPIA Port B TX Data bus <br> 2. UTOPIA Port A RX Address 4 <br> This pins can be used as a GPIO. |
| AD4 | Z | 1. txb_data [15] <br> 2. txa_addr[4] | $\begin{aligned} & 1.0 \\ & 2.1 \end{aligned}$ | TTL, 4 mA (F) | 1. UTOPIA Port B TX Data bus <br> 2. UTOPIA Port A TX Address 4 <br> This pin can be used as a GPIO. |
| AD3 | Z | txb_par | 0 | 4 mA (F) | UTOPIA Port B TX Parity |
| P4 |  | $\begin{aligned} & \text { 1. rxb_clav } \\ & \text { 2. } \frac{\text { rxb_enb }}{} \end{aligned}$ | $\begin{aligned} & 1.1 \\ & 2.1 \end{aligned}$ | TTL (F) | 1. UTOPIA Port B RX Cell Available (in ATM) <br> 2. UTOPIA Port B RX Enable (in PHY) |
| P5 | $\begin{aligned} & 1 . Z \\ & 2 . Z \end{aligned}$ | 1. $\overline{\mathrm{rxb}}$ _enb <br> 2. rxb_clav | $\begin{aligned} & 1.0 \\ & 2.0 \end{aligned}$ | 4 mA (F) | 1. UTOPIA Port B RX Enable (in ATM). This pin must be pulled-up externally. <br> 2. UTOPIA Port B RX Cell Available (in PHY). This pin must be pulled-down externally. |
| V2 |  | rxb_soc | I | TTL (F) | UTOPIA Port B RX Start of Cell |
| T3, T4, R2, R3, R4, R5, P2, P3 |  | rxb_data [7:0] | I | TTL (F) | UTOPIA Port B RX Data bus |
| U3, U4, U5, T2 |  | 1. rxb_data [11:8] <br> 2. rxa_addr [3:0] | $\begin{aligned} & 1.1 \\ & 2.1 \end{aligned}$ | TTL (F) | 1.UTOPIA Port B RX Data bus <br> 2. UTOPIA Port A RX Address bus |
| V4, V5, U1, U2 |  | 1. rxb data [15:12] <br> 2. txa_addr [3:0] | $\begin{aligned} & 1.1 \\ & 2.1 \end{aligned}$ | TTL (F) | 1. UTOPIA Port B RX Data bus <br> 2. UTOPIA Port A TX Address bus |

Table 8 - UTOPIA Interface Pins (continued)

| Pin | rst | Name | I/O | Type | Description |
| :---: | :---: | :---: | :---: | :---: | :---: |
| V3 |  | rxb_par | I | TTL (F) | UTOPIA Port B RX Parity |
| AG10 | Z | txc_clk | I/O | TTL, 4 mA (F) | UTOPIA Port C TX Clock |
| AF7 | Z | rxc_clk | I/O | TTL, 4 mA (F) | UTOPIA Port C RX Clock |
| AE10 |  | $\text { 1. } \frac{\text { txc clav }}{\mathrm{txc}}$ | $\begin{aligned} & 1.1 \\ & 2.1 \end{aligned}$ | TTL (F) | 1. UTOPIA Port C TX Cell Available (in ATM). <br> 2. UTOPIA Port C TX Enable (in PHY). |
| AF10 | $\begin{aligned} & \text { 1. Z } \\ & \text { 2. } Z \end{aligned}$ | 1. $\overline{\text { txc_enb }}$ <br> 2. txc_clav | $\begin{aligned} & 1.0 \\ & 2.0 \end{aligned}$ | 4 mA (F) | 1. UTOPIA Port C TX Enable (in ATM). This pin must be pulled-up externally. <br> 2. UTOPIA Port C TX Cell Available (in PHY). This pin must be pulled-down externally. |
| AH9 | Z | txc_soc | 0 | 4 mA (F) | UTOPIA Port C TX Start of Cell |
| AG9, AF9, AE9, AH8, AG8, AF8, AE8, AH7 | Z | txc_data[7:0] | 0 | 4 mA (F) | UTOPIA Port C TX Data bus |
| AG7 | Z | txc_par | 0 | 4 mA (F) | UTOPIA Port C TX Parity |
| AH6 |  | $\text { 1. } \frac{\mathrm{rxc} \text { clav }}{\mathrm{rxc} \mathrm{enb}}$ | $\begin{aligned} & 1.1 \\ & 2.1 \end{aligned}$ | TTL (F) | 1. UTOPIA Port C RX Cell Available (in ATM) <br> 2. UTOPIA Port C RX Enable (in PHY) |
| AJ6 | $\begin{aligned} & \text { 1. Z } \\ & \text { 2. Z } \end{aligned}$ | 1. $\overline{\mathrm{rxc}} \mathrm{enb}$ <br> 2. rxc_clav | $\begin{aligned} & 1.0 \\ & 2.0 \end{aligned}$ | 4 mA (F) | 1. UTOPIA Port C RX Enable (in ATM). This pin must be pulled-up externally. <br> 2. UTOPIA Port C RX Cell Available (in PHY). This pin must be pulled-down externally. |
| AG6 | 1.2 | rxc_soc | 1 | TTL (F) | UTOPIA Port C RX Start of Cell |
| AF6, AE6, AF5, AF4, AE2, AE3, AF3, AD1 |  | rxc_data[7:0] | 1 | TTL (F) | UTOPIA Port C RX Data bus |
| AH10 |  | rxc_par | I | TTL (F) | UTOPIA Port C RX Parity |

Table 8 - UTOPIA Interface Pins (continued)

| Pin | rst | Name | I/O | Type | Description |
| :--- | :---: | :---: | :---: | :---: | :--- |
| J25 |  | pllvss_110 | I |  | VSS pin for the CT PLL. See Figure 5, "PLL Pin Connections," <br> on page 25 for recommended connections. |
| F28 |  | Pllvdd_110 | I |  | VDD pin for the CT PLL. See Figure 5, "PLL Pin Connections," <br> on page 25 for recommended connections. |
| H26 |  | Plllp2_110 | I |  | Loop-filter pin for the CT PLL. See Figure 5, "PLL Pin <br> Connections," on page 25 for recommended connections. |
| G28 |  | pllagn_110 | O |  | Analog Ground pin for the CT PLL. See Figure 5, "PLL Pin <br> Connections," on page 25 for recommended connections. |
| AC5 |  | pllvss_300 | I |  | VSS pin for the FC PLL. See Figure 5, "PLL Pin Connections," <br> on page 25 for recommended connections. |
| AG3 |  | pllvdd_300 | I |  | VDD pin for the FC PLL. See Figure 5, "PLL Pin Connections," " <br> on page 25 for recommended connections. |

Table 9 - Phase Lock Loop (PLL) Pins

| Pin | rst | Name | I/O | Type | Description |
| :--- | :---: | :---: | :---: | :---: | :---: |
| E7 |  | proc_out | O |  | Process Monitor Pin Output. Must not be connected. |

Table 10 - Process Monitor Pins

| Pin | Location |
| :---: | :---: |
| A2 | GND |
| A3 | VDD |
| A4 | GND |
| A5 | VDD |
| A6 | $\begin{aligned} & \text { 1. rxa_clav } \\ & \text { 2.r rxa_enb } \end{aligned}$ |
| A7 | GND |
| A8 | GND |
| A9 | GND |
| A10 | VDD |
| A11 | VDD |
| A12 | GND |
| A13 | N/C |
| A14 | VDD |
| A15 | GND |
| A16 | VDD |
| A17 | ct_d[30] |
| A18 | GND |
| A19 | VDD |
| A20 | VDD |
| A21 | GND |
| A22 | GND |
| A23 | GND |
| A24 | ct_d[5] |
| A25 | GND |
| A26 | GND |
| A27 | VDD |
| A28 | GND |
| A29 | VDD |
| B1 | GND |
| B2 | VDD |
| B3 | N/C |


| Pin | Location | Pin | Location |
| :---: | :---: | :---: | :---: |
| B9 | N/C | C16 | ct_d[23] |
| B10 | recov_d | C17 | ct_vdd5_3 |
| B11 | recov_h | C18 | ct_d[24] |
| B12 | N/C | C19 | ct_d[12] |
| B13 | N/C | C20 | ct_d[22] |
| B14 | N/C | C21 | ct_d[3] |
| B15 | N/C | C22 | ct_d[9] |
| B16 | N/C | C23 | ct_d[1] |
| B17 | ct_d[31] | C24 | ct_d[7] |
| B18 | ct_d[26] | C25 | mc_rx |
| B19 | ct_d[27] | C26 | ct_frame_b |
| B20 | ct_d[15] | C27 | ct_frame_a |
| B21 | ct_d[16] | C28 | ct_c8_a |
| B22 | ct_d[17] | C29 | VDD |
| B23 | ct_d[14] | D1 | GND |
| B24 | ct_d[6] | D2 | phya_rx_led |
| B25 | ct_vdd5_0 | D3 | rxa_data[4] |
| B26 | ct_d[2] | D4 | rxa_data[3] |
| B27 | mc_tx | D5 | rxa_data[2] |
| B28 | ct_c8_b | D6 | N/C |
| B29 | GND | D7 | N/C |
| C1 | VDD | D8 | N/C |
| C2 | phya_alm | D9 | N/C |
| C3 | N/C | D10 | recov_b |
| C4 | N/C | D11 | recov_f |
| C5 | rxa_data[1] | D12 | N/C |
| C6 | rxa_clk | D13 | N/C |
| C7 | N/C | D14 | N/C |
| C8 | N/C | D15 | N/C |
| C9 | N/C | D16 | ct_d[19] |
| C10 | recov_c | D17 | ct_d[28] |


| Pin | Location |
| :---: | :---: |
| D23 | ct_d[11] |
| D24 | N/C |
| D25 | ct_c16+ |
| D26 | GND |
| D27 | ct_c2 |
| D28 | ct_c4 |
| D29 | GND |
| E1 | GND |
| E2 | rxa_data[7] |
| E3 | rxa_data[6] |
| E4 | rxa_data[5] |
| E5 | N/C |
| E6 | GND |
| E7 | proc_out |
| E8 | N/C |
| E9 | N/C |
| E10 | recov_a |
| E11 | recov_e |
| E12 | N/C |
| E13 | N/C |
| E14 | N/C |
| E15 | N/C |
| E16 | ct_d[25] |
| E17 | ct_d[20] |
| E18 | ct_d[21] |
| E19 | ct_d[13] |
| E20 | ct_d[4] |
| E21 | ct_d[18] |
| E22 | ct_d[29] |
| E23 | ct_c16- |
| E24 | mc_clock |

Table 11 - Pin Names Listed by Location

| Pin | Location | Pin | Location | Pin | Location |
| :---: | :---: | :---: | :---: | :---: | :---: |
| B4 | N/C | C11 | recov_g | D18 | ct_vdd5_2 |
| B5 | rxa_data[0] | C12 | N/C | D19 | ct_d[8] |
| B6 | 1. $\overline{\text { rxa_enb }}$ <br> 2. rxa_clav | C13 | N/C | D20 | ct_vdd5_1 |
| B7 | N/C | C14 | N/C | D21 | ct_d[0] |
| B8 | N/C | C15 | N/C | D22 | ct_d[10] |
| F1 | rxa_data[12] | J3 | txa_data[1] | L29 | VDD |
| F2 | rxa_data[11] | J4 | txa_data[0] | M1 | GND |
| F3 | rxa_data[10] | J5 | $\begin{aligned} & \text { 1. txa_clav } \\ & \text { 2. } \overline{\text { txa_enb }} \end{aligned}$ | M2 | txa_data[13] |
| F4 | rxa_data[9] | J25 | pllvss_110 | M3 | txa_data[12] |
| F5 | rxa_data[8] | J26 | dmem_a[9] | M4 | txa_data[11] |
| F25 | dmem_a[2] | J27 | dmem_a[10] | M5 | N/C |
| F26 | ct_mc | J28 | dmem_a[4] | M25 | dmem_d[0] |
| F27 | dmem_a[1] | J29 | GND | M26 | dmem_d[2] |
| F28 | Pllvdd_110 | K1 | VDD | M27 | dmem_d[12] |
| F29 | dmem_a[0] | K2 | txa_data[6] | M28 | N/C |
| G1 | GND | K3 | txa_data[5] | M29 | GND |
| G2 | rxa_par | K4 | txa_data[4] | N1 | rxb_clk |
| G3 | rxa_data[15] | K5 | txa_data[3] | N2 | txa_soc |
| G4 | rxa_data[14] | K25 | dmem_a[11] | N3 | txa_par |
| G5 | rxa_data[13] | K26 | dmem_a[17] | N4 | txa_data[15] |
| G25 | ct_netref1 | K27 | dmem_d[8] | N5 | txa_data[14] |
| G26 | ct_netref2 | K28 | dmem_d[9] | N11 | GND |
| G27 | dmem_a[3] | K29 | VDD | N13 | GND |
| G28 | pllagn_110 | L1 | VDD | N15 | GND |
| G29 | GND | L2 | txa_data[10] | N17 | GND |
| H1 | GND | L3 | txa_data[9] | N19 | GND |
| H2 | 1. $\overline{\text { txa_enb }}$ <br> 2. txa_clav | L4 | txa_data[8] | N25 | dmem_d[3] |
| H3 | txa_clk | L5 | txa_data[7] | N26 | dmem_d[13] |
| H4 | rxa_soc | L11 | GND | N27 | dmem_d[4] |
| H5 | phya_tx_led | L13 | GND | N28 | dmem_d[14] |
| H25 | dmem_a[5] | L15 | GND | N29 | dmem_d[5] |
| H26 | Plllp2_110 | L17 | GND | P1 | VDD |
| H 27 | dmem_a[7] | L19 | GND | P2 | rxb_data[1] |


| Pin | Location |
| :---: | :---: |
| E25 | N/C |
| E26 | ct_fr_comp |
| E27 | ct_sclk |
| E28 | ct_sclkx2 |
| E29 | GND |
| P26 | dmem_d[6] |
| P27 | dmem_par[1] |
| P28 | dmem_d[7] |
| P29 | VDD |
| R1 | GND |
| R2 | rxb_data[5] |
| R3 | rxb_data[4] |
| R4 | rxb_data[3] |
| R5 | rxb_data[2] |
| R11 | GND |
| R13 | GND |
| R17 | GND |
| R19 | GND |
| R25 | mem_clk_negative_i |
| R26 | mem_clk_positive_i |
| R27 | dmem_par[0] |
| R28 | dmem_a[12] |
| R29 | GND |
| T1 | VDD |
| T2 | 1. rxb_data[8] <br> 2. rxa_addr[0] |
| T3 | rxb_data[7] |
| T4 | rxb_data[6] |
| T5 | phyb_tx_led |
| T25 | N/C |
| T26 | dmem_a[13] |
| T27 | dmem_a[14] |
| T28 | dmem_a[18] |
| T29 | GND |

Table 11 - Pin Names Listed by Location (continued)

| Pin | Location | Pin | Location | Pin | Location | Pin | Location |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| H28 | dmem_a[8] | L25 | dmem_a[6] | P3 | rxb_data[0] | U1 | 1. rxb_data[13] <br> 2. txa_addr[1] |
| H29 | GND | L26 | dmem_d[10] | P4 | $\begin{aligned} & \text { 1. rxb_clav } \\ & \text { 2. } \overline{\text { rxb_enb }} \end{aligned}$ | U2 | 1. rxb_data[12] <br> 2. txa_addr[0] |
| J1 | GND | L27 | dmem_d[1] | P5 | 1. $\overline{\mathrm{rxb}} \quad \mathrm{enb}$ <br> 2. rxb_clav | U3 | 1. rxb_data[11] <br> 2. rxa_addr[3] |
| J2 | txa_data[2] | L28 | dmem_d[11] | P25 | dmem_d[15] | U4 | 1. rxb_data[10] <br> 2. rxa_addr[2] |
| U5 | 1. rxb_data[9] <br> 2. rxa_addr[1] | W26 | 1. cmem_a[18] <br> 2. cmem_cs[1] | AB28 | cmem_d[1] | AE11 | $\overline{\text { reset }}$ |
| U11 | GND | W27 | cmem_a[1] | AB29 | GND | AE12 | inmo_d[15] |
| U13 | GND | W28 | cmem_a[0] | AC1 | GND | AE13 | GND |
| U15 | GND | W29 | VDD | AC2 | 1. txb_data[14] <br> 2. rxa_addr[4] | AE14 | inmo_d[7] |
| U17 | GND | Y1 | VDD | AC3 | txb_data[13] | AE15 | mem_clk_negative_o |
| U19 | GND | Y2 | txb_data[3] | AC4 | txb_data[12] | AE16 | inmo_d[1] |
| U25 | dmem_r/w | Y3 | txb_data[2] | AC5 | pllvss_300 | AE17 | global_tri_state |
| U26 | dmem_ $\overline{\text { bws }}[0]$ | Y4 | txb_data[1] | AC25 | cmem_d[4] | AE18 | inmo_a[14] |
| U27 | dmem_ $\overline{\text { bws }}[1]$ | Y5 | txb_data[0] | AC26 | cmem_d[10] | AE19 | inmo_a[10] |
| U28 | dmem_cs[0] | Y25 | cmem_cs[0] | AC27 | cmem_d[3] | AE20 | cpu_mode[0] |
| U29 | dmem_a[15] | Y26 | cmem_ $\overline{\text { bws }}[1]$ | AC28 | cmem_d[9] | AE21 | inmo_a[3] |
| V1 | GND | Y27 | cmem_ $\overline{\mathrm{bws}}[0]$ | AC29 | VDD | AE22 | N/C |
| V2 | rxb_soc | Y28 | cmem_r/w | AD1 | rxc_data[0] | AE23 | N/C |
| V3 | rxb_par | Y29 | VDD | AD2 | txb_soc | AE24 | cmem_a[15] |
| V4 | 1. rxb_data[15] <br> 2. txa_addr[3] | AA1 | GND | AD3 | txb_par | AE25 | N/C |
| V5 | 1. rxb_data[14] <br> 2. txa_addr[2] | AA2 | txb_data[7] | AD4 | 1. txb_data[15] <br> 2. txa_addr[4] | AE26 | cmem_d[13] |
| V25 | dmem_cs[3] | AA3 | txb_data[6] | AD5 | GND | AE27 | cmem_d[6] |
| V26 | dmem_cs[2] | AA4 | txb_data[5] | AD25 | GND | AE28 | cmem_d[12] |
| V27 | dmem_cs[1] | AA5 | txb_data[4] | AD26 | cmem_d[5] | AE29 | GND |
| V28 | dmem_a[16] | AA25 | cmem_d[0] | AD27 | cmem_d[11] | AF1 | GND |
| V29 | VDD | AA26 | cmem_a[4] | AD28 | mem_clk_o | AF2 | N/C |
| W1 | VDD | AA27 | cmem_a[3] | AD29 | mem_clk_i | AF3 | rxc_data[1] |
| W2 | 1. txb_clav <br> 2. $\overline{\text { txb_enb }}$ | AA28 | cmem_a[2] | AE1 | GND | AF4 | rxc_data[4] |
| W3 | 1. $\overline{\text { txb_enb }}$ <br> 2. txb_clav | AA29 | GND | AE2 | rxc_data[3] | AF5 | rxc_data[5] |
| W4 | txb_clk | AB1 | GND | AE3 | rxc_data[2] | AF6 | rxc_data[7] |

Table 11 - Pin Names Listed by Location (continued)

| Pin | Location | Pin | Location | Pin | Location |
| :---: | :---: | :---: | :---: | :---: | :---: |
| W5 | phyb_rx_led | AB2 | txb_data[10] | AE4 | N/C |
| W11 | GND | AB3 | txb_data[9] | AE5 | N/C |
| W13 | GND | AB4 | txb_data[8] | AE6 | rxc_data[6] |
| W15 | GND | AB5 | phyb_alm | AE7 | trst |
| W17 | GND | AB25 | N/C | AE8 | txc_data[1] |
| W19 | GND | AB26 | cmem_d[2] | AE9 | txc_data[5] |
| W25 | N/C | AB27 | cmem_d[8] | AE10 | 1. txc_clav <br> 2. txc_enb |
| AF14 | inmo_d[6] | AG11 | interrupt2 | AH8 | txc_data[4] |
| AF15 | mem_clk_positive_o | AG12 | inmo_d[13] | AH9 | txc_soc |
| AF16 | inmo_d[0] | AG13 | inmo_d[10] | AH10 | rxc_par |
| AF17 | inmo_wr_r/w | AG14 | inmo_d[5] | AH11 | interrupt1 |
| AF18 | inmo_a[13] | AG15 | inmo_d[3] | AH12 | inmo_d[12] |
| AF19 | inmo_a[9] | AG16 | inmo_rdy_ndtack | AH13 | inmo_d[9] |
| AF20 | inmo_a[6] | AG17 | inmo_ale | AH14 | inmo_d[4] |
| AF21 | inmo_a[2] | AG18 | inmo_a[12] | AH15 | inmo_d[2] |
| AF22 | cpu_mode[3] | AG19 | inmo_a[8] | AH16 | inmo_rd_ds |
| AF23 | N/C | AG20 | inmo_a[5] | AH17 | $\overline{\text { inmo_cs }}$ |
| AF24 | cmem_a[14] | AG21 | inmo_a[1] | AH18 | inmo_a[11] |
| AF25 | cmem_a[10] | AG22 | cpu_mode[1] | AH19 | inmo_a[7] |
| AF26 | cmem_par[0] | AG23 | cmem_a[17] | AH20 | inmo_a[4] |
| AF27 | cmem_d[14] | AG24 | cmem_a[13] | AH21 | inmo_a[0] |
| AF28 | cmem_d[7] | AG25 | cmem_a[9] | AH22 | cpu_mode[2] |
| AF29 | GND | AG26 | cmem_a[7] | AH23 | cmem_a[16] |
| AG1 | VDD | AG27 | cmem_par[1] | AH24 | cmem_a[12] |
| AG2 | txb_data[11] | AG28 | cmem_d[15] | AH25 | cmem_a[8] |
| AG3 | pllvdd_300 | AG29 | GND | AH26 | cmem_a[6] |
| AG4 | tms | AH1 | GND | AH27 | cmem_a[5] |
| AG5 | tdi | AH2 | GND | AH28 | GND |
| AG6 | rxc_soc | AH3 | VDD | AH29 | VDD |
| AG7 | txc_par | AH4 | tdo | AJ1 | VDD |
| AG8 | txc_data[3] | AH5 | tck | AJ2 | GND |
| AG9 | txc_data[7] | AH6 | $\begin{aligned} & \text { 1. rxc_clav } \\ & \text { 2. } \frac{\text { rxc_enb }}{} \end{aligned}$ | AJ3 | VDD |
| AG10 | txc_clk | AH7 | txc_data[0] | AJ4 | GND |


| Pin | Location |
| :---: | :---: |
| AF7 | rxc_clk |
| AF8 | txc_data[2] |
| AF9 | txc_data[6] |
| AF10 | 1. $\overline{\mathrm{txc}}$ _enb <br> 2. txc_clav |
| AF11 | mclk_src |
| AF12 | inmo_d[14] |
| AF13 | inmo_d[11] |
| AJ5 | VDD |
| AJ6 | 1. $\overline{\text { rxc_enb }}$ <br> 2. rxc_clav |
| AJ7 | GND |
| AJ8 | GND |
| AJ9 | GND |
| AJ10 | VDD |
| AJ11 | VDD |
| AJ12 | GND |
| AJ13 | inmo_d[8] |
| AJ14 | VDD |
| AJ15 | GND |
| AJ16 | VDD |
| AJ17 | inmo_a_das |
| AJ18 | GND |
| AJ19 | VDD |
| AJ20 | VDD |
| AJ21 | GND |
| AJ22 | GND |
| AJ23 | GND |
| AJ24 | cmem_a[11] |
| AJ25 | GND |
| AJ26 | GND |
| AJ27 | VDD |
| AJ28 | GND |
| AJ29 | VDD |
|  |  |

Table 11 - Pin Names Listed by Location (continued)

| Type | Input | Output | I/O | Power | Ground | N/C | Total |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| UTOPIA Port A | 21 | 12 | 12 |  |  |  | 45 |
| UTOPIA Port B | 21 | 14 | 10 |  |  |  | 45 |
| UTOPIA Port C | 12 | 12 | 2 |  |  |  | 26 |
| Clock recovery | 0 | 0 | 8 |  |  |  | 8 |
| CPU Bus | 26 | 3 | 16 |  |  |  | 45 |
| TDM Bus (H.100/H.110 bus) | 1 | 9 | 39 | 4 |  |  | 53 |
| Control memory | 0 | 23 | 18 |  |  |  | 41 |
| Data memory | 0 | 26 | 18 |  |  |  | 44 |
| Data and Control Memory Clocks | 3 | 3 | 0 |  |  |  | 6 |
| Test | 6 | 0 | 0 |  |  |  | 6 |
| PLL | 1 | 1 | 0 | 2 | 2 |  | 6 |
| Power |  |  |  | 40 |  |  | 40 |
| Ground |  | 1 |  |  |  | 88 |  |
| Miscellaneous |  |  |  |  |  | 88 |  |
| No Connects | $\mathbf{1 0 4}$ | $\mathbf{1 0 4}$ | $\mathbf{1 2 3}$ | $\mathbf{4 6}$ | $\mathbf{9 0}$ | $\mathbf{4 9}$ | $\mathbf{5 0 3}$ |
| Total: |  |  |  |  |  |  |  |

Table 12 - Pinout Summary


Figure 5 - PLL Pin Connections

### 4.0 Functional Description*

### 4.1 CPU Interface

The MT90503 CPU module provides an interface permitting programmability from an external microprocessor. The CPU module permits read/write access from MT90503: internal registers, internal and external memories. The CPU interface comprises of

- Direct Access Select (DAS) as the MSB bit concatenated with a 15 -bit address bus.
- 16-bit data bus.
- 2 interrupt signals.
- associated control signals.
- little endian format, unless otherwise specified.

The CPU interface can be configured to operate with either Intel or Motorola CPUs. The MT90503 supports both 8 -bit or 16 -bit data bus and multiplexed or non-multiplexed address/data pins. If the CPU is operating in 15-bit byte mode addressing with the LSB of its address bus as a byte field, then the inmo_a [14:0] pins of the MT90503 can be connected to the a[15:1] pins of the CPU. If both the MT90503 and the CPU are in 15-bit word mode addressing, then the inmo_a[14:0] pins should be connected to the a[14:0] pins of the CPU.

A reduced set of registers 'CPU Interface Registers' (0000h to 000Ah) are employed to optimize access time and to permit the CPU to execute indirect read/write accesses via these registers. The CPU also engages these registers to perform direct read/write accesses. The MT90503 and CPU timing relationship is described in section 9.0 on page 207.

The CPU Control Register ( 0100 h ) provides a software reset capability that allows the CPU to reset the MT90503 except for the CPU interface. The CPU interface can only be reset by a hardware reset.

### 4.1.1 CPU Interrupts

The CPU interface provides a programmable global interrupt capability. The interrupt signal names are 'interrupt1' and 'interrupt2', pins AH11 and AG11 respectively. Both interrupts have programmability to select their polarity (open collector drive) via registers 'interrupt1_conf' and 'interrupt2_conf' addresses 0224h and 0226h respectively. Interrupt1 accommodates a capability to program a minimum acceptable period between interrupts. The period is programmed in $\mu \mathrm{s}$ units via 'interrupt1_conf' register. This provides a 'frequency interrupt controller' facility and masks the assertion of further interrupts until the specified period has elapsed. The mask period will commence when the interrupt1_treated[15], register interrupt_flags address 0220h is set. When Interrupt2 is enabled it is always activated when an interrupt condition occurs. Interrupt pins are always tri-stated when inactive.

The operation of the CPU interrupt network is common for all modules. When an interrupt is asserted an interrupt flag is set to identify the module where the interrupt was generated. Each module has one or more Interrupt Enable Status Registers where a set interrupt enable bit identifies the source of the interrupt. On completion of the ISR the interrupt must be cleared as the interrupt will remain asserted until it is de-asserted by the user. All Interrupt Enable Status Registers have a symmetrical Status Register. Hence, the bit positioning of the interrupt enables and the associated status bits are identical.

### 4.1.1.1 Example Interrupt Flow

Upon the initialization of the Global Interrupt pins the following methodology is adopted to identify the source of the interrupt. For this example Interrupt2 is employed and the CPU module will be the source of the interrupt.

[^0]
### 4.1.1.1.1 Interrupt Initialisation

- Set interrupt polarity, register interrupt2_conf[15:14].
- Enable Interrupt2 for the CPU module, register interrupt2_enable[0] 022Ch. The MT90503 will generate an interrupt on interrupt2 according to the modules enabled in interrupt2_enable.
- Set the individual CPU interrupt sources by enabling the respective bits in the 'status0_ie' 0104h register. Within the `status0_ie' register there are two possible interrupt sources: internal_read_timeout_ie and inmo_read_done_ie. In the MT90503 Register Description the interrupt bits are labelled IE (Interrupt Enable) in the 'Type' column. This register offers the facility to mask/ disable unwanted interrupts.


### 4.1.1.1.2 Interrupt Servicing

When interrupt2 is asserted ('interrupt2' pin):

- Read the interrupt flags to ascertain the module raising the interrupt. The CPU module interrupt flag is located in register interrupt_flags[0] 0220h, this bit is named cpureg_interrupt_active.
- If the cpureg_interrupt_active bit is set, locate the source of the CPU interrupt by reading the 'status0' register 0102h, either internal_read_timeout and/or inmo_read_done.
- The associated status register 'status0' 0102h contains internal_read_timeout and inmo_read_done bits. Therefore, to de-assert the interrupt the user must write a 1 to register 0102h bits 3 or 4, internal_read_timeout and inmo_read_done respectively. Only then will the interrupt be de-asserted.


### 4.1.2 Intel/Motorola Interface

The MT90503 CPU interface supports both Intel and Motorola microprocessors, in both 8 -bit or 16 -bit data bus and multiplexed or non-multiplexed address/data pins. The MT90503 supports 8 MB of addressable space, therefore indirection addressing is necessary. The microprocessor interface directly addresses five control words, used for indirection accessing. The indirection register contents are shown in tables 13 to 17 inclusively. The timing relationship pertaining to the CPU Interface Registers and Extended Access is defined in section 9.0 on page 207.

| 0000 h | Control Register |
| :--- | :---: |
| 0004 h | Read/Write Data Register |
| 0008 h | Address High Register |
| 000 Ah | Address Low Register |
|  |  |


| $\underset{[\overline{3}: 0]}{\text { cpu_mode }}$ | Interface Type | address pins | data pins | direct_access | ale |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0000 | Intel, 16 bit data bus, non-multiplexed | inmo_a[14:0]** (word address) | inmo_d[15:0] | inmo_a_das | inmo_ale* |
| 0001 | Intel, 16 bit data bus, multiplexed | inmo_d[15:1] <br> (word address) | inmo_d[15:0] | inmo_a_das | inmo_ale* |
| 0010 | Intel, 8 bit data bus, non-multiplexed | inmo_a[14:0] (byte address) | inmo_d[7:0] | inmo_a_das | inmo_ale* |
| 0011 | Intel, 8 bit data bus, multiplexed |  <br> inmo_d[7:0] <br> (byte address) | inmo_d[7:0] | inmo_a_das | inmo_ale* |
| 0100 | Motorola, 16 bit data bus, non-multiplexed | inmo_a[14:0] (word address) | inmo_d[15:0] | inmo_a_das | inmo_ale* |
| 0101 | Motorola, 16 bit data bus, multiplexed | inmo_d[15:1] <br> (word address) | inmo_d[15:0] | inmo_a_das | inmo_ale* |
| 0110 | Motorola, 8 bit data bus, non-multiplexed | inmo_a[14:0] (byte address) | inmo_d[7:0] | inmo_a_das | inmo_ale* |
| 0111 | Motorola, 8 bit data bus, multiplexed |  <br> inmo_d[7:0] <br> (byte address) | inmo_d[7:0] | inmo_a_das | inmo_ale* |
| 1xxx | Reserved |  |  |  |  |
| * The inmo_ale pin is interpreted in all modes. However, it is not necessary in the non-multiplexed modes and can be tied to VCC. **The address placed on the inmo_a[14:0] pins is a word address in 16 -bit mode and a byte address in 8 -bit mode. The address, when placed on the inmo_d pins, is always a byte address. |  |  |  |  |  |

Table 13-CPU Interface Mode Selection

| Field | Bit | Type | Reset | Description |
| :--- | :---: | :---: | :---: | :--- | :--- |
| read_burst_length | $6: 0$ | RW | 01 h | Number of words to prefetch. <br> Setting <br> Words <br> 00 h |

Table 14-Control Register (0000h)

| write_enable | $13: 12$ | RW | Oh | Active high write enables. <br> $00=$ read access. <br> $01=$ write to lower byte. <br> $10=$ write to upper byte. <br> 11 = write to entire word. <br> This field is ignored for: extended direct reads and all byte wide <br> extended direct accesses. |
| :--- | :---: | :---: | :---: | :--- |
| extended_parity | $15: 14$ | RW | Oh | Read/Write Parity bits. |

Table 14-Control Register (0000h)

| Field | Bit | Type | Reset | Description |
| :---: | :---: | :---: | :---: | :--- |
| extended_data[15:0] | 15:0 | RW | 0000 h | The extended indirect read/write data word register. Invalid <br> for extended direct access. |

Table 15 - Read/Write Data Register (0004h)

| Field | Bit | Type | Reset | Description |
| :--- | :---: | :---: | :---: | :--- |
| extended_a[32:20] | $12: 0$ | RW | 000 h | Upper extended address [32:20]. |
| Reserved | $15: 13$ | RO | $0 h$ | Reset to 000. |

Table 16 - Address High Register (0008h)

| Field | Bit | Type | Reset | Description |
| :---: | :---: | :---: | :---: | :---: |
| extended_a[19:4] | 15:0 | RW | 0000 h | Lower extended address [19:4]. In extended direct <br> addressing, bits 19:16 are employed for 16 bit data bus <br> and bits 19:15 are employed for 8 bit data bus. |

Table 17 - Address Low Register (000Ah)

### 4.1.2.1 Extended Indirect Accessing

Extended Indirect Accessing solely employs the registers 0000h to 000Ah to access the 8 MB of addressable memory space.

Synopsis: the user writes the access address to registers 0000h, 0008h and 000Ah. Then the MT90503 will read/write to that address and fetch/place the data value from/to register 0004h. For all extended indirect accesses the INMO_A_DAS bit will be held low.

### 4.1.2.1.1 Extended Indirect Writes

The following steps must be executed to perform an extended indirect write:

1 Write the upper address, extended_a[32:20], to register 0008h. This write may be not be required if previous value holds true.

2 Write the lower address, extended_a[19:4], to register 000Ah. This write may be not be required if previous value holds true.

3 Write the write data, extended_data[15:0], to register 0004h. This write may be not be required if previous value holds true.

4 Write write_enable, extended_parity, access_req='1' and extended_a [3:1] in a single access to register 0000h.

5 Read the access_req bit located in the Control Register[8] to determine when the write has completed.
The software will set access_req [8] in register 0000h (Step 4 above) and the hardware will reset it when the data write has completed. Therefore, the user can poll this bit to determine when the data write has completed.

### 4.1.2.1.2 Extended Indirect Reads

1 Write the upper address, extended_a[32:20], to register 0008h. This write may be not be required if previous value holds true.

2 Write the lower address, extended_a[19:4], to register 000Ah. This write may be not be required if previous value holds true.

3 Write write_enable = 00, access_req='1' and extended_a [3:1] in a single access to register 0000h.
4 Wait until access_req is cleared, then read the data from the data field extended_data[15:0], register 0004h.
5 Optional parity check may be ascertained by performing a read on the extended_parity[15:14], register 0000h.

The software will set access_req[8] register 0000h and the hardware will reset it when the data is ready to be read from register 0004h.

### 4.1.2.2 Extended Direct Accessing

Extended Direct Accessing employs the high and low address registers to perform page addressing. The address within the page is provided directly by the CPU address bus. Similarly the data is fetched/placed directly on the CPU data bus.

Synopsis: the user writes the access address to registers 0008h and 000Ah but this performs only the page addressing. Upon assertion of the address within the page the MT90503 will read/write the data with respect to that address. The INMO_A_DAS bit is set when the data read/write occurs. When operating the CPU interface in direct mode with a 16-bit data bus, extended_a[19:16], are employed for the lower address word register

000Ah. However, when operating the CPU interface in direct mode with an 8-bit data bus bits, [19:15] are used for the lower address word.

### 4.1.2.2.1 Extended Direct Writes

1 Write the upper address, extended_a[32:20], to register 0008h. This write may be not be required if previous value holds true.

2 Write the lower address extended_a[19:16] or [19:15] to register 000Hh. The remaining bits [15:4] or [14:4] are ignored. This write may be not be required if previous value holds true.

3 Write write_enable[13:12] (where applicable) and extended_parity[15:14]. The extended parity write is optional.

4 Write the data value to the address within the corresponding memory page with the INMO_A_DAS bit set.

### 4.1.2.2.2 Extended Direct Reads

1 Write the upper address, extended_a[32:20], to register 0008h. This write may be not be required if previous value holds true.

2 Write the lower address, extended_a[19:16] or [19:15], to register 000Ah. The remaining bits [15:4] or [14:4] are ignored. This write may be not be required if previous value holds true.
3 Assert the lower address within the memory page and fetch the read data with INMO_A_DAS set.
4 An optional read may be performed to obtain the parity values, extended_parity[15:14] register 0000h.

### 4.1.3 MT90503 Reset Procedure

The following reset procedure is required to power-up the MT90503. The reset procedure must be adhered at power-up employing the reset pin. Post power-up, the reset procedure can be performed from step 3.

1 Assert the reset pin for at least 1000 MCLK cycles.
2 De-assert the reset pin. All accesses in the remaining reset procedure will employ indirection.
3 Initialise the mem_clk_* bits and write_cache_enable bit via the Control Register bits [13:10], address 0100h.

4 Write the mclk_src frequency at register led1[6:0] and LED Flash Frequency at registers led1[15:7], register address 0120h. Write the LED Flash Frequency units (i.e. ms or $\mu \mathrm{s}$ ) at register led2[0], address 0122h.

5 Initialise the PLL via register pll_conf, address 0128h.
6 Set bits nreset_* in the Control Register bits [7:0], address 0100h.
7 Initialize the data and control memory types via registers 0240h, 0242h, 0244h, 0248h, 024Ah, and 024Ch
8 Select UTOPIA clocking methodology at register addresses 0230h, 0232h, 0234h and 0236h.
9 Configure the interrupts' active_level for interrupt1 and interrupt2 in register 0224h and 0226h.

### 4.2 TDM Module

The general architecture of the TDM bus consists of three main elements.

- The TDM bus interface which is coupled directly to the bus pins and manages the timing requirements of the H. 100/H. 110 interface.
- The datapath management for transporting the bytes from the TDM bus to the circular buffers (located in the external data memory).
- The TDM clocking mechanism, this enables the MT90503 to be H. 100/H. 110 bus master capable and communicates with the clock recovery module.


### 4.2.1 TDM Bus Interface

The MT90503 TDM Module interfaces with all 32 data streams of the H.100/H. 110 bus. The maximum data rate of 8.192 Mbps determines the total bus capacity of 4096 TDM TSST. The MT90503 can process up to: 2048 transmitting TDM TSST and 2048 receiving TDM TSST within one frame of $125 \mu \mathrm{~s}$.

One less TDM channel is carried for each CAS channel that is desired. If all TDM channels have CAS, then 1024 transmit and 1024 receive channels is the limit. The MT90503 can drive out data on any particular TSST (Time Slot/Stream), or read in data from any particular TSST.

Each individual time slot or DS0 on the $\mathrm{H} .1 \times 0$ bus is assigned by a unique TSST number, based on the following equation:

$$
\text { TSST }= \begin{cases}\text { timeslot * } 32+\text { stream } & \text { for } 8 \mathrm{MHz} \text { streams } \\ \text { (timeslot * } 2+1 \text { ) } 32+\text { stream } & \text { for } 4 \mathrm{MHz} \text { streams } \\ \text { (timeslot * } 4+3 \text { ) } 32+\text { stream } & \text { for } 2 \mathrm{MHz} \text { streams }\end{cases}
$$

The timeslot ranges from 0 to 31 for 2 MHz streams, from 0 to 63 for 4 MHz streams, or from 0 to 127 for 8 MHz streams.

The 16 lowest data streams are capable of running at a data rate lower than 8.192 MHz . The streams are grouped in fours and each quartet must run at the same data rate. Streams [3:0], [7:4], [11:8] and [15:12] can each run at $2.048,4.096$ or 8.192 MHz as a group. This allows backward compatibility with older, slower TDM buses. In the reduced rate, the data is still latched using the CT_C8_A or _B clock-edge, but using every second or every fourth clock-edge. The streams numbered [31:16] must always run at 8.192 MHz .

The TDM bus can also be configured to use only 4, 8, or 16 streams. This configuration requires less processing of TDM bytes and, also allows a reduced mclk speed. The 4 -stream mode is useful for conducting tests, while the 16 -stream mode allows mclk speed to be cut in half while still allowing half the bandwidth of a full H. 100/H. 110 bus.

When the H. $100 / \mathrm{H} .110$ bus is carrying CAS signalling bits, every even stream carries regular TDM data bytes, and is associated with an odd stream that carries CAS bits and the indication of the beginning of the multiframe. For TSSTs that are input to the MT90503, the multiframing is provided by the framer. For TSSTs that are generated by the MT90503, the output multiframing is provided by the MT90503; there is one multiframe which is common to all the T1 channel signals and one multiframe which is common to all the E1 channel signals. On the input, the multiframing is independent for each TSST and the MT90503 synchronises all of them to its multiframe.


Figure 6 - TDM Serial to Parallel/Parallel to Serial (SPPS) Converter

Each TSST can be used independently for CAS or regular TDM bytes. For example, if time slot 0 on stream 0 requires signalling bits, then they will be carried on time slot 0 , stream 1 . However, time slot 1 stream 1 can be used to carry normal TDM data, even if the time slot preceding it is used for signalling. T1 and E1 channel formats can be supported simultaneously. The multiframe bit's polarity and position is programmable and global, and will be accepted at both the input and output. The CAS bits occupy bits [7:4] or [3:0] of the CAS byte, depending on the position within the byte of the multiframe indicator. The only difference between CAS support in E1 or T1 mode is the number of frames in the multiframe. T1 contains 24 frames within its multiframes while E1 contains 16. The CAS bits are latched at the input when the multiframe bit is active, and sent out on the output side during the entire multiframe. When the multiframe bit is at an active polarity, the associated TDM byte is the first byte of the multiframe.


Figure 7 - CAS and MFS Transport on the TDM Bus

Figure 7, CAS and MFS Transport on the TDM Bus shows the direct interface of the MT90503 with the TDM bus.

The TDM interface is capable of looping back up to 2048 TSSTs with a latency of two TDM frames, or $250 \mu \mathrm{~s}$. This loopback is indicated by the control structures associated to each TSST. The loopback allows any two time-slots and streams to be connected together, and work independently of the speed of each of the streams.

The MT90503 can interface with 32 data streams as mentioned in previous paragraphs. Since each TSST can be used to either send or receive data, the MT90503 must tri-state the CT_D pins that are receiving data. To avoid overlap between a sending and a receiving TSST, this must be done before the next one begins.

### 4.2.2 TDM Bus Clocking Mechanism

The MT90503 can operate as either a primary or secondary H.100/H. 110 master clock source. The H. 100/ H. 110 clock (CT_C8) is driven by either CT_C8_A or CT_C8_B. The MT90503 can generate the primary bus clock and all its associated bus clocks (ct_fr_comp, sclk, sclkx2, CT_C16-, CT_C16+, CT_C2 and CT_C4) by using the CT_NETREF signal. The CT_NETREF may be received from another H.100/H. 110 bus compatible device or it can be determined by the clock recovery module. In master mode, the MT90503 will generate all the compatibility clocks that are necessary for communicating with MVIP and SCSA TDM interfaces. These include the FR_COMP, C2, C4 and C16 signals for communicating with MVIP buses, and the SCLK and SCLKx2 signals for communicating with Signal Computing System Architecture SCSA buses.

The TDM module monitors both CT_C8_A and CT_C8_B signals for clock properties and failure. In every operating mode, both the CT_C8_A ${ }^{-}$and $C T \_C 8 \_\bar{B}$ signals are monitored for clock failure. There are two methods of detecting a clock failure on the bus. Firstly, if the rising edge of the clock does not appear within $\pm 35$ ns window of the expected period, the clock is flagged as being invalid. Secondly, if a single frame does not contain exactly 1024 clock cycles, then the clock will fail.

The MT90503 can be configured to switch to the backup master clock (CT_C8_A or CT_C8_B), upon the detection of a master clock failure. The MT90503 can be programmed as a secondary master and therefore switch to the backup master clock when the external primary master clock source has failed. If the MT90503 switches from a secondary master to master it will re-synchronise the TDM module with the backup clock and generate the compatibility clocks. The MT90503 continually monitors both CT_C8_A and CT_C8_B clocks for a failure condition. The MT90503 can be programmed to automatically switch to its backup clock in case of a failure.

The MT90503 is an H. 100/H. 110 master-capable device and, therefore, is able to generate both CT_C8 clocks and CT_FRAME and compatibility signals. As long as the primary signals on the bus are valid, the MT90503 will synchronise its output to them using a PLL. The MT90503 can also generate the CT_C8 and CT_FRAME signals independently of the PLL, using the local 16 MHz clock. In this case, the output signals have no phase relation with those present on the bus.

When the TDM bus interface module has written the received data to the circular buffers, it increments the Time-Slot Memory Pointer that is sent to the TX_SAR.

### 4.2.3 TDM Datapath

When writing to the external data memory, no return value is expected. However, when reading data, a return value, upon storing the data into an RX TSST, is expected and is required to be written back to the structure memory. The underrun count, which may or may not have been updated, is always written back along with all the other bits of the second word.
When the data is received at the time slot memory, it is sent via the TDM datapath to the external data memory. According to which time slot is currently being used, the TDM control structures are read.


Figure 8 - TDM Data Path Controller
When accesses are performed from the TDM channel association memory, the write back to the RX half of the circular buffer is also performed. When this write back is done, the underrun detection bit is written to ' 0 ', and a null-octet is inserted in place of the old data byte.

The TDM datapath module can be configured to work with $32 \mathrm{H} .100 / \mathrm{H} .110$ streams, 16 -streams, 8 -streams, or 4 -streams in order to allow lower clock speeds on the MT90503. The bandwidth requirements for supporting each of the above configurations are described in the Bandwidth section.

### 4.2.4 TDM Channel Association Structures

TDM channel association structures are located at $0 \times 8000$ to $0 \times B F F E$ in internal memory. Each structure corresponds to a TDM channel by its TSST order.

### 4.2.4.1 Non CAS Operation



Figure 9 - TDM Channel Association Structures: TX Channel non-CAS mode
The control structure in Figure 10, TDM Channel Association: RX Channels (Non CAS mode) on page 37 indicates the mode and options selected for the TDM channel. In the general TX structure, the RBW (RX Byte

Write) bits provide the format of the byte that is to be written over the RX byte in the circular buffer. While it is possible to write over the low-byte, there are three possibilities provided if writing over the byte is chosen: writing a null byte or generating one of two silence patterns. The detailed description of the dual-direction buffer is provided in the TDM Circular Buffers section.

| b15 b14 b13 b12 b11 b10 b9 b8 b7 b6 b5 b4 b3 b2 b1 b0 |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| +0 | TX/RX Circular Buffer Address and Size |  |  |  |  | OE |
| +2 | Mode | C | U | 1 | UR Cour |  |

TX/RX Circular Buffer Address and Size: Address and size of the circular buffer in the data memory from which data bytes will be
OE. ${ }^{\wedge}$ Voice Stream Output Enable. This bit, when set, enables the driving of the Voice Data on the stream.
Mode: Channel Mode of operation.
"0000"=Normal PCM;
"1000"=Low Latency Loopback
Reserved others=Reserved.
C: Cut VC Status Enable. When '0', the UR Count field is freerunning( 256 consecutive underruns will béronsidered as a Cut VC). When ' 1 ', the UR Count field is used as a consecutive underrun indicator that stops incrementing at 255.
$\mathbf{U}$ : Underrun Status Enable. When this bit is set and an underrun is detected, the status counters and bits will report this event.
UR Count 8 bit underrun count. When the $\mathbf{C}$ bit is cleared, acts as a freerunning underrun error counter. When the $\mathbf{C}$ bit is set, it is used as a cut VC detector. In this mode, the counter is cleared each time a valid byte is received. It increments (to a maximum of FFh) each time a byte underrun is detected. When a transition from FEh to FFh occurs, the Cut VC status register, counter and id fields are updated. Must be initialized to FFh by software before enabling VC when $\mathbf{C}$ bit is set. Must be initialized to 00h by software before enabling VC when $\mathbf{C}$ bit is cleared.

Figure 10 - TDM Channel Association: RX Channels (Non CAS mode)

In the receive direction, the control structure retains some of the recurrent features from the transmit structure. The TX circular buffer/size field is the same, which is normal given that the TX and RX circular buffers are common. The information on the address and size is encoded in the same way as in the TX structure.

For the reception structure, the high mode bit codes whether the RX channel is transmitting a channel received from ATM or if it is retransmitting information taken from the TDM bus and written into a circular buffer. If the high Mode bit is ' 0 ', the channel is ATM; if it is ' 1 ', the channel is TDM. By supporting a low latency TDM loopback, the MT90503 conforms to the H.100/H. 110 specification.

The $C$ and $U$ bits serve as status enable bits. These bits are R/W, and tell the register module whether the register counters and status bits should report errors on this VC. The $U$ bit is the underrun status enable: when this bit is set and an underrun is detected, the TDM register module will increment the global underrun counter and set the underrun detect status bit.

The $C$ bit serves at the cut VC detect status enable. When this bit is at ' 0 ', the UR count acts as a free running 256 -underrun counter. In this case, it serves to compile the total number of underruns that have occurred. When the $C$ bit is at ' 1 ', then the UR count serves as a detector for cut VCs. The counter resets to 0 each time a valid TDM byte is received. If an underrun is received, the counter is incremented by one. If the counter ever increments from 254 to 255 , then 255 consecutive underruns have been received, indicating there is 32 ms of absent data. This is interpreted as a cut VC and will generate an interrupt.

Cell-loss integration periods of greater than 32 ms can be supported by software. Upon UR Count reaching 255 , and subsequent interrupt to the CPU, the software can check that the count has not returned to zero (which happens if the cells start to arrive again) at some interval equivalent to the cell-loss integration period.

### 4.2.4.2 CAS Operation



TXIRX Circular Buffer Address and Size: Address and size of the circular buffer in the data memory to which data bytes will be written. I: Initialized Bit. Written by ' 0 ' by software, set by hardware when the channel starts being treated.
Mode: Channel Mode of operation.
"0100"=E1 Strict Multiframing;
"0101"=E1 FASTCAS;
"0110"=T1 Strict Multiframing;
"0111"=T1 FASTCAS;
others=Reserved.
RBW: RX Byte Write. Selects byte that will be written in the low bytes of the data memory word where the TX TDM bytes are written.
"00"=Do not write over the low byte;
" 01 " $=$ Write a null byte (usually FFh);
"10"=Write silence pattern A;
"11"=Write silence pattern B.
CASWE CAS write enable.
"x0"=Leave CAS bits in RX Circular Buffer untouched;
"x1"=WriteRX CAS URin RX Circular Buffer in case of underrun;
" $0 x$ "=Write CAS bits present on the associated odd stream to the TX Circular Buffer;
" 1 x "=WriteTX CAS Forcefield in order to bypass the CAS bits from the associated odd stream.
TX CAS Force TX CAS value that must be written to bypass the CAS bits from the associated odd stream.
EI: CAS Enable Ignore. If set, the CAS Enable bit (SFS Bit) will be ignored. Instead, the CAS will be latched at one point for each 16/24 consecutive bytes (for E1 and T1 respectively).
RX CAS UR RX CAS value that will be sent in case of an underrun situation on the RX_SAR side. CM: Cas Monitor. When ' 1 ', any change in the TX CAS value will be reported to the CPU.
Last TX CAS Last Value of the TX CAS received from the TDM bus.
FS Offset Frame Offset that must be added in order for internal and external multiframes to coincide. Initialize to "00000" by software.

Reserved

Figure 11 - TDM Channel Association: TX Channels (CAS mode)

The RX Byte Write bits remain the same to allow the insertion of null patterns as for regular channels. The Mode bits, indicate the possible configurations and combinations of multiframing and CAS insertion: "0100" is E 1 , " 0110 " is T 1 , and toggling the lowest bit of either of the numbers indicates that the FASTCAS method of transmitting the data and CAS bytes is being used. FASTCAS is used to lower the latency of Circuit Emulation cells. When using the FASTCAS mode, the multiframe pointer is not used and multiframe integrity is therefore not maintained.

In addition, special fields for CAS are added in the two additional words of the structure. The CAS WE (CAS Write Enable) bits are used to determine if the CAS is written back.

The TX_CAS field is used so the CPU can insert values of CAS, in place of values from the TDM bus; the RX CAS value is used to compensate for underruns.

The Last TX CAS field is used to detect when the CAS value received on the TDM bus has changed. When a CAS value is received, it is written back to the Last TX CAS field in the structure. Whenever a new value arrives, it is compared to the last value. If there is a difference, a CAS change signal is sent to the registers, and the new CAS value is written to the CAS change buffer in the external memory.

The Frame Offset field is used to store the offset between the internal and external multiframes in the TX direction. When the TDM bus first obtains an external multiframe, it writes the offset field to the correct value between 0 and 23. Subsequently, each time that a byte is written to the circular buffer, the offset field is read, and an offset is established between the internal TDM pointer and the pointer that will be used.

```
b15b14b13b12b11b10 b9 b8 b7 b6 b5 b4 b3 b2 b1 b0
```



```
TX/RX Circular Buffer Address and Size: Address and size of the circular buffer in the data memory from which data bytes will be
OE. Voice Stream Output Enable. This bit, when set, enables the driving of the Voice Data on the even stream.
Mode: Channel Mode of operation.
"0100"=E1 Strict Multiframing;
"0101"=E1 FASTCAS;
"0110"=T1 Strict Multiframing;
"0111"=T1 FASTCAS; others=Reserved.
C: Cut VC Status Enable. When ' 0 ', the UR Count field is freerunning ( 256 consecutive underruns will be cosidered as a Cut VC). When ' 1 ', the UR Count field is used as a consecutive underrun indicator stops incrementing at 255.
U: Underrun Status Enable. When this bit is set and an underrun is detected, the status counters and bits will report this
CU: CAS Underrun Enable. When this bit is set and an CAS underrun is detected, the status counters and bits will report this
UR Count 8 bit underrun count. When the \(\mathbf{C}\) bit is cleared, acts as a freerunning underrun error counter. When the \(\mathbf{C}\) bit is set, it is used as a cut VC detector. In this mode, the counter is cleared each time a valid byte is received. It increments (to a maximum of FFh) each time a byte underrun is detected. When a transition from FEh to FFh occurs, the Cut VC status register, counter and id fields are updated. Must be initialized to FFh by software before enabling VC when \(\mathbf{C}\) bit is set. Must be initialized to 00h by before enabling VC when \(\mathbf{C}\) bit is cleared.
\(\mathbf{C M}\) : CAS Monitor. When ' 1 ', any change in RX CAS value will be reported to the CPU. When an underrun occurs, a CAS change will never be reported.
CO: CAS Output Enable. This bit, when set, enables the driving of the CAS value out on the associated odd TDM steam.
Last RX CAS This is the value of the last RX CAS sent on the TDM bus. This value is not written if an underrun has occured for this CAS value,
ME: Manual CAS Insert. When ' 1 ', the Man RX CAS is sent instead of the CAS contained in the memory.
Man RX CAS CAS Value that is sent onto the TDM bus if ME = ' 1 '. \(\quad\) Reserved
```

Figure 12 - TDM Channel Association: RX Channels (CAS mode)

The channel association structure for multiframe channels is slightly different, with only one bit being positioned differently in the structure. While the circular buffer address and size field is still present and functions in the same way, the mode bits are coded in the same way as they are in the transmit structure for multiframe channels. " 0100 " is E1, " 0110 " is T1, and toggling the lowest bit of either of the numbers indicates that the FASTCAS method of transmitting the data and CAS bytes are employed.

Four new fields are added to allow for CAS management on the RX side with multiframing:

- Last RX CAS value that indicates the previous CAS bits received on ATM.
- The CAS monitor bit indicates if a change in the value of CAS is to be reported by the module.
- Man RX CAS is the value of CAS that can be inserted in the place of the one received. This value is only used if the ME bit is set.
- ME bit Manual CAS insert.

For both TX and RX channels, a change in the value of CAS on a channel where the CAS monitor bit is high will cause the new value to be written to the CAS change buffer in the external memory. This buffer contains all the values that have been detected as different from the previous value, along with the 12 -bit time slot and stream value of the TDM channel to which they are associated.

|  b15 b14 b13 b12 b11 b10 b9 b8 <br>  b7 b6 b5 b4 b3 b2 b1 b0 <br>  TSST Number On Which CAS Changed  New CAS      |  |  |
| :---: | :---: | :---: |
|  |  |  |

Figure 13 - CAS Change Structure in Control Memory
The Circular Buffer Address/Size field is 15 bits and the addresses are 14 bits each, an indirect method of indicating the size is used. This is done by coding the size by the first ' 1 ' that is encountered by reading the field starting from the left. For example, if the lowest bit of the field is ' 1 ', the buffer is 128 bytes (mapped on a 256 byte boundary). If the lowest bit was ' 0 ' but the second-lowest bit is ' 1 ', then the buffer is 256 bytes (mapped on a 512 byte boundary), etc.


Figure 14-TX/RX Circular Buffer and Size Field

### 4.2.5 TDM Circular Buffers

The TDM module writes the data received from the bus to circular buffers that are assigned on a per-channel basis. The circular buffers' size is variable and is specified in two places, i.e., once in the TDM channel association structure for which each channel has an entry, and once in the TX_SAR or RX_SAR structure which reads or writes to that channel.

The TX part of the circular buffer (data read from the TDM bus and written to the circular buffer) is contained in bits 15:8 of each circular buffer word. The RX part of the buffer is contained in the remaining bits, 7:0. This alignment allows for underrun detection and null-pattern insertion in the case of underruns. It also allows the detection of cut VCs. A cut VC is indicated when 256 consecutive underrun bytes are detected.

In the non-multiframing mode, the circular buffers are used as a single block. This means that for each channel, up to 1024 bytes of data can be stored in the channel's circular buffer. On the RX side, up to 128 ms of data can be stored for retrieval; therefore, at $\mathrm{n}=1$ with a packet size of 48 bytes, up to $\pm 61 \mathrm{~ms}$ of CDV can be compensated.

In multiframing mode, each 32-byte division of the circular buffer is used to store the information of one multiframe. In T1 operation, the first 24-bytes of the 32-byte division are used to remember the TDM data from the multiframe. In E1 operation, the first 16-bytes are used for the same purpose. In addition, in byte 31 of each division, the CAS value of the multiframe is stored. In this way, the multiframe portion of the TDM pointer is valid for the TDM data and for the CAS data; only the lowest bits need to be changed to " 11111 " in order to point to CAS.

When write backs are performed, a null byte or one of two silent patterns (background noise) can be inserted in the RX byte. The silent patterns are read from buffers of up to 64 KB . These silent patterns can be used to generate noisy silences.


Figure 15 - Silent Pattern Buffer A/B in Control Memory


Figure 16 - TDM Circular Buffer (one MultiFrame in T1 mode)

The CAS bits are always kept in the highest four bits of the last byte.
In addition to storing the data in the circular buffers, the parity bits of the data memory are used to detect underruns generated on the ATM link. When the RX_SAR writes data to the circular buffers, it writes a ' 1 ' in the parity bit of the memory's low byte. When the TX_SAR writes data bytes in the buffer, it writes a ' 0 ' to the parity bit of the memory's low byte. Therefore, when the TDM reads the data from the circular buffer, it expects to read a ' 1 ' from the parity bit. If it reads a ' 0 ', this means that an underrun has occurred because not enough cells are being received, or there is too much CDV.


Figure 17 - TDM Circular Buffer (one Super Frame in E1 mode)

### 4.2.6 TDM Circular Buffer Pointers

The convention that a write pointer points to a byte that is currently being written. In like manner, a read pointer points to a byte that is currently being read. Therefore, a read and write pointer to the same buffer must never cross, or be equal to one another.

### 4.2.6.1 Non-multiframing mode

The global TDM_write_pointer points to the byte from the TDM bus presently being written to the circular buffers. As in all modes, the TDM_read_pointer is incremented before the TDM_write_pointer, because bytes that come from the TDM bus are written after the frame is completed. Meanwhile, bytes read from the data memory and sent onto the bus are already transmitted by the time the frame is complete.

The TDM_write_pnt sent to the TX_SAR is the same as global TDM_write_pointer. The TDM_read_pnt sent to the RX_SAR is one frame ahead of the global TDM_read_pointer. Since the TDM_read_pnt is one more than the global TDM_read_pointer, and the TDM_read_pointer is incremented before the TDM_write_pointer, the TDM_read_pnt will always have a lead of either one or two over the TDM_write_pnt.


Figure 18 - TDM Circular Buffers (Normal mode)
Figure 18 TDM Circular Buffers (Normal mode) shows that very few of the positions within the circular buffer are invalid for the SAR pointers.

### 4.2.6.2 E1/T1 Multiframing (standard)

The TDM pointers change to respect the multiframing standards required by E1 and T1 operation. The pointer has two halves. One half indicates the number of the multiframe that the TDM is writing or reading; the other half is the number of the frame within the multiframe. In the T1 mode, this fraction goes from 0 to 23 , while in the E1 mode it ranges from 0 to 15 . The same global standards of pointer definition still apply as the pointer still points to an invalid byte.

The TDM_write_pnt has a global value, which is the value of the internal multiframe pointer. However, due to the multiframing offset, the TDM_write_pnt can vary by almost an entire frame. Therefore, if the global TDM_write_pointer points to 0.0 (Multiframe 0 , frame 0 ) the individual TDM channels may be written anywhere from 0.0 to 0.23 in the T1 mode, and from 0.0 to 0.15 in the E1 mode. The TDM channels may be written almost an entire frame ahead of the global pointer. The pointer sent to the TX_SAR is the same as the global TDM_write_pointer.

The TDM_read_pnt is two multiframes ahead of the TDM_write_pnt. The need for this lead arises from the write backs and from CAS. Since the TDM_write_pnt can be writing up to 15 frames ahead, a lead of at least one multiframe must be given. In addition to this, the CAS must be considered. CAS is written during the multiframe of the tx, which can be up to one multiframe ahead of the global pointer. When CAS is read from the $r x$ side it is read at the beginning of the multiframe. Therefore, to ensure that there is always a difference of at least one in the CAS pointers, the delta between the TDM pointers is established at two multiframes.

Furthermore, the read pointer that is sent to the RX_SAR is almost one multiframe ahead. This is necessary because CAS is sent at the end of the multiframe on the ATM side, and is read at the beginning of the multiframe on the TDM side. Therefore, to ensure that the CAS of the multiframe is written by the RX_SAR before the TDM reads it, the pointer passed to the RX_SAR is incremented by one multiframe minus one frame. Therefore, if the TDM_read_pointer would have been 2.5 (two multiframes, five frames), the TDM_read_pnt passed to the RX_SAR is 3.4 (three multiframes, four frames).


Figure 19 - TDM Circular Buffer (Complete Buffer in E1 mode, Strict Multiframing)

Figure 19 TDM Circular Buffer (Complete Buffer in E1 mode, Strict Multiframing) shows the TDM circular buffer in the E1 mode. The buffer is only half used for the TDM data. The CAS is located at the end of the buffer, and the write and read pointers are offset by a full three multiframes.


Figure 20 - TDM Circular Buffer (Complete Buffer in E1 mode, FASTCAS)

Figure 20 TDM Circular Buffer (Complete Buffer in E1 mode, FASTCAS) is an E1 circular buffer in the FASTCAS mode. Although the TDM read and write pointers are still significantly offset, the delay has been reduced to nearly zero. This is because the rx_sar_write_pnt and the TDM_read_pnt need only be offset by a single byte, as in non-multiframing mode.


Figure 21 - TDM Circular Buffer (Complete Buffer in T1 mode, Strict Multiframing)

Figure 21 TDM Circular Buffer (Complete Buffer in T1 mode, Strict Multiframing) shows a T1 circular buffer in standard multiframing mode. The T1 circular buffer is identical to an E1 circular buffer except that the multiframes are 24 bytes instead of 16 bytes.


Figure 22 - TDM Circular Buffer (Complete Buffer in T1 mode, FASTCAS)

Figure 22 TDM Circular Buffer (Complete Buffer in T1 mode, FASTCAS) shows a T1 circular buffer in FASTCAS mode. The T1 circular buffer is identical to an E1 circular buffer except for the fact that the multiframes are 24 bytes instead of 16 bytes.

### 4.3 TX_SAR Module

### 4.3.1 Overview

The purpose of the TX_SAR Module is to assemble TDM data bytes into ATM cells to be transmitted to the UTOPIA Module. The TX_SAR Module has no external interfaces, and does not control any of the MT90503's pins. However, it is the connecting block between the TDM module and the UTOPIA Module in the data transmission direction.

When the TX_SAR assembles ATM cells, it reads bytes from the circular buffers in data memory and uses these to assemble the ATM cells.

In order to transmit CBR ATM cells correctly, a timing algorithm is used to ensure that the ATM cells are assembled and transmitted at the defined data rate. This timing algorithm is implemented using the transmit event scheduler. The transmit event scheduler is a construct in control memory that identifies the time ATM cells need to be transmitted, and the circular buffer from which data for the ATM cell is to be retrieved. ATM cells are capable of carrying a variable number of virtual channels. Therefore, the rate at which cells are transmitted is variable.

### 4.3.1.1 Support and Trunking for Different Types of ATM Cells

The type of ATM cells the TX_SAR supports includes AAL1, CBR-AALO, and AAL5-VTOA. The transmit event scheduler is capable of supporting the different data rates which are inherent to the different types of ATM cells. Figure 23 - ATM Cell Formats shows the different formats for the ATM cells that the TX_SAR Module supports.
The TX_SAR is also capable of supporting hyperchannels (trunking) of up to 2048 channels per VC. In addition, it can also support AALO and AAL5 trunking, providing the number of channels is divisible by the payload bytes.


Figure 23 - ATM Cell Formats

### 4.3.2 TX_SAR Event Schedulers

### 4.3.2.1 Overview

The purpose of the transmit event scheduler is to ensure that the MT90503 assembles and transmits ATM cells at the appropriate time. This timing function is implemented to reduce data overruns and underruns.

Before arriving at the TX_SAR, TDM channel data is written into a variable-length ( 128 to 1024 bytes) transmit circular buffer, a construct in external data memory, by the TDM Module. Each TDM channel is assigned its own transmit circular buffer (section 4.2.6 TDM Circular Buffer Pointers). The transmit event schedulers determines when the information will be assembled into ATM cells.

The 15 transmit event schedulers are identical, and all have individual configuration registers. Each transmit event scheduler is divided into a programmable number of frames. When configured correctly, each transmit event scheduler frame is constrained to an average of $125 \mu \mathrm{~s}$ (i.e. the time required for one byte to be received/transmitted on each TDM channel).

### 4.3.2.2 The Transmit Event Scheduler Process

Please refer to Figure 24 - Transmit Event Scheduler Process in conjunction with this subsection.

The MT90503's 15 transmit event schedulers are all maintained in a designated memory block in the external control memory. Each transmit event scheduler can be programmed to support one of the four types of cells (AAL1 with or without pointer, CBR-AALO, or AAL5-VTOA). Each of the transmit event schedulers can be enabled or disabled via registers ( $0610 \mathrm{~h}-0616 \mathrm{~h}$ ) in order to use less bandwidth if a certain format of ATM cell is not required. Figure 24 - Transmit Event Scheduler Process provides an example of one of the 15 transmit event schedulers in the MT90503. All transmit event schedulers have exactly the same properties, and all perform the same functions. They can be configured individually to handle different VC configurations.

Each of the 15 transmit event schedulers is made up of a number of "frames", with each frame containing a number of events. An event, if executed, consists of the assembling and placing in a UTOPIA output FIFO of one ATM cell. The base address, length, and number of events can be programmed for each scheduler frame. Each transmit event scheduler has its own Scheduler Base Address which, in conjunction with an internal frame counter, will locate the events in a specific frame.

In order for an ATM cell to be assembled and transmitted at the appropriate time, the following transmit event scheduler process steps are required:

1 On the reception of a frame pulse, the scheduler scans through the events of the current frame.
2 For each valid event in the current frame, a cell is assembled and transmitted to the UTOPIA module.
3 On the reception of the next frame pulse ( $125 \mu \mathrm{~s}$ later) or upon the completion of the final valid event in the current frame, whichever is later, the scheduler increments current frame, to point to the next frame in the scheduler.

4 The scheduler reads each frame sequentially, returning to frame 0 upon completion of the final frame in the scheduler.

An analogy can be made of the above process to a continuous "spinning wheel", where there is a continual selection and reading of events for each transmit event scheduler frame. Table 19 - Scheduler Event Fields provides a description of the fields for one of the transmit event schedulers. Refer to Figure 24 - Transmit Event Scheduler Process to locate the fields.


Figure 24 - Transmit Event Scheduler Process

### 4.3.2.3 Transmit Event Scheduler Fields Description

Table 18 - Field Description for the Transmit Event Scheduler provides a description of the fields for the transmit event scheduler. Please refer to Figure 24 - Transmit Event Scheduler Process to locate the fields.

| Field | Name of Field | Bits Used | Description of Field |
| :---: | :---: | :---: | :---: |
| I | Scheduler Initialised Bit | +0/b15 | This bit is reset by software immediately before enabling the transmit event scheduler. Hardware will set this bit the first time it reads the scheduler information structure. When the I bit is read at ' 0 ' by hardware and the scheduler is used for T1/E1 support, the transmit event scheduler's Current Frame will be written between $0-23$ for T 1 and $0-15$ for E 1 in order for frame 0 to correspond to the first byte of a multiframe. |
| Turn Num | Free run counter of scheduler wraps | +0/b13:b11 | Turn Num is a counter used for the implementation of multiframing. The scheduler re-synchronises itself with the multiframe count when Turn Num is 0 . |
| Current Frame | Current Frame | +0/b10:b0 | Current Frame is used to record the frame position of the transmit event scheduler. <br> Range 0 to (Last Frame) |
| Num Env | Number of events per frame in the scheduler | +2/b7:b5 | The number of events per frame is $2^{(N u m ~ E n v)+1}$. The same number applies to all frames in the scheduler. Range:000 to 101 (2 to 63 events per frame). All others reserved. |
| Read 2m Events | Read 2m First Events | +2/b4:b0 | Indicates to the scheduler how many events in each frame must be read. The same number applies to all frames in the scheduler. |
| Scheduler Base Address | Pointer to the beginning of the transmit event scheduler | +4/b14:b0 | This field is appended with "00000" as the LSBs, to form a 20-bit address. <br> Note: A frame must never cross a boundary of its own size in memory. Therefore, if the transmit event scheduler has more than 8 events per frame ( 32 bytes per frame), then some LSBs of this field must be 0 . |
| Last Frame (M) | Last Frame (M) | +6/b10:b0 | Last Frame $=$ (number of frames in the scheduler) - 1. If Last Frame $=0$, the scheduler is one frame long, if Last Frame $=$ 1, the scheduler is two frames long, etc. |

Table 18 - Field Description for the Transmit Event Scheduler

The number of events per frame in the scheduler, Num Env, determines how much memory is allocated for each frame in the scheduler: 2 words per event. Within this memory, the scheduler will read only the first 2 m
events, as stored in the Read 2 m Events field. Of the events that are read, only those with the Scheduler Num not equal to '1111' (see Table 19) are executed. No action is taken on events that are read whose Scheduler Num is '1111'.

### 4.3.2.4 Scheduler Events Fields Description

Table 19 - Scheduler Event Fields provides a description of the fields for the transmit event scheduler entries. Figure 24 - Transmit Event Scheduler Process to locate the fields

| Field | Name of Field | Bits <br> Used | Description of Field |
| :--- | :--- | :---: | :--- |
| Scheduler Num | Transmit Event <br> Scheduler <br> Number | $+0 / b 5: b 2$ | This field contains the scheduler number. <br> 0000 to 1110 = valid scheduler numbers <br> $1111=$ invalid event. No action will be taken on this event <br> The transmit event schedulers are read sequentially. |
| S | Start Bit | $+0 / b 0$ | This field indicates the first event that will be carried out <br> when the VC is initialised. <br> Software must set the S bit of only one event once <br> programming of the scheduler is complete. <br> This bit is only relevant until the I bit is set in the TX Control <br> Structure. After the I bit is set, all events with valid scheduler <br> numbers will be executed as they are encountered. |
| TX Structure <br> Address Pointer | TX Structure <br> Address Pointer | $+2 / b 14: b 0$ | This field is the pointer to the TX_SAR Structure used to <br> assemble an ATM cell each time this event is read. This field <br> is appended with "00000" as the LSBs to form a 20-bit <br> address. |

Table 19 - Scheduler Event Fields

### 4.3.2.5 Bandwidth Limitations for Transmit Scheduler Events

| Transmission Speed | Maximum Number of Events per Frame |
| :---: | :---: |
| 25 Mbps | 7 |
| 155 Mbps | 45 |
| 622 Mbps | 183 |

Table 20 - Maximum number of Events per Frame for Common Transmission Speeds

The maximum number of events per frame corresponds to the maximum number of events that can occur in one frame of all the schedulers simultaneously.

If the frames are not synchronised (as in Figure 25 - Unsyncrhonised Schedulers), the maximum number of events per frame is the sum of the maximum events in any frame of each scheduler (in this case, 32 events + 21 events $=53$ events is the maximum number of events per frame). If the schedulers are synchronised (as in Figure 26 - Synchronised Schedulers), the maximum number of events per frame is the worst in any particular frame (in this case, $23+17=40$ events is greater than $6+24=30$ events, so 40 events is the maximum number). If the schedulers are partially synchronised (as in Figure 27 - Partially Synchronised Schedulers), the maximum number of events per frame is the worst in any frame that can possibly align (in this case, $20+5=$ 25 events is less than $27+12=39$ events which is greater than $20+16=36$ events and $27+7=34$ events so 39 events is the maximum number).


Figure 25 - Unsyncrhonised Schedulers


Figure 26 - Synchronised Schedulers


Figure 27 - Partially Synchronised Schedulers

### 4.3.3 Out of Bandwidth Error

The TX_SAR out-of-bandwidth error indicates that a particular region of the scheduler, or the scheduler as a whole, is overloaded. This error occurs when the TX_SAR lagging by $F$ frames or more, where $F$ is an integer written to register 0608h. The TX_SAR has $125 \mu \mathrm{~s}$ to execute all the events in the frame it is servicing at the current time. If all the events in the frame are not executed within $125 \mu \mathrm{~s}$, the TX_SAR continues until the end of the current frame before moving on to events in the next frame. This method reduces the amount of time it has to service the following frame. If the first frame took $150 \mu \mathrm{~s}$ to service, then there would only be $100 \mu \mathrm{~s}$ left to service the following frame. Therefore, if the second frame is also very full, the TX_SAR would not have the time to finish servicing it and will be late again for the following frame.

If frames take an excessively long time to service, the MT90503 may become several frames late. If the MT90503 detects that it is F or more frames late, an out-of-bandwidth error will be declared and the MT90503 will skip the next F frames of events. This action will cause latency in the transmission delay and corrupt data on the VCs until the channel reconverges. The channel reconverges once every 375 frames ( 47 ms ) in "AAL1 with pointer" schedulers, and less often if multiframing is used. The cells that were supposed to be assembled in the F missing frames are skipped, and new cells are assembled with the old data. The out-of-bandwidth error
arises from an error with the mapping of the transmit event scheduler. Despite this error, the MT90503 will continue with its operations, however, some cells may be lost.

### 4.3.3.1 Percent of Bandwidth Register

The TX_SAR incorporates a Percent of Bandwidth register ( 0510 h ), which indicates the maximum number of mclk cycles utilised by the TX_SAR process per frame. The Percent of Bandwidth register counts the number of molk cycles from the time that a frame is received, to the time that all the entries for that frame in the schedulers have been completely handled. This value will be compared to the current maximum value obtained by the TX_SAR process. If the most current value is higher, it is retained and written into the Percent of Bandwidth register. The value in the Percent of Bandwidth register can be cleared by the software. If the TX_SAR takes more than $125 \mu \mathrm{~s}$ to assemble the cells of a particular frame, the Percent of Bandwidth register's value will be greater than the number of molk cycles in the $125 \mu \mathrm{~s}$ time frame.

### 4.3.3.2 Distribution of Events by Software

The events in the transmit event scheduler need to be distributed as evenly as possible. For example, an "AAL1 with pointer" 3 -channel VC has three events per 46.875 frames. These events must be spaced out with a distance of 15 or 16 frames between events to prevent irregular data distribution resulting in transmission latency and/or data integrity problems. Consequently, a scheduler 8 * $46.875=375$ frames in length containing 24 events is required to map the three events per 46.875 frames evenly. The mapping of this information is accomplished by external software.

One of the key features of the transmit event schedulers is their programmable length. Since different cells require different sizes of schedulers, the schedulers are capable of handling any scheduler length from 1 to 2048 frames.

Refer to Table 21 for examples of transmit event scheduler sizes.

| ATM Cell Type | Number of Frames | Transmit Event Scheduler size in KB <br> (e.g. 32 events per frame) |
| :--- | :---: | :---: |
| CBR-AALO | 48 | $6^{*}$ |
| AAL1 with Pointer | 375 | 47 |
| AAL1 without Pointer | 47 | 5.875 |
| AAL5-VTOA/CBR-AALO | 240 | 30 |
| E1 with CAS | 750 | 94 |
| T1 with CAS | 1125 | 141 |
| ${ }^{*}$ Note: 6 KB $=32$ events per frame * 48 frames per scheduler * 4 bytes per event |  |  |

Table 21 - Examples of typical Transmit Event Scheduler Sizes

### 4.3.4 Mapping of the Transmit Event Scheduler

With the "AAL1 with pointer", the format for the mapping of the transmit event scheduler is asymmetric. The "AAL1 with pointer" format expects a cell every 46.875 frames per channel, a transmit event scheduler with 47 frames that skips the last frame one turn out of eight is required. The MT90503 overcomes this difficulty by creating an extended transmit event scheduler of 375 frames, which is 8 * 46.875 . Therefore, the irregularity of the transmit event scheduler is corrected and events can be mapped appropriately.

Due to this multiplication, the number of events mapped in the transmit event scheduler for any VC is always a multiple of eight. This means the first event in the transmit event scheduler is always a p-byte event and contains a zero value pointer indicating the start of a structure.

### 4.3.5 TX_SAR Control Structures

TX_SAR control structures are constructs in control memory which contain ATM cell information. When a frame event is read from one of the 15 transmit event schedulers, an ATM cell is assembled. The scheduler event contains the base address that points to the control structure that is used for assembling the ATM cell. The control structure contains all of the fields that are required for assembling the ATM cell. The destination field in the control structure is responsible for telling the UTOPIA Module the destination VC of the assembled ATM cell. For detailed information regarding the destination field, refer to Table 22, 'Description of the Fields for the TX_SAR Control Structure", on page 60.

Figure 28 - TX_SAR Event Scheduler Pointer Flow and Control Structure shows a functional block diagram of an example of transmit event scheduler interconnections and pointer flow to the TX_SAR control structure.


Note: All structures are in the Control memory unless otherwise specified
Figure 28 - TX_SAR Event Scheduler Pointer Flow and Control Structure

### 4.3.5.1 TX_SAR Control Structure Fields

Figure 29 - TX_SAR Control Structure shows the TX_SAR control structure format and data fields held in control memory. Table 22, Description of the Fields for the TX_SAR Control Structure describes the data fields.

The header for the ATM cell to be constructed is derived from words +8 and +A in the TX_SAR control structure. The final 8 bits of the cell header, the Header Error Check (HEC), are calculated and inserted by the UTOPIA module. The channel entries, beginning at +16 h, are 14 -bit pointers to the circular buffers in external data memory, one for each channel in the VC.


Figure 29-TX_SAR Control Structure

| Field | Name of Field | Byte Address Offset /Bits Used | Description of Field |
| :---: | :---: | :---: | :---: |
| First Entry | First Entry | +0/b15:b12 | First entry gives the position within the structure of the pointer to the first channel. The field is a word pointer and is constant with a value of $0 \times B$. To allow for future structure updates, its value is programmable. <br> Channel 0's address is located at base_add + (2 * First Entry). |
| A | Structure <br> Active | +0/b11 | This bit indicates whether or not the VC is active. The control structure becomes active when $A$ is set. The $A$ bit is set by software. If the A bit is reset, the control structure will always be ignored. If the A bit is set, hardware may ignore it for other reasons, such as the control structure not being initialised. |
| SRTS | Synchronous <br> Residual Time Stamp | +0/b8:b7 | The SRTS bits indicate the generation of SRTS within the VC. SRTS is either absent, enabled, or enabled and master. Many VCs can be programmed to transmit SRTS, but only one can request a new SRTS value. If many VCs transport SRTS, they must all be of the same size, i.e., the same number of channels, to ensure the validity of the values. If SRTS is generated on the VC, it is packaged within the AAL1 byte of cells $1,3,5$ and 7 of an 8 -cell cycle. <br> The SRTS field should never be enabled on CBR-AAL0 or AAL5-VTOA VCs. If SRTS is enabled on multiple VCs, the events must be mapped in such a way that an event that generated cell \# 0 for the master VC occurs before the event that generates cell \# 0 for all the other VCs. <br> The SRTS field bits are encoded as follows: <br> $00=$ No SRTS <br> 01 = Reserved <br> 10 = Send SRTS <br> $11=$ Send SRTS, and request a new SRTS value on Seq $=7$ |
| AAL | Adaptation Layer | +0/b6:b5 | The AAL bits indicate the ATM format to be used to assemble the cells in this structure. <br> The AAL field bits are decoded as follows: <br> $00=$ CBR AALO. Consists of CBR data <br> 01 = AAL5-VTOA. Includes a CRC and a payload size indicator at the end of the cell <br> 10 = AAL1 without pointer. Includes a Sequence number packaged with the cell, as well as the possibility of transmitting SRTS on the VC <br> 11 = AAL1 with pointer. Includes a Sequence number packaged with the cell, and the possibility of transmitting SRTS on the VC |

Table 22 - Description of the Fields for the TX_SAR Control Structure

| Field | Name of Field | Byte Address Offset /Bits Used | Description of Field |
| :---: | :---: | :---: | :---: |
| CE | Circuit Emulation | +0/b4:b2 | The CE field bits indicate the presence of a multiframing standard used within the VC. <br> The multiframing standard changes the way the information is read from the TDM circular buffers, as well as the standard that is used to generate the p-byte. In addition, setting the CE bits to a CAS setting will indicate to the TX_SAR that CAS signaling bits must be inserted at the end of the AAL1 control structure. <br> The field bits are decoded as follows: <br> $000=$ No CAS, no multiframe <br> 001 = Reserved <br> 01x = Reserved <br> $100=\mathrm{T} 1$ without CAS (T1 type circular buffer, but no CAS will be sent in the ATM cells) <br> 101 = T1 with CAS (T1 type circular buffer, with CAS sent in the ATM cells <br> $110=\mathrm{E} 1$ without CAS (E1 type circular buffer, but no CAS will be sent in ATM cells) <br> 111 = E1 with CAS (E1 type circular buffer, with CAS sent in the ATM cells |
| BS | TX/RX <br> Circular <br> Buffer Size | +0/b1:b0 | These bits encode the size of the TDM circular buffer that is to be read from. In T1 and E1 modes, a portion of space in the circular buffer is required to store the CAS values; in the T1 mode $25 \%$ is required, and in the E1 mode $50 \%$ is required. Half of the TDM circular buffer is used for TX data and half for RX data. <br> The field bits are decoded as follows: <br> $00=128$ words <br> $01=256$ words <br> $10=512$ words <br> $11=1024$ words |
| Last Entry | Last Entry | +2/b15:b4 | Last Entry is the word offset from the 8 kB boundary containing the structure to the last circular buffer pointer in the structure. |
| TXD | TX_SAR <br> Destination Field | +2/b3:b0 | The TXD field is used to tell the UTOPIA Module the destination of the assembled ATM cell. <br> The TX_SAR Destination field is decoded as follows: <br> $0000=$ Discard ATM cell <br> $0 \times X 1=$ Send to TXA port <br> $0 \times 1 X=$ Send to TXB port <br> 01XX = Send to TXC port <br> Others = Reserved <br> Broadcasting to multiple ports is allowed. |
| Payload Size | Payload Size | +4/b15:b10 | For fully-filled cells, regardless of type, the payload size field is 30h. For partially-filled cells, the payload size indicates the number of TDM bytes to be placed in each ATM cell. The field range is from 4 h to 2 Fh . <br> Note: 2Fh is an illegal value for partially-filled AAL1cells For partially-filled AAL5 VTOA, this field must be set to 8 h , $10 h, 18 h, 20 h$, or 28 h. |

Table 22 - Description of the Fields for the TX_SAR Control Structure (continued)

| Field | Name of Field | Byte Address Offset/Bits Used | Description of Field |
| :---: | :---: | :---: | :---: |
| Offset | Offset | +4/b9:b3 | Offset is used when the VC first starts up, and whenever an event with the $S$ bit is set in the transmit event scheduler entry. This shows the delta that must exist between the TX_SAR read pointer and the TDM write pointer within the circular buffer. The value for this offset will change depending on the number of channels in the VC , and on the multiframing standard used. <br> Offset's value is programmed with the maximum number of bytes of a given channel in an ATM cell, plus three. If an error is produced when programming Offset, a global tx_slip will be flagged in the register 0502h, indicating an erroneous configuration and the possibility of corrupted data. The Offset between the tx_wrt_pnt (TDM) and the tx_sar_read_pnt that must be present prior to assembling any ATM cell whose event in the transmit event scheduler has the S bit set. This offset is coded as an integer from normal VCs. It is coded as an multi-frame [1:0] and frame [4:0] number for E1 and T1 VCs. |
| GFC, VPI, VCI, OAM, PTI, CLP | Header information | $\begin{aligned} & +8,+\mathrm{A} / \\ & \mathrm{b} 15: \mathrm{b0} \end{aligned}$ | Header information for the cells to be assembled using this TX_SAR control structure. |
| P-Byte Counter | P-Byte Counter | +C/b14:b0 | The $p$-byte counter field is used for the generation of the p-byte within the cell, or within the multiframe structure. <br> The p-byte counter is decremented each time a byte of data is sent, including a CAS byte, but not including an information byte such as the AAL1 byte or the pointer-byte. Whenever the counter reaches 0 and must decrement, its value is reset to $p$ byte Max field which must be set to 0 by the software. Whenever a p-byte needs to be generated, the seven LSBs of the $P$-byte Counter are the value of the p-byte, with parity added as the MSB. The decrementing continues until the value of the P-byte Counter reaches 0 and wraps around again, ending the multiframe and beginning a new one. |
| Current Entry | Current Entry | +E/b15:b4 | The Current Entry field tells the SAR what the first channel is to be assembled in this cell. The valid values for this field are contained between the First Entry and the Last Entry fields inclusively. The Current Entry field also points to words, and is initialised by software to the word pointed to by the First Entry field. the The Current Entry field is similar the Last Entry field, such that it is defined as the offset between the Current Entry field and the 8 kB boundary in which the structure is contained. |
| 1 | Structure Initialised Bit | +E/b3 | The I bit is set by the hardware when a scheduler entry flagged with the Start bit asserted high is encountered, i.e., as soon as the first cell is sent on this VC. <br> This bit is cleared by the software upon initialisation. |
| Seq | AAL1 <br> Sequence <br> Number | +E/b3:b0 | These bits are reset by software. |

Table 22 - Description of the Fields for the TX_SAR Control Structure (continued)

| Field | Name of Field | Byte Address Offset /Bits Used | Description of Field |
| :---: | :---: | :---: | :---: |
| $\begin{aligned} & \text { tx_sar_read_ } \\ & \text { pnt } \end{aligned}$ | Tx_sar_read _pnt | +10/b15:b6 | This value is a pointer to the next byte to be read in the TX/RX Circular Buffers. |
| Transmitted Cell Counter | Free running transmitted cell counter | $\begin{aligned} & +12,+14 / \\ & \text { b15:b0 } \end{aligned}$ | The Transmitted Cell Counter increments each time a cell is transmitted on the VC. <br> The time needed for the counter to wrap-around decreases proportionately with the number of channels in the VC. This field should be reset by software and is used for monitoring. |
| V | TDM Channel Valid | $\begin{aligned} & \text { +16 to end/ } \\ & \text { b15 } \end{aligned}$ | When this bit is high the channel is active. When this bit is ' 0 ', the programmable null byte, found in register 0420h, will be transmitted. |
| Channel N's <br> TXIRX <br> Circular <br> Buffer <br> Pointer | Channel N's <br> TXIRX <br> Circular <br> Buffer <br> Pointer | $\begin{aligned} & +16 \text { to end/ } \\ & \text { b13:b0 } \end{aligned}$ | This field is a pointer to the TX/RX Circular Buffer associated with this channel. "0000 0000" will be appended to this field as the LSBs to form a 22-bit address in data memory. |

Table 22 - Description of the Fields for the TX_SAR Control Structure (continued)

### 4.3.6 Miscellaneous TX_SAR Features

### 4.3.6.1 T1 with CAS and E1 with CAS Cell Format Mapping

For the transmission of T1 with CAS and E1 with CAS, the number of frames required is 9000 and 6000 respectively. But, because the number of frames mapped in a full size transmit event scheduler is always a multiple of eight, the MT90503 handles large size schedulers by dividing the size by eight, and mapping eight times less events ( 1125 for T1 with CAS and 750 for E1 with CAS). This results in a latency of one frame.

### 4.3.6.2 Support of Partially-Filled Cells

The MT90503 is capable of supporting partially-filled cells, as long as the number of channels in the VC is smaller or equal to the cell fill. A single transmit event scheduler can be created to accommodate several sizes of AAL1, CBR-AAL0 or AAL5-VTOA ATM cells. AAL1 with pointer ATM cells can be used by this format, since the number of data bytes per ATM cell is always constant. For example, the 240 -entry transmit event scheduler is capable of accommodating partially-filled cells of $4-, 5-, 6-, 8-, 10-, 12-, 15-, 16-, 20-, 24-, 30$-, or 40 -bytes per cell.

### 4.3.6.3 TX_SAR FIFO

The TX_SAR transmits cells contained in its data cell FIFO when there are no events to be processed. The TX SAR's data FIFO can be used to store AALO cells as well as OAM cells, therefore, CPU-based OAM cell generation is supported.

### 4.4 RX_SAR Module

The RX_SAR module performs processing on ATM cells received from the UTOPIA module. Cells placed in the RX_SAR input FIFO by the UTOPIA module are read, processed, and then written into the appropriate multicell circular buffer in external control memory. The processing involves identifying the VC corresponding to the cell, examining the cell for errors, determining where to place the data, and monitoring the status of circular buffers. The RX_SAR module also directs data cells to the data cell FIFO from which the CPU can read them.

The RX_SAR module does not connect to any external pins, interfacing instead with the UTOPIA module, CPU interface, and external memory controller. Global pointers are shared between the RX_SAR and TDM modules.

ATM cells that are received by the MT90503 are processed by the UTOPIA module and can be directed to any of three TX output FIFOs or to the 32-cell RX_SAR input FIFO. Those cells that are forwarded to the RX_SAR are directed to the SAR portion (cells to be formatted into TDM streams), to the data cell portion (to be examined by the CPU), or to both.

For cells directed to the SAR portion, the RX_SAR uses control information in the RX_SAR Control Structures (Section 4.4.2) to extract the payload data from the received cell and store it into TDM channel RX Circular Buffers located in external data memory.

### 4.4.1 Treatment of Data Cells

Data cells, such as those containing OAM information, are placed in a programmable length FIFO in external control memory. The length of the FIFO is stored in register 070Eh. The CPU can read a data cell at any time, after obtaining the address of the FIFO (register 070 Ch ) and the read pointer (register 0708h).

The CPU can be alerted to the presence of data cells via an interrupt that triggers if either of two events occur: the interrupt can be generated when the FIFO becomes more than half full or the interrupt can be generated if a data cell has been present in the FIFO for longer than a programmable period of time (registers 0720h, 0722h). This interrupt can be enabled through register 0220h.

When ready to process the information, the CPU obtains a read pointer to the information from register 0708h and reads the information through 26 word accesses.

Cells with the OAM bit set in the PTI portion of the header can be directed to the data cell FIFO on a per VC basis. The same is true for non-OAM cells. In addition, unknown non-OAM cells, and/or unknown OAM cells can also be sent to the data cell FIFO (all unknown non-OAM cells are directed to the same location(s)).

### 4.4.2 Control Structure

For each VC directed to the SAR portion of the RX_SAR, an RX_SAR control structure exists in external control memory. The structure, similar to that of the TX_SAR, contains information on how to process a cell including:

- what type of traffic is being carried (AAL1, CBR-AAL0, AAL5-VTOA)
- the size of the circular buffers for the data
- how to act in the case of an overrun or underrun
- the multiframing standard in place
- expected presence of CAS
- enabled errors, the size of the payloads
- information necessary to detect and correct for errors
- a pointer to each circular buffer, one for each channel in the VC


Figure 30 - RX_SAR Control Structure

| Field | Name of Field | Byte Address Offset/Bits Used | Description of Field |
| :---: | :---: | :---: | :---: |
| First Entry | First Entry | +0/b15:b12 | First entry field gives the position within the structure of the pointer to the first channel. The field is a word pointer and is constant with a value of $0 \times B$. To allow for future structure updates, its value is programmable. <br> Channel 0's address is located at base_add + (2 * First Entry). |
| AAL | Adaptation Layer | +0/b8:b7 | The AAL bits indicates the ATM format used to assemble the cells in this structure. The AAL bits are decoded as follows: $\begin{aligned} & 00=\mathrm{CBR} \text { AAL0 } \\ & 01=\mathrm{AAL5}-\mathrm{VTOA} \\ & 10=\mathrm{AAL} 1 \text { without pointer } \\ & 11=\text { AAL1 with pointer } \end{aligned}$ |
| BS | TX/RX <br> Circular Buffer Size | +0/b6:b5 | These bits encode the size of the entire TX/RX circular buffer that is to be written into. In T1 and E1 modes, a certain portion of the space in the buffer is required to store the CAS values; in the T1 mode $25 \%$ is required and in the E1 mode $50 \%$ is required. <br> The eight MSBs of each word in the TX/RX Circular buffer are used for TX data and the remaining 8 bits for RX data. These bits are decoded as follows: $\begin{aligned} & 00=128 \text { words } \\ & 01=256 \text { words } \\ & 10=512 \text { words } \\ & 11=1024 \text { words } \end{aligned}$ |
| IM | Initialisation Method | +0/b4 | Initialisation method for rx_sar_write_pnt '0' = initialise rx_sar_write_pnt to nearest boundary: either tdm_read_pnt + (Max used bytes in TX/RX Circular buffer) 1 OR tdm_read_pnt + 1 <br> '1' = initialise rx_sar_write_pnt to (Max used bytes in TX/RX Circular buffer)/2 <br> This initialisation method will be used when the first cell is received and each time a slip (an overrun or underrun) occurs. |
| CE | Circuit Emulation | +0/b3:b1 | The circuit emulation bits indicate the multiframing standard used within the VC. <br> These bits are decoded as follows: <br> $000=$ No CAS, no multiframe <br> 001 = Reserved <br> 01x = Reserved <br> $100=$ T1 without CAS <br> 101 = T1 with CAS <br> $110=E 1$ without CAS <br> 111 = E1 with CAS |
| FC | FASTCAS Enable | +0/b0 | FASTCAS is enabled when FC is asserted high. If FASTCAS is used, then the receive pointer used for regular TDM bytes and for CAS bytes is not the same. CAS is written one multiframe before TDM bytes in FASTCAS. When FC is deasserted, regular multiframing is employed: CAS is written in the same multiframe as TDM bytes. |

Table 23 - Description of the Fields for the RX_SAR Structure

| Field | Name of Field | Byte Address Offset/Bits Used | Description of Field |
| :---: | :---: | :---: | :---: |
| Last Entry | Last Entry | +2/b15:b4 | Last Entry is the word offset from the 8 kB boundary containing the structure to the last circular buffer pointer in the structure. |
| SE | Slip Error Report Enable | +2/b3 | When set, overrun and underrun slips will generate an RX_SAR Error Report Structure (Figure 32) in the error FIFO in control memory. |
| LE | Cell Loss and Insertion Error Report Enable | +2/b2 | When set, single cell losses, multiple cell losses and cell misinsertions will generate an RX_SAR Error Report Structure (Figure 32) in the error FIFO in control memory. |
| AE | AAL1 Byte Report Enable | +2/b1 | When set, AAL1 CRC-3 and CRC-3/Seq Num parity errors will generate an RX_SAR Error Report Structure (Figure 32) in the error FIFO in control memory. |
| PE | P-Byte Error Report Enable | +2/b0 | When set, p-byte parity errors, p-byte absent and p-byte framing errors will generate an RX_SAR Error Report Structure (Figure 32) in the error FIFO in external control memory. |
| Max Used Bytes in RX Circular Buffer | Max Used Bytes in RX Circular Buffer | +4/b15:b6 | This field indicates the range of valid positions in bytes of the rx_sar_write_pnt relative to the tdm_rx_read_pnt before a cell is received. In no case can this value be 0. |
| Payload Size | Payload Size | +4/b5:b0 | For fully-filled cells, regardless of type, the payload size field is 30 h . For partially-filled cells, the payload size indicates the number of TDM bytes in each ATM cell. The field range is from 4 h to 2 Fh . <br> Note: 2Fh is an illegal value for partially-filled AAL1cells For partially-filled AAL5 VTOA, this field must be set to 8 h , $10 \mathrm{~h}, 18 \mathrm{~h}, 20 \mathrm{~h}$, or 28 h . |
| PBI | P-Byte Initialization | +6/b15 | Initialized to '0' by software. |
| P-Byte Counter | P-Byte Counter | +6/b14:b0 | This field is a counter used to detect p-byte framing errors. When a p-byte is detected, this counter is loaded with the pbyte's value. Each time a TDM byte is received, the counter is decremented. When the counter reaches the Current Entry it should be reset to the First Entry and the rx_sar_write_pnt should point to the beginning of the multiframe if applicable. When it decrements below 0 , the counter resets to its maximum. Any time a p-byte is received, its value should match the value in this field, or a p-byte framing error will be detected. |
| Current Entry | Current Entry | +8/b11:b0 | Current Entry indicates which TDM channel the RX_SAR is currently writing to the circular buffers. The current entry is defined as being the offset from the 8 KB boundary containing the structure and the "TX/RX Circular Buffer Pointer" being read from the structure. Current Entry is initialised to the value of First Entry, increments up to Last Entry, and then wraps around to First Entry. This field should be initialised by software to OBh. Note: This field should not be written to while the VC is active. |

Table 23 - Description of the Fields for the RX_SAR Structure (continued)

| Field | Name of Field | Byte Address Offset/Bits Used | Description of Field |
| :---: | :---: | :---: | :---: |
| $\begin{aligned} & \hline \text { rx_sar_wr } \\ & \text { ite_pnt } \end{aligned}$ | RX SAR write pointer | +A/b15:b4 | This is a pointer to the location where the next byte will be written in each TX/RX Circular Buffer. It is common to all TX/ RX Circular buffers controlled by the RX_SAR Control Structure. <br> In the E1/T1 mode, this field is divided in multiframe [6:0] and frame[4:0]. Only the lower part of the field is used to point to the TX/RX Circular Buffer. |
| I | Structure Initialised Bit | +A/b3 | This bit must be reset by software before enabling the VC in the LUT. The bit is set by hardware after receiving the first cell. |
| Last Seq | Last Seq | +A/b2:b0 | Last received AAL1 sequence number. |
| $\begin{aligned} & 2^{\text {nd }} \text { Last } \\ & \text { Seq } \end{aligned}$ | $2^{\text {nd }}$ Last Seq | +C/b2:b0 | Second-last received AAL1 sequence number. |
| Slip Free rx_sar_wr ite_pnt | Slip-free RX_SAR write pointer | +10/b15:b4 | The Slip-free rx_sar_write_pnt is the same as the rx_sar_write_pnt except that slips (overruns and underruns) do not affect its value. |
| Received Cell Counter | Received Cell Counter | $\begin{aligned} & +12 / \mathrm{b} 15: \mathrm{b0} \\ & +14 / \mathrm{b} 15: \mathrm{b0} \end{aligned}$ | A 32-bit free running cell counter used for monitoring activity on a VC and for statistical purposes. <br> This field can be initialised to '0' by software. |
| V | TDM Channel Valid | +16 to end/b15 | When this bit is high the channel is active. When this bit is reset, all received bytes on the channel are discarded. |
| Channel <br> N's TX/ <br> RX <br> Circular <br> Buffer <br> Pointer | Channel N's TXIRX Circular Buffer Pointer | $\begin{aligned} & +16 \text { to end/ } \\ & \text { b13:b0 } \end{aligned}$ | This field is a pointer to the TX/RX Circular Buffer associated with this channel. "0000 0000" will be appended to this field as the LSBs to form a 22-bit address in external data memory. |

Table 23 - Description of the Fields for the RX_SAR Structure (continued)

Memory containing the control structures is divided into 8 KB blocks. Zero or more control structures can exist in a block. Control structures must be fully contained in a single block. The pointers Current Entry and Last Entry are relative to the 8 KB block boundary in which their structure resides.

The Buffer Size (BS) field indicates the size of the Receive Circular Buffers. Though the Buffer Size is indicated in words, the received data occupies only the lower byte of each word; the data to be transmitted occupies the upper bytes. Selection of size: $128,256,512$, or 1024 words, depends on the amount of available memory and on the CDV for the VC. The receive half of the buffer must be capable of holding twice the maximum CDV (peak CDV) plus the packetisation size of the cells in the VC plus two additional bytes. Additional space must be added if E1 or T1 formats are employed (the buffer must be twice as big for E1 and a $4 / 3$ of the size for T1). To convert the maximum CDV from ms into bytes, a data rate of 8000 bytes/s must be applied. The packetisation size is defined as the maximum number of bytes a channel can contain in a single cell of the VC.

To find the maximum CDV supported by the buffer, the value of the "Max Used Bytes In Circular Buffer" field must be divided by two and multiplied by $125 \mu \mathrm{~s} /$ byte. Because a larger buffer will cause more delay through the RX_SAR, the choice of the value of "Max Used Bytes In Circular Buffer" must be made as a compromise between the CDV supported and the delay inserted by the RX_SAR. In the T1/E1 multiframe mode, the value of "Max Used Bytes In Circular Buffer" must be an integer number of frames and is counted in a multiframe/ frame fashion. The five MSBs select a number of multiframes and the five LSBs a number of frames. In strict muliframing, the five LSBs are always set to '00000'; they can be any value in FASTCAS.

Setting the multiframing standard in the circuit emulation (CE) field will change the way the information is read from the TDM circular buffers, as well as the standard that is used to interpret the p-byte. In addition, setting the CE bits to a CAS setting will indicate to the RX_ SAR that CAS signalling bits should be expected at the end of the AAL1 structure.

The payload size represents the number of TDM bytes in the cell and does not include pointer or AAL1-bytes.

| Format | TDM Bytes | Bytes Transmitted | Payload Size |
| :--- | :---: | :---: | :---: |
| AAL1 with pointer, non-p-type cell | 16 | 17 | 16 |
| AAL1 with pointer, p-type cell | 16 | 18 | 16 |
| CBR-AAL0 | 16 | 16 | 16 |

Table 24 - Payload sizes for various cell formats

### 4.4.3 Errors

| Category | Error | Coverage | Example |
| :---: | :---: | :---: | :---: |
| AAL1 Byte | CRC error | AAL1 cells |  |
|  | CRC/Sequence Number Parity error | AAL1 cells | CSI $=1, C R C=010$, Seq. num $=110$, parity bit $=1$. Even parity not observed: a parity error has occurred. |
| Cell loss/misinsertion | Single cell loss | AAL1 cells | Sequence numbers 1, 2, 4 received. Cell 3 has been lost |
|  | cell misinsertion | AAL1 cells | Sequence numbers 1, 2, 4, 3 received. A cell misinsertion has occurred |
|  | Multiple cell loss | AAL1 cells | Sequence numbers 1, 2, 5 received. Cells 3,4 have been lost. |
| P-byte errors | P-byte parity error | AAL1 cells | P-byte $=0100110$, parity bit $=0$. Even parity not observed: a p-byte parity error has occurred |
|  | P-byte out-of-range | AAL1 cells | P-byte received $=0 \times 52$, Structure length for that $\mathrm{VC}=$ $0 \times 18$. P-byte is too large |
|  | P-byte framing error | AAL1 cells | P-byte is $0 \times 15$ when $0 \times 17$ was expected. |
|  | P-byte absent | AAL1 cells |  |
| Slip errors | Overrun | all cells | See Figure 31. |
|  | Underrun | all cells | See Figure 31. |

Table 25-RX_SAR errors

Within the four categories of errors, there are eleven possible errors. All but slip errors (overruns and underruns) pertain only to AAL1 cells.

For the following errors, the only action taken is the generation of an error report structure (Section 4.4.4):

- CRC errors
- CRC/sequence number parity errors
- multiple cell loss
- p-byte parity
- out-of-range
- framing
- absent errors

It is the responsibility of the CPU and associated software to act upon notification that an error has occurred. The exceptions to this policy are

- single cell loss (a dummy cell is inserted and treated before accepting the received cell; the received cell counter is incremented by only one)
- cell misinsertion (the misinserted cell is discarded and the received cell counter is not incremented)
- overruns (data is lost and read/write pointers are adjusted)
- underruns (bytes are inserted according to information set in registers 0420h and read/write pointers are adjusted).


Note: For a description of Max Used Bytes in Circular Buffer see Table 23 (+4/b15:b6).

Figure 31-Overrun and Underrun Examples

### 4.4.4 Error Report Structure

Four categories of errors can be enabled in the RX_SAR structure:

- slip errors
- cell loss \& misinsertion errors
- AAL1 byte errors
- p-byte errors

For each cell containing an error for which the error category is enabled, an 8-byte error report structure is generated and stored in a FIFO in external control memory. From here, the CPU can read the FIFO and treat the errors.

As per the Data Cell FIFO, the CPU can read the error FIFO at any time after obtaining the address to the FIFO (registers 0714h, 0716h) and the read pointer to the RX_SAR Error Report Structure (register 0710h). Again, an interrupt can be generated that will trigger if either of two events occurs: when the FIFO becomes more than half full or if an error report structure has been present in the FIFO for longer than a programmable period of time (register 0724h, 0726h). This interrupt can be enabled in register 0220h.

Errors of more than one type on the same cell will result in one error report structure being created indicating all types of errors on that cell. The exception to this is cell misinsertion. Since the cell is discarded once the misinsertion has been detected, no further errors can be found on the misinserted cell.


Figure 32-RX_SAR Error Report Structure

| Field | Name of Field | Byte Address Offset/bits Used | Description of Field |
| :---: | :---: | :---: | :---: |
| RX SAR Structure Base Address [19:5] | RX_SAR <br> Structure Base <br> Address [19:5] | +0/b14:b0 | Base address of the RX_SAR Control Structure that caused the Error Report Structure to be generated. Appended with " 00000 " as the LSBs to form a 20-bit address. |
| Seq | Sequence Number | +2/b2:b0 | Sequence number of the cell that caused the Error Report Structure to be generated. |
| Last Seq | Last Sequence Number | +2/b5:b3 | Sequence number of the cell which preceded the one that caused the Error Report Structure to be generated. |
| Received Cell Counter | Received Cell Counter | $\begin{aligned} & +4 / \mathrm{b} 15: \mathrm{b0} \\ & +6 / \mathrm{b} 15: \mathrm{b0} \end{aligned}$ | Cell number of the cell that caused the Error Report Structure to be generated. |
| PA | P-Byte Absent Error | +0/b15 | A p-byte was expected but was not detected in the cell that caused the error. |
| PF | P-Byte Framing Error | +2/b15 | The p-byte that was detected did not match the pbyte that was expected. |
| PR | P-Byte Range Error | +2/b14 | The p-byte that was detected was out-of-range; the 7-bit p-byte must not be 94-126 and must be less than the sequence length of the corresponding VC. 127 signifies a dummy offset value. |
| PP | P-Byte Parity Error | +2/b13 | The p-byte detected did not match its parity bit. |
| AC | AAL1 Byte CRC-3 Error | +2/b12 | There was an error detected in the CRC-3. |
| AP | AAL1 Byte Parity Error | +2/b11 | The AAL1 byte detected did not match its parity bit. |
| OR | Overrun Slip | +2/b10 | The circular buffer associated with a channel receiving data from the cell was overrun. |
| UR | Underrun Slip | +2/b9 | The circular buffer associated with a channel receiving data from the cell was underrun. |
| M | Cell Misinsertion | +2/b8 | The sequence number of the cell indicated that a cell was misinserted. This error is always preceded by a single cell loss. |
| ML | Multiple Cell Loss | +2/b7 | The sequence number of the cell indicated that multiple cells were lost. |
| SL | Single Cell Loss | +2/b6 | The sequence number of the cell indicated that a single cell was lost. |

Table 26 - Description of the Fields for the RX_SAR Error Report Structure

### 4.5 UTOPIA Module

### 4.5.1 Overview

The purpose of the UTOPIA module is to provide an external interface with the ATM domain. The MT90503 complies with The ATM Forum's specifications: af-phy-0017.000 and af-phy-0039.000. The MT90503 uses octet-level handshaking on its UTOPIA interface.

The UTOPIA module is responsible for accepting cells from four input interfaces, examining the cells and, based on the source and information in the header, sending the cell to one or more of the four output interfaces. The UTOPIA module also calculates and appends the HEC to outgoing ATM cells.


Figure 33 - UTOPIA Module
The UTOPIA interface consists of three ports, labelled A, B, and C. Port A is a Level-2 ATM (single PHY), single PHY or multi-PHY port and can operate at 50 MHz . Port B is a Level-2 ATM (single PHY) or single PHY port and can operate at 50 MHz . It is restricted to an 8 -bit data bus when port A is in multi-PHY mode. Port C is a Level-1 ATM or PHY port.

In addition to ports A, B, and C, the UTOPIA accepts cells from the TX_SAR and routes cells to the RX_SAR and to the data cell FIFO in external control memory.

### 4.5.2 UTOPIA Interfaces

Each of the three ports is divided into two portions: a receive portion and a transmit portion. The TX_SAR and the receive portions are each connected to a 4-cell FIFO. These FIFOs are read on a round-robin basis by the Cell Router (See "Cell Router" on page 76.).

The RX_SAR and the transmit portions are each connected to a 32 -cell FIFO.

The ports are configurable with the following options:

- Port A's transmit portion can be ATM, PHY, with a 16 -bit or 8 -bit data bus.
- Port A's receive portion can be ATM or PHY, with a 16-bit or 8-bit data bus.
- Port A can be Level-2 multi-PHY.
- Port B's transmit portion can be ATM or PHY, with a 16 -bit or 8-bit data bus.*
- Port B's receive portion can be ATM or PHY, with a 16-bit or 8-bit data bus.*
- Port C's transmit portion can be ATM or PHY, with an 8-bit data bus.
- Port C's receive portion can be ATM or PHY, with an 8-bit data bus.
*When Port A is in Level-2 multi-PHY mode, Port B must have an 8-bit data bus.
Each receive interface can be independently enabled or disabled. If disabled, the receive interface will stop accepting cells after the current cell has been received.

When the transmit portions of a port are in PHY mode, the SOC, data bus, and parity output pins can be tristated when the port is not selected. This allows the MT90503 to share a data bus, SOC, and parity lines with other devices (i.e. independent ENB signals and CLAV signals for each PHY device, controlled by a single ATM device).

In the case of a receive PHY, the generation of the rx_clav signal is independent of the state machine. The rx_clav signal is asserted high at any time when a complete cell can be received. Thus as soon as the first byte of a cell is received, and there is no room for another cell in the input FIFO, the rx_clav signal will be asserted low. In the case of a Level-2 PHY, the rx_clav's will only be driven when the address was placed on the bus during the previous cycle.

### 4.5.3 Errors on received cells

If the MT90503 receives a short cell on any one of its three ports, the cell will be discarded, and a new cell will be started when the second SOC signal is set.

If the SOC is not set after the 53rd byte of a received cell, subsequent bytes are ignored until a a new SOC is received.

Data received on all of the three ports is examined for parity errors and an interrupt is raised if an error is found. Cells are not discarded if a parity error is detected. Register 0304h indicates on which port the parity error is detected.

The ATM HEC is not examined on received cells.

### 4.5.4 Transmit and Receive State Machines for ATM and PHY Modes



Figure 34 - ATM Mode State Machines


Figure 35 - PHY Mode State Machines

### 4.5.5 Cell Router

Cells are read on a round-robin basis from the four input FIFOs. Cells from the TX_SAR that contain a '0' in the MSB of the TXD field (see Figure 36 on page 76) are written into the output FIFOs designated by the three LSBs of the TXD field. Cells from the TX_SAR that contain a ' 1 ' in the MSB of the TXD field and cells from the RXA, RXB, and RXC ports are handled based on the VPI/VCI of the cell. The TXD field for a cell is the same as the TXD field in the TX_SAR control structure that created the cell (see Figure 29 on page 59).

Cells written into the RX_SAR output FIFO can be directed to the SAR portion (cells to be formatted into TDM streams), to the data cell portion (to be examined by the CPU), or to both. This is indicated by the RXD field in cells written into the RX_SAR output FIFO.


Figure 36 - Cell Format for cells in internal UTOPIA input and output cell FIFOs


Figure 37 - Cell Router Flow

### 4.5.6 Match \& Mask for cell routing

Cells are designated as "known" or "unknown" based on the result of a match \& mask (configurable for each port). For a cell to be considered "known" all VPI/VCI bits whose corresponding mask bit is '1' must have the value contained in the corresponding bit of the match register. The match and mask registers for Port A are 0328h, 032Ah, 032Ch and 032Eh. Likewise the match and mask register for Port B are 0342h, 0344h, 0348h and 034Ah. For Port C, the registers are 0368h, 036Ah, 036Ch and 036Eh.


Figure 38 - Match \& Mask Example
In addition, the Cell Router can be configured to eliminate null cells (those with VPI = 0 and $\mathrm{VCI}=0$ ). For the purpose of null cell elimination, the NNI can be included on a per-port basis (see registers 0300h and 0302h).

Unknown non-OAM cells and/or unknown OAM cells can be discarded or directed to one or more output FIFOs. All unknown non-OAM cells from a port are discarded or directed to the same location(s) and all unknown OAM cells from a port are discarded or directed to the same location(s). Unknown cells can be directed differently for each port on which they were received. Unknown cells cannot be sent to the SAR portion of the RX_SAR. The routing of unknown cells is set in registers 03A2h and 03A4h.

Known cells are handled according to the LUT (Look-Up Table) entry for the cell's VPI/VCI.

### 4.5.6.1 Look-Up Tables Entries

LUT entries direct cells with known VPI/VCIs to either be discarded or placed in one or more of five possible destinations: the four output FIFOs and the data cell FIFO in external control memory, by way of the RX_SAR FIFO. OAM cells can be directed independently of non-OAM cells with the same VPI/VCI. OAM cells cannot be directed to the SAR portion of the RX_SAR.

LUT entries can be either 4- or 8-bytes long (short or long LUT entries, set in register 302h). All look-up table entries in all three LUTs are the same size. 8 -byte entries are only required if header translation is to be performed for one or more VCs. Cells undergoing header translation have their NNI bits, the remaining VPI bits and/or the VCI bits replaced by the corresponding bits in the LUT entry and are then either discarded or sent to one or more of the possible destinations. VCs that undergo header translation are not directed to the SAR portion of the RX_SAR.

Clock recovery information can be gathered from up to two VCs by setting bit A in one LUT entry and bit B in the same or another LUT entry. A maximum of one VC can have bit A set and a maximum of one VC can have bit B set.


Figure 39 - Short and Long Look-Up Table Entries

### 4.5.6.2 LUT Addressing

A LUT base address exists for each of the three ports (registers 0320h, 0340h, 0360h). The LUT base addresses for two or more ports can be the same. An identifier for a VC is created by concatenating any number of LSBs from the VPI and LSBs from the VCI, to a maximum of 16 bits. The number of VCI bits used is programmed in registers 0324h, 0344h and 0364h for Port A, B, C respectively. The total number of bits in the identifier is programmed in registers $0322 \mathrm{~h}, 0342 \mathrm{~h}$ and 0362 h for Port $\mathrm{A}, \mathrm{B}$ and C respectively. The identifier is then appended with either two or three zeros (for either short or long LUT entries).

Finally, this value is used as the byte-pointer to the LUT entry, offset from the LUT base address for that port.


Figure 40 - VPI/VCI Concatenation and LUT Entry Address Example

### 4.5.6.3 UTOPIA Clocks

Each of the three ports must have a clock to operate the receive interface and a clock to operate the transmit interface. Two or more clocks may have the same source. These clocks can either be input to the MT90503 from an external source or output from the MT90503, from one of three internal UTOPIA clocks. For each port the transmit clock and receive clock must be configured to be either both input or both output. An exception is Port C where both transmit clock and receive clock must be input only.

The source of the each of the three internal UTOPIA clocks can be one of eight clocks: mclk, fast_clk, or any of the six UTOPIA clocks (rxa, rxb, rxc, txa, txb, and txc). The selected clock is divided by $n$, an integer from 1 to 16 , and can be inverted.

Other parts of the UTOPIA module, including the look-up engine, the TX_SAR portion and the RX_SAR portion operate off of mclk.

### 4.5.7 LED Operation

The UTOPIA module generates two LED signals for Port A (pins D2, H5) and two LED signals for Port B (pins W5, T5) in order to indicate the status of the A and B ports. The status conditions are: idle, presence of traffic, or PHY alarm. When a port is in an idle state, both its LEDs are illuminated. If $R X$ traffic (other than null cells) is flowing, then the RX LED for that port will flash; If TX traffic (other than null cells) is flowing, then the TX LED for that port will flash. If a PHY alarm is detected, the TX LED is on and the RX LED is off. The polarity of the LED signals is active-low, i.e., a ' 0 ' will turn the LED on. The frequency of the LEDs is programmed in registers 0120 h and 0122h while the LEDs are enabled in register 0302h.


Figure 41 - UTOPIA Clock Generation

### 4.5.8 UTOPIA FIow Control

The UTOPIA module contains the ability to prevent cells in the 4 -cell input FIFOs (RXA, RXB, RXC, and TX_SAR) from being handled by the UTOPIA module in the case that the 32 -cell output FIFOs (TXA, TXB, TXC, and RX_SAR) exceed programmable levels.

An input FIFO will be blocked when the level of any output FIFO exceeds the level set for that combination of output FIFO and input FIFO. The levels can be set independently to 1 to 31 cells or to 0 , which means no flow control will be exerted (see registers 0338h-033Ah, 0358h-035Ah, 0378h-037Ah). Cell arrival counters and cell departure counters for each port, stored in registers (0330h-0336h, 0350h-0356h, 0370-0376h, 0390-0396h), are used to monitor the fill levels of each output FIFO.

### 4.5.9 External Interface Signals

Due to the different possible configurations of the UTOPIA ports, the functions of some pins change, depending on the configuration. Some unused data pins when port A and/or port B are in 8 -bit mode become general purpose inputs and/or outputs. When Level-2 addressing is in place for port A, portions of the port B data buses are used for as port A addressing pins. The function of the clav (cell available) and enb (enable data transfer) pins alternate when the port is in ATM mode or PHY mode (Figure 42 on page 81).

Please note that the I/O direction of the pins remains the same.


Figure 42 - External UTOPIA Interface

### 4.6 Clock Recovery Module

### 4.6.1 Overview

The purpose of the clock recovery module is to synchronise the TDM clock domain of the MT90503 with other devices on the network through information transmitted across the ATM link. Clock recovery is necessary only when the MT90503 is operating as the TDM clock master.

The clock recovery system is composed of several sub-components:

### 4.6.1.1 Two Point Generation Modules

The point generation modules permit SRTS and/or adaptive clock recovery to be performed by the MT90503. The two modules provide the flexibility to have a back up clock recovery process operating from another VC and/or of the complementary type. These modules generate points which are written to control memory and subsequently used by the clock recovery algorithm.

### 4.6.1.2 One SRTS (synchronous residual time stamp) Generating Module

The SRTS generating module, is employed to generate the 4-bit outgoing RTS value.

### 4.6.1.3 Three Integer Divisor Clock Modules

The integer divisor clock (idclk) modules are designed to manipulate any incoming or internal clocks and produce an 8 kHz idclk. They provide the flexibility of inverting the clock before or after division, option of setting the duty cycle to $50 \%$, 16 -bit division and checking of the frequency (within a desired range). Along with the idclk_loss signals, idclk's are suited to being sent to an external PLL.

### 4.6.1.4 Two Precise Clock Modules

The precise clock (pclk) modules are used to divide mclk down to pclk_a and pclk_b. Each module has a pclk_loss indicator which can be employed in tandem with the pclk signal to be routed to an external PLL. The division of mclk is performed with a 16-bit integer and a 16-bit fraction, allowing for precise specification of the pclk frequency. These modules are ideally suited to being programmed by the clock recovery algorithm to generate the recovered clock.

### 4.6.1.5 Eleven Multiplexers

The 11 multiplexers are: recov_a to recov_h, ct_netref1, ct_netref2 and local_netref_16m. recov_a through recov_h and ct_netrefx are general l/O pins. Local_netref_16m is an internal node, used as the master clock for the TDM section of the MT90503.

### 4.6.2 Multiplexers

There are eleven multiplexers (see Table 28, "Source Selection," on page 85) with 36 possible inputs each (see Table 29, "idclk_a Register," on page 86):

| Signal | Address | Bits |
| :---: | :---: | :---: |
| recov_a | 0860 h | $13: 8$ |
| recov_b | 0860 h | $5: 0$ |
| recov_c | 0862 h | $13: 8$ |
| recov_d | 0862 h | $5: 0$ |
| recov_e | 0864 h | $13: 8$ |
| recov_f | 0864 h | $5: 0$ |
| recov_g | 0866 h | $13: 8$ |
| recov_h | 0866 h | $5: 0$ |
| ct_netref1 | 0868 h | $13: 8$ |
| ct_netref2 | 0868 h | $5: 0$ |
| local_netref_16m | 086 Ah | $5: 0$ |

Table 27 - Multiplexer Registers


Figure 43 - Multiplexer

| Source Select Number | Hex Values | Source | Description |
| :---: | :---: | :---: | :---: |
| 000000 | $0 \times 00$ | Logical 0 | Ground |
| 000001 | $0 \times 01$ | Logical 1 | Vcc |
| 000010 | $0 \times 02$ | Tri-state | high-impedance |
| 000011 | $0 \times 03$ |  | Reserved |
| 000100 | 0x04 | pclk_a | precise clock |
| 000101 | 0x05 | pclk_valid_a | precise clock valid |
| 000110 | $0 \times 06$ | pclk_b | precise clock |
| 000111 | $0 \times 07$ | pclk_valid_b | precise clock valid |
| 001000 | 0x08 | idclk_a | integer divisor clock |
| 001001 | $0 \times 09$ | idclk_valid_a | integer divisor clock valid |
| 001010 | $0 \times 0 \mathrm{~A}$ | idclk_b | integer divisor clock |
| 001011 | $0 \times 0 \mathrm{~B}$ | idclk_valid_b | integer divisor clock valid |
| 001100 | 0x0C | idclk_c | integer divisor clock |
| 001101 | 0x0D | idclk_valid_c | integer divisor clock valid |
| 001110 | 0x0E |  | Reserved |
| 001111 | 0x0F |  | Reserved |
| 010000 | $0 \times 10$ | recov_a | external I/O pin |
| 010001 | $0 \times 11$ | recov_b | external I/O pin |
| 010010 | $0 \times 12$ | recov_c | external I/O pin |
| 010011 | $0 \times 13$ | recov_d | external I/O pin |
| 010100 | $0 \times 14$ | recov_e | external I/O pin |
| 010101 | 0x15 | recov_f | external I/O pin |
| 010110 | 0x16 | recov_g | external I/O pin |
| 010111 | $0 \times 17$ | recov_h | external I/O pin |
| 011000 | $0 \times 18$ | ct_c8 | active clock: ct_c8_a or _b |
| 011001 | $0 \times 19$ | ct_c8_a | TDM clock |
| 011010 | $0 \times 1 \mathrm{~A}$ | ct_c8_b | TDM clock |
| 011011 | 0x1B | ct_frame | active clock: ct_frame_a or _b |
| 011100 | 0x1C | ct_frame_a | TDM clock |
| 011101 | $0 \times 1 \mathrm{D}$ | ct_frame_b | TDM clock |
| 011110 | 0x1E | ct_netref1 | TDM clock |
| 011111 | $0 \times 1 \mathrm{~F}$ | ct_netref2 | TDM clock |
| 100000 | $0 \times 20$ | ref_vca | cell arrival on VC A |
| 100001 | $0 \times 21$ | ref_vcb | cell arrival on VC B |
| 100010 | $0 \times 22$ | phy_alm_a | PHY Alarm UTOPIA A |
| 100011 | $0 \times 23$ | phy_alm_b | PHY Alarm UTOPIA B |
| 100100 | $0 \times 24$ | mclk | master clock |

Table 28 - Source Selection

### 4.6.3 Integer Divisor Clocks (idclk)

There are three idclk modules in the MT90503. Each module consists primarily of a 16-bit integer clock divider, but also has several clock manipulation circuits. These modules have the ability to flag an interrupt (if enabled - 0884h, 08A4h, 08C4h) if the input frequency is above or below a desired range, invert the clock's polarity (before and after division) and set the percentage duty cycle to $50 \%$. idclk_loss_x signals the loss (or incorrect frequency) of an input synchronisation clock, allowing the clock to easily be output to an external PLL. idclk_a is configured in the register range of 0880h-0894h as described in Table 29.

These modules are ideally suited to dividing down the highly accurate reference clock ( $f_{n}$ ) required when performing SRTS clock recovery.

| Register | Bits | Name | Description |
| :---: | :---: | :--- | :--- |
| 0880 h | $[0]$ | divisor_load_now | set to load new value |
| 0880 h | $[1]$ | divisor_reset | 0 for reset, 1 for normal operation |
| 0880 h | $[2]$ | even_duty_cycle_select | when 1, $50 \%$ duty cycle is generated |
| 0880 h | $[3]$ | input_invert_select | invert clock before dividing |
| 0880 h | $[4]$ | output_invert_select | invert clock after dividing |
| 0880 h | $[13: 8]$ | input_source_select | see Table 28 |
| 0882 h | $[5: 0]$ |  | status registers |
| 0884 h | $[5: 0]$ |  | interrupt enabling |
| 0886 h | $[5: 0]$ |  | manual setting of status registers |
| 0888 h | $[5: 0]$ | ext_loss_source_select | external output of input clock status. See Table 28. |
| 0888 h | $[6]$ | ext_loss_source_polarity | if set to '1', source loss is active high |
| 0888 h | $[7]$ | output_loss_polarity | if '1' output loss is active high |
| 088 Ah | $[15: 0]$ | clk_div | denominator of clock divider |
| 0890 h | $[15: 0]$ | freqchck_div | denominator of frequency check divider |
| 0892 h | $[15: 0]$ | freqchck_max_mclk_cycles | max \# of mclk's between rising edges of the input clock divided by <br> freqchck_div, if failure occurs, freq_too_low (0882h) will be set. |
| 0894 h | $[15: 0]$ | freqchck_min_mclk_cycles | min \# of mclk's between rising edges of the input clock divided by <br> freqchck_div, if failure occurs, freq_too_high (0882h) will be set. |

Note: idclk_b and idclk_c have a corresponding set of registers in the ranges of 08A0h - 08B4h and 08C0h - 08D4h respectively.

Table 29 - idclk_a Register


Figure 44 - Integer Clock Processor

### 4.6.4 Precise Clocks (pclk)

The MT90503 has two digital PLL (pclk) modules. Using mclk as a source, the module will divide it with a 16 -bit integer and optional 16 -bit fraction. The 16 -bit fraction allows more precise specification on the output frequency. Using the optional 16-bit faction in a typical configuration, mclk $=80 \mathrm{MHz}$, pclk_int_a $=10000$ and pclk_a $=8 \mathrm{kHz}$, will increase the precision from 100 ppm to 1.5 ppm . Also, using the fractional divider will reduce the maximum jitter to one mclk period ( 12.5 ns for 80 MHz mclk ). Set pclk_frc to 0 if no jitter insertion by the pclk module desired. The divider can be programmed dynamically and has a maximum response time of $125 \mu \mathrm{~s}$.

The following registers are applicable:

| Register | Name | Description |
| :---: | :--- | :--- |
| 0820h[5] | adapsrts0_pclk_loss | indicates state of clock |
| 0820h[6] | adapsrts0_pclk_divisor_load_now | when set, pclk_div and pclk_frc are loaded into digital PLL. |
| 0820h[7] | adapsrts0_pclk_divisor_reset | when '0' digital PLL is in reset state |
| 0830h | adapsrts0_pclk_div | integer divider of pclk_a |
| 0832h | adapsrts0_pclk_frc | fractional divider of pclk_a |
| 0820h[5] | adapsrts1_pclk_loss | indicates state of clock |
| 0840h[6] | adapsrts1_pclk_divisor_load_now | when set, pclk_div and pclk_frc are loaded into digital PLL. |
| 0840h[7] | adapsrts1_pclk_divisor_reset | when '0' digital PLL is in reset state |
| 0850h | adapsrts1_pclk_div | integer divider of pclk_b |
| $0852 h$ | adapsrts1_pclk_frc | fractional divider of pclk_b |

Table 30 - pclk registers

The following equation illustrates the derived frequency of pclk from mclk:

$$
f_{p c l k}=\frac{f_{m c l k}}{p_{\text {clk }}^{d i v}}+\frac{p c l k_{f r c}}{65536}
$$

### 4.6.5 Point Generation

The function of the point generation module is to place points in external memory which have been generated by either the SRTS or Adaptive clock recovery methods. These points express the rates of the master device's TDM clock (through a time stamp or the cell rate), the rate of the slave (performing the clock recovery) device's master clock and the rate of the slave device's pclk. This allows the clock recovery algorithm to evaluate the respective rates and make corrections to the pclk in order to synchronise with the master device.

There are two point generation modules. Each can be configured for SRTS or adaptive clock recovery. The two modules each have a separate point generation process, separate timing reference and each is associated with one pclk module (i.e. pclk_a with adapsrts0 and pclk_b with adapsrts1). This allows switching between clock recovery types and/or sources on the fly. The A and B bits of the UTOPIA look-up table determine which VC generates ref_vca and ref_vcb (see section 4.5.6.1 on page 78).

| Register | Bits | Name | Module $^{\text {a }}$ | Description |
| :---: | :---: | :---: | :---: | :--- |
| 0820 h | 0 | adaptive_enable | A | '1' activates adaptive clock recovery |
| 0820 h | 1 | rx_srts_enable | S | '1' activates SRTS clock recovery |
| 0820 h | 2 | ignore_crc | B | '1' ignores CRC of AAL1 byte |
| 0820 h | 3 | ignore_parity | B | '1' ignores parity bit of AAL1 byte |

Table 31 - adapsrts0 Registers ${ }^{\text {b }}$

| Register | Bits | Name | Module ${ }^{\text {a }}$ | Description |
| :---: | :---: | :---: | :---: | :---: |
| 0820h | 4 | ignore_seq_num | B | '1' ignores sequence number of AAL1 byte |
| 0822h | 0 | aal1_crc_error | B | status bit indicating CRC error |
| 0822h | 1 | aal1_bad_parity | B | status bit indicating parity error |
| 0822h | 2 | single_cell_lost | B | status bit indicating single cell loss |
| 0822h | 3 | multi_cell_lost | B | status bit indicating multiple cell loss |
| 0822h | 4 | cell_misinserted | B | status bit indicating a cell misinsertion error |
| 0822h | 5 | timeout_flag | B | status bit indicating interval since last cell exceeds time_out_period (0828h) (reset when timeout ceases) |
| 0822h | 6 | timeout_pulse | B | status bit indicating the interval between two cells has exceeded the time_out_period (0828h) |
| 0822h | 7 | rx_srts_remote_overflow | S | status bit indicating the interval between remote SRTS values received was too short and one was lost |
| 0822h | 8 | rx_srts_local_overflow | S | status bit indicating the interval between local SRTS values received was too short and one was lost |
| 0824h | [8:0] |  |  | Interrupt Enables |
| 0826h | [5:0] | ref_input_select | A | timing reference (usually ref_vca or ref_vcb), see Table 28. |
| 0826h | [13:8] | rx_fnxi_input_select | S | selects fnxi input, see Table 28. |
| 0828h | [15:0] | time_out_period | B | time-out period between two cells (units -1024 mclk cycles) |
| 082Ah | [7:0] | adap_pnt_elim_x | A | "keep 1 point out of X " and write to external memory |
| 082Ch | [15:0] | srts8m8c_div_p | S | pclk must be divided by K in order to match the interval of 8 SRTS carrying cells, where $K=P / Q, P$ is normally the number of frames in the scheduler, i.e. 375 for fully filled structured AAL1 |
| 082Eh | [15:0] | srts8m8c_div_q | S | $Q$ is the number of channels open |

Table 31 - adapsrts0 Registers ${ }^{\text {b }}$ (continued)
a. adaptive relevant register, S - SRTS relevant register, B - relevant to both
b. A corresponding set of registers exists for adapsrts1 from 0840h-084Eh.


Figure 45 - Adaptive Clock Recovery

### 4.6.6 Adaptive Clock Recovery

Adaptive clock recovery is a method which generates a clock based on the rate at which AAL1 cells are arriving. The device acting as the master does not have to structure its cells differently or add any information, it simply transmits CBR data in AAL1 cells. The slave device performs the adaptive clock recovery, based on a clock recovery algorithm using points placed in external memory by the point generation module.

To perform adaptive clock recovery, the point generation module of the MT90503 is normally configured with the cell arrival event (ref_vca or ref_vcb) as its timing reference (source $0 \times 20$ or $0 \times 21$ of Table 28, "Source Selection," on page 85). The VC's which are defined as vca and vcb are recorded in the UTOPIA LUT (see Figure 39 on page 78). The input multiplexer gives the flexibility to use any clock desired. The cell arrival event is received directly from the UTOPIA look-up module. The look-up engine generates a VC-specific pulse and passes on the AAL1 byte of the received cell.

The AAL1 byte is composed of a sequence number (SN), CRC-3 sequence number protection (SNP) and a parity bit. The adaptive module checks for CRC and parity errors. If errors are found, the cell is ignored and the appropriate bit of register 0822h or 0842h is flagged. The sequence number is also verified to determine if cells have been lost. Single cell losses can be compensated for and the single_cell_lost register bit ( 0822 h or 0842h [2]) will be set. Multiple cell losses cannot be compensated, but will be flagged as either multi_cell_lost (0822h or 0842h [3]) or cell_misinserted (0822h or 0842h [4]).


Figure 46 - Adaptive Cell Reception Flow
The point recording portion of the module records three fields. The 'number of the cell' is a 32 -bit field, the 'number of mclk cycles' counted at time of cell arrival is a 32 -bit field and the 'number of pclk cycles' counted when the cell received is a 48 -bit field composed of a 32 -bit integer and 16 -bit fraction. This information, from the three counters (cell, mclk and pclk), 112 bits in total, is written to external control memory in a circular buffer reserved for clock recovery information. (See External Memory Point Format on page 93.) Ratios of this data are compared by the clock recovery algorithm to the desired ratios and correspondingly, corrections are made to the pclk frequency dividers.

It is possible to eliminate some of the timing reference cells sent to memory. This may be done in order to conserve processing power. It is especially useful if the clock recovery VC has a high number of channels (i.e. a high rate of cell arrival). The 8 -bit registers, adap_pnt_elim_x (082Ah and 084Ah) can be programmed to "keep one point out of $X$ ".

Each adaptive module has its own associated pclk generator, allowing the wander of each VC to be tracked with respect to the pclk frequency.

### 4.6.6.1 SRTS Clock Recovery

The Synchronous Residual Time Stamp (SRTS) method of clock recovery is standardised in ITU-T I.363.1, ANSI T1.630 and Bellcore's patent ${ }^{1}$ (U.S. Patent 5260978 (11/93)).

The SRTS method uses a stream of residual time stamps (RTS) to express the difference between a common reference clock ( $f_{n}$ ) and a local service clock ( $f_{s}$ - derived from the local TDM clock, ct_c8_x).

[^1]The point generation modules can both be configured for SRTS clock recovery and receive data simultaneously from 2 VCs. Like adaptive clock recovery, the data is retrieved from the UTOPIA look-up module. The SRTS values however, are spread over 8 cells. As in the adaptive mode, CRC errors, parity errors and missing cells are reported to their respective registers ( 0822 h \& 0842 h [4:0]). This component of the SRTS recovery that receives SRTS data on a VC from an outside source generates "remote" data

The SRTS clock recovery method requires an accurate external reference clock ( $f_{n}$ ) (e.g. a stratum 3 clock). This clock drives the 4 -bit counter fnxi_cnt. This count is compared to a count driven by the precise clock digital PLL. In order to match the interval of 8 SRTS carrying cells, pclk ( 8 kHz ) must be multiplied by K. K is proportional to the number of frames in the scheduler (P) ( 375 for fully filled structured AAL1) and inversely proportional to the number of channels open (Q) (respectively of registers 082Ch and 082Eh for point generation module 0 (adapsrts0)). i.e. K=P/Q. This component of the SRTS clock recovery that compares the pclk generated clock with that of the fnxi clock generates "local" data.

These "local" and "remote" values are written to external memory for the CPU to access.


Figure 47 - Rx SRTS Clock Recovery Module

### 4.6.7 SRTS Transmission

Similar to the SRTS receive side, the generation of SRTS data must take into account the number of frames in the scheduler ( $P$ of register 0818h) and the number of channels in the VC ( $Q$ of register 081Ah). A single VC may be used to carry SRTS or the SRTS values may be broadcast on multiple VCs. These VCs must, however, be of the same format, consistent with the master SRTS VC. These VC's are configured in the Tx SAR (see Table 22 on page 60).


Figure 48 - Tx SRTS Clock Recovery Module

| Register | Bits | Name | Description |
| :---: | :---: | :---: | :--- |
| 0810 h | 0 | enable | '1' enables Tx SRTS |
| 0810 h | $[1: 2]$ | bus_clk_sel | source of timing: '10' - ct_c8_a,'11' - ct_c8_b,'0x' - follows active <br> clock, ct_c8_a or_b |
| 0810 h | $[8: 13]$ | fnxi_input_select | selects source of fnxi see Table 28 |
| 0812 h | 0 | overflow | set if values sent by Tx SRTS is greater than the number of values <br> read by the Tx SAR |
| 0812 h | 1 | underflow | set if the number of Tx SRTS values is less than that read by the Tx <br> SAR |
| 0814 h | $[0: 1]$ |  | Interrupt Enables |
| 0818h | $[0: 15]$ | srts8m8c_div_p | P = number of frames in the scheduler |
| 081Ah | $[0: 15]$ | srts8m8c_div_q | Q = number of channels open |

Table 32 - Tx SRTS Registers

### 4.6.8 External Memory Point Format

Figure 49 on page 94 indicates the format of the circular buffers in external memory that contain SRTS and adaptive point information. In adaptive clock recovery, a point is the information pertinent to the reception of a single cell. In SRTS, a point is the information corresponding to reception of an SRTS value (gathered over eight cells). If the received SRTS value is corrupt (due to errors in the received cells) the valid bit ( V ) will be 0 .


Figure 49 - Clock Recovery Information Buffers

| Field | Name of Field | Bits Used | Description |
| :--- | :--- | :--- | :--- |
| V | Valid Bit | b12 | Set if the accompanying RX SRTS value is valid |
| RX SRTS | Remote SRTS value | b11:b8 | The 4-bit SRTS value from the incoming ATM cells on the <br> designated VC. |
| Local SRTS | Local SRTS value | b3:b0 | The 4-bit SRTS value calculated based on pclk and fnxi |

Table 33 - SRTS Pointer Buffer, Field Description

### 5.0 Memory

### 5.1 Memory Overview

The MT90503 requires external memory for two purposes: control and data memory. The control memory contains information required for: TX_SARs and RX_SARs control structures, transmission schedulers, look up values to map VCs to RX structures. The control memory also stores data cell FIFO's. CAS buffers and clock recovery data is also stored in the control memory. The data memory is employed to store network traffic data for a maximum of 2048 bi-directional TDM channels.

The MT90503 interfaces with the external Data SSRAM via the following pins: 19 address pins (dmem_a), 4 memory bank/chip selection pins (dmem_cs), 16 data pins, (dmem_d), 2 parity pins (dmem_par), 1 R/W(low) pin used for late write memories (dmem_rw), 2 data memory byte write select pins (dmem_bws) and a memory clock (mem_clk). The external Control SSRAMs interface pins: 19 address pins (cmem_a), cmem_a[18] can be configured as a: memory bank/chip selection pin (cmem_cs[1]) or an address pin cmem_a[18], 1 dedicated memory bank/chip selection pin (cmem_cs[0]), 16 data pins, (cmem_d), 1 R/W(low) pin used for late write memories (cmem_rw), 2 control memory byte write select pins (cmem_bws).

The data memory supports up to 4 memory banks up to 512 k words per bank determining a data memory limit of 4 MB . The data memory clock speed gamut is 40 MHz to 80 MHz . Note: recommended mem_clk speed is 80 MHz . The option of a reduced memory capability is also supported. The following SSRAM sizes can be employed: $128 \mathrm{kB}, 256 \mathrm{kB}, 512 \mathrm{kB}$ and 1 MB . The data bus consists of 18 data bits where two data bits are dedicated as parity bits. The parity bits are used to detect underruns in the circular buffers generated on the ATM link and data error detection. The parity check can be disabled to permit non-parity memory compatibility. The MT90503 supports 1, 2, 3 or 4 banks of external memory, each bank having a total capacity ranging from $64 \mathrm{k} \times 18$ bits to $512 \mathrm{k} \times 18$ bits. Therefore the MT90503 can operate with external memory ranging from 128kB to 4MB. The above data memory configuration is initialised via: Data Memory Parity 0 Register, Data Memory Parity 1 Register and Data Memory Configuration Register (0248h, 024Ah \& 024Ch respectively).

The control memory maximum capacity is 512 k words and supports 2 memory banks. The reduced memory capability is supported in the same manner as the data memory. However, if all 19 address bits are employed then the use of 1 memory bank is permitted. Therefore the MT90503 can operate with external control memory ranging from 128 kB to 1 MB . The MT90503 dose not use the parity bits supplied by the control memory. Parity bits can be generated within the MT90503 and are used for error detection. The above control memory configuration is initialised via: Control Memory Parity 0 Register, Control Memory Parity 1 Register and Control Memory Configuration Register (0240h, 0242h \& 0244h respectively).

The MT90503 supports both Pipelined and Flow Through SSRAM employing either: 'normal' or 'Zero Bus Turnaround' (ZBT) operation. When PECL clock is employed 'Late-Write' is supported in 'normal' operation.

The memory clock (mem_clk) interface must be configured to either PECL or TTL. The interface can be initialised via the 'CPU Control Register' address 0100h.

The external memory controller can interface with several types of SSRAM, but they must support synchronous bus enabling. The SSRAM chip must only enable its data output buffers one cycle after a read (two for pipelined SSRAM), irrespective of the state of the asynchronous output enable pin. A read is indicated by mem_rw high and the appropriate mem_cs asserted. The SSRAM can be a registered input type (Synchronous, Synchronous Flow through or Synchronous Burst) or a registered input/output type (Synchronous pipelined). Although the MT90503 uses synchronous access feature of the memory, it does not use the burst access feature of the memory.

Specific Synchronous SRAM devices may require a turnaround cycle with respect to the bidirectional data bus. The MT90503 can be configured to insert a turnaround cycle between read access and write access. A turnaround cycle can be inserted between read access and read access to other memory banks. The turnaround cycle configuration is initialised via: Control Memory Configuration Register and Data Memory Configuration Register ( 0244 h \& 024Ch respectively). It should be noted that turnaround cycles restrict the
memory bandwidth and therefore the operation MT90503. Maximum throughput is achieved with full clock speed on MCLK and without pipelined synchronous RAM and turnaround cycles.

### 5.2 Memory Map

The location of the absolute starting and ending addresses of the internal and external memories are shown in Table 34 - MT90503 Memory Map. The complete set of internal registers is listed in section 5.5 "Detailed Register Description", on Page 99.

| Start Address | End Address | Name |
| :---: | :---: | :---: |
| 0100h | 01FEh | CPU Registers |
| 0200h | 02FEh | Main Registers |
| 0300h | 03FEh | UTOPIA Registers |
| 0400h | 04FEh | TDM Registers |
| 0500h | 05FEh | TX_SAR Registers |
| 0600h | 06FEh | Scheduler Registers |
| 0700h | 07FEh | RX_SAR Registers |
| 0800h | 08FEh | Clock Registers |
| 0900h | 09FEh | Miscellaneous Registers |
| 0A00h | OAFEh | H. 100 Registers |
| 1000h | 10FEh | TX SAR Input FIFO |
| 1400h | 14FEh | UTOPIA Port A Input FIFO |
| 1800h | 18FEh | UTOPIA Port B Input FIFO |
| 1C00h | 1CFEh | UTOPIA Port C Input FIFO |
| 2000h | 27FEh | RX SAR Output FIFO |
| 3000h | 33FEh | UTOPIA Port A Output FIFO |
| 4000h | 43FEh | UTOPIA Port B Output FIFO |
| 5000h | 53FEh | UTOPIA Port C Output FIFO |
| 8000h | BFFEh | TDM Channel Association Memory |
| 200000h | 2FFFFEh | External Control Memory |
| 400000h | 7FFFFEh | External Data Memory |

Table 34 - MT90503 Memory Map

### 5.3 Memory Controllers

Two memory controllers for external SSRAM exist in the MT90503: one for data memory and one for control memory. The memory controller blocks of the MT90503 reside between the internal blocks and the external memory. They receive memory access requests from the internal blocks (TDM, TX_SAR, RX_SAR, and CPU interface modules) and service them by reading data from, or writing data to, the external memories.

### 5.3.1 Data Memory

The data memory contains

- TX/RX circular buffers: one per channel, each with a programmable size of $128,256,512$, or 1024 words.

The data memory can be up to 4 MB in size. This allows 2048 TDM channels to each have a maximum-size circular buffer: 2048 TDM channels * 2048 bytes/channel $=4 \mathrm{MB}$. A parity bit, necessary to detect underruns on incoming TDM data, can be disabled, allowing non-parity memory to be used. The parity bit, when enabled, is also used for error detection.

The data memory can be distributed between one and four banks. 16 to 19 address bits are required to access the 128 KB to 1 MB banks of data memory.

### 5.3.2 Control Memory

The control memory contains

- TX control structures: one per VC, minimum of 12 words each, maximum of 48 KB total
- RX control structures: one per VC, minimum of 14 words each, maximum of 56 KB total
- Transmit event schedulers: up to fifteen, typical sizes are 6 to 150 KB per enabled scheduler
- Look-up table: three LUTs, each with one entry per known VC, 4 or 8 bytes per entry
- Data cell FIFO: two FIFOs, each programmable in length from 4 to 16384 cells, 32 words per cell
- CAS change buffers: 1 to 32768 words in size
- Silent tone buffers: 1 to 32768 words in size
- Clock recovery point buffers: 9216 words in size
- Error message buffer: programmable length of 0 to 65536 error report structures, 4 words per structure.

The control memory can be up to 1 MB in size. A parity bit is used for error detection. The control memory can be distributed in either one or two banks. As with the data memory, 16 to 19 address bits are employed; when 19 address bits are used, only one memory bank can be supported.

### 5.3.3 Data Memory Controller

There are five agents which interact with the data memory controller: TX_SAR, RX_SAR, TDM transmit, TDM receive, and CPU (through the CPU interface). Of these five agents, all but the CPU send their accesses to their internal cache, one for each agent, from which the data is written into the data memory when the data memory bus is available. The internal caches are capable of buffering up to 128 words each.

| Agent | Access types |
| :---: | :---: |
| TX_SAR | Reads |
| RX_SAR | Writes |
| TDM transmit | Writes |
| TDM receive | Reads |
| CPU | Reads and writes |

Table 35-Types of Data Memory accesses for each agent
The CPU has the highest priority on the write accesses and writes whenever it is flagged to do so. The priority arbitration used for accesses to the CPU is described in more detail the CPU module section (see See 4.1 on page 26).

The memory controller generates the CRC-32 needed for each AAL5-VTOA cell.

### 5.3.4 Control Memory Controller

Unlike the data memory controller, the control memory controller does not contain internal caches. Agents' read/write requests are granted on a cycle-by-cycle basis and each agent waits until its request has been completed. Each agent has its own access port allowing it to communicate address, data, r/w, and write enable information with the memory controller. The agents are identified in Table 36-on page 98.

Arbitration between the agents for control memory access is as follows: an agent can request a read or write at any time. The control memory will continue in its current mode (read or write) until no requests of that type are pending. Then, it will switch and service requests of the other mode. The CPU has the highest priority if the memory controller is handling accesses of the mode that the CPU requests. Other agents are granted access on a first-come first-served basis if the memory controller is handling accesses of the mode the agent requests.

| Agent | Access types |
| :---: | :---: |
| TX_SAR | Reads and writes |
| RX_SAR | Reads and writes |
| UTOPIA (LUT) | Reads and writes |
| CPU | Reads and writes |

Table 36 - Types of Control Memory Accesses For Each Agent

### 5.4 Register Overview

This section describes the MT90503's internal registers. An 8 KB memory block is reserved for the register mapping.

The register descriptions are grouped in the following sections:

- CPU Module Interface
- Main Registers
- UTOPIA Module
- TDM Module
- TX_SAR Module
- RX_SAR Module
- Clock Registers
- Miscellaneous Registers
- H.100/H. 110 Bus Registers


### 5.5 Detailed Register Description

### 5.5.1 CPU Registers

Address: 100h
Label: control
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| nreset_registers | 0 | RW | Controls the reset for the MAINREG module and certain CPU <br> functions. '11' out of reset. '0' in reset. Should not be removed <br> unless mclk is present. |
| nreset_chip_mclk_src | 1 | RW | Resets all other parts of the chip. '1' out of reset. '0' in reset. Should <br> not be removed unless mclk is present. |
| nreset_txa_clk_mclk_src | 2 | RW | nreset for UTOPIA TXA clock (syncronized on mclk_src). This <br> reset should not be removed unless the corresponding clock is <br> present. |
| nreset_txb_clk_mclk_src | 3 | RW | nreset for UTOPIA TXB clock (syncronized on mclk_src). This <br> reset should not be removed unless the corresponding clock is <br> present. |
| nreset_txc_clk_mclk_src | 4 | RW | nreset for UTOPIA TXC clock (syncronized on mclk_src). This <br> reset should not be removed unless the corresponding clock is <br> present. |
| nreset_rxa_clk_mclk_src | 5 | RW | nreset for UTOPIA RXA clock (syncronized on mclk_src). This <br> reset should not be removed unless the corresponding clock is <br> present. |
| nreset_rxb_clk_mclk_src | 6 | RW | nreset for UTOPIA RXB clock (syncronized on mclk_src). This <br> reset should not be removed unless the corresponding clock is <br> present. |
| nreset_rxc_clk_mclk_src | 7 | RW | nreset for UTOPIA RXC clock (syncronized on mclk_src). This <br> reset should not be removed unless the corresponding clock is <br> present. |
| reserved | 8 | RW | For future use: reserved for mem_clk oe (TTL) |

Table 37-CPU Control Register

Address: 100h
Label: control
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| reserved | 9 | RW | For future use: reserved for PECL oe |
| mem_clk_o_enable | 10 | RW | Enables the mem_clk TTL output to toggle, active high |
| mem_clk_pecl_enable | 11 | RW | Enables the mem_clk PECL output to toggle, active high |
| mem_clk_input_sel | 12 | RW | '0' = mem_clk_i pin, '1' = mem_clk PECL pins |

Table 37 - CPU Control Register (continued)

Address: 102h
Label: status0
Reset Value: 0000h

| Label | Bit Position | Type | Description |
| :--- | :--- | :--- | :--- |
| reserved | 0 | RO | Reserved. Always read as "0". |
| reserved | 1 | RO | Reserved. Always read as "0". |
| reserved | 2 | RO | Reserved. Always read as "0". |
| internal_read_timeout | 3 | ROL | Goes high if a read access has been active for 32k mclk_src <br> cycles without completion. Fatal chip error. |
| inmo_read_done | 4 | ROL | Indicates that an extended indirect access has completed. <br> This is used to indicate to the host that read data is ready. |
| reserved | $15: 5$ | ROL | Reserved. Always read as "0000_0000_000" |

Table 38-CPU Status Register

Address: 104h
Label: status0_ie
Reset Value: 0 $\overline{000} \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| reserved | 0 | RO | Reserved. Always read as "0". |
| reserved | 1 | RO | Reserved. Always read as "0". |
| reserved | 2 | RO | Reserved. Always read as "0". |
| internal_read_timeout_ie | 3 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| inmo_read_done_ie | 4 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| reserved | $15: 5$ | RO | Reserved. Always read as "0000_0000_000". |

Table 39-CPU Interrupt Enable Register

Address: 10EH
Label: counters
Reset Value: 0000h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| emul_mode | $1: 0$ | EMO | Indicates state of counters and status bits. "00" = normal <br> mode, "01" = reset, " $1 \times$ " = test mode. These bits are only <br> present for tests and should never be used. |
| reserved | $15: 2$ | RW | Reserved. Must be "0000_0000_0000_00" |

Table 40 - CPU Counter Register

Address: 120h
Label: led1
Reset Value: 3FD0h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| mclk_src_freq[6:0] | $6: 0$ | RW | mclk_src Frequency in MHz |
| led_flash_freq[8:0] | $15: 7$ | RW | Determines the time in ms that the LEDs will be turned off <br> to indicate link activity. |

Table 41 - LED1 Register

Address: 122h
Label: led2
Reset Value: 0000h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| led_test_mode | 0 | RW | If '0', the LED Flashing time will be determined in ms. If <br> '1', the LED Flashing time will be determined in us. |
| reserved | $15: 1$ | RW | Reserved. Must be "0000_0000_0000_000" |

Table 42 - LED2 Register

| Address: 128h <br> Label: pll_conf <br> Reset Value: 00A1h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit Position | Type |  |
| pll_div_x | $2: 0$ | RW | Divides mclk_src before entering the REF pin of the <br> fast_clk PLL. Together, pll_div_x and pll_div_y determine <br> the speed of fast_clk, which must be between 160 and 200 <br> MHz. |
| pll_div_y | $5: 3$ | RW | Divides the feedback path from the output pin of the <br> fast_clk PLL. Together, pll_div_x and pll_div_y determine <br> the speed of fast_clk, which must be between 160 and 200 <br> MHz. |
| pll_div_z | $10: 6$ | RW | Divides fast_clk to generate mem_clk |
| pll_bypass | 11 | RW | If '1', mclk_src divided by pll_div_x becomes fast_clk, <br> bypassing the fast_clk PLL |
| nreset_pll_async | 12 | RW | Resets the module that divides mclk_src before being <br> used as the REF pin of the fast_clk PLL |
| reserved | $15: 13$ | RO | Reserved. Always read as "000" |

Table 43 - PLL Configuration Register

Address: 130h
Label: inmo a gpi0
Reset Value: 0000 h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| inmo_a_in[14:0] | $14: 0$ | RO | Current level of the corresponding pin |
| reserved | 15 | RO | Reserved. Always read as "0" |

Table 44 - Intel/Motorola Address Register

## Address: 132h

Label: inmo_a_gpi1
Reset Value: 0000h

| Label | Bit Position | Type |  |
| :---: | :---: | :---: | :--- |
| inmo_a_rise0 | 0 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise1 | 1 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise2 | 2 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise3 | 3 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise4 | 4 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise5 | 5 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise6 | 6 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise7 | 7 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise8 | 8 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise9 | 9 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise10 | 10 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise11 | 11 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise12 | 12 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise13 | 13 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| inmo_a_rise14 | 14 | ROL | This bit is set if the corresponding pin goes for '0' to '1' |
| reserved | 15 | RO | Reserved. Always read as "0:" |

Table 45 - Intel/Motorola Address Rise Register

Address: 136h
Label: inmo_a_gpi2
Reset Value: 0000 h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| inmo_a_fall0 | 0 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall1 | 1 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall2 | 2 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall3 | 3 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall4 | 4 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall5 | 5 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall6 | 6 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall7 | 7 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall8 | 8 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |

Table 46 - Intel/Motorla Address Fall Register

Address: 136h
Label: inmo a gpi2
Reset Value: 0000h

| Label | Bit Position | Type | Description |
| :---: | :---: | :--- | :--- |
| inmo_a_fall9 | 9 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall10 | 10 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall11 | 11 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall12 | 12 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall13 | 13 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_a_fall14 | 14 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| reserved | 15 | RO | Reserved. Always read as "0" |

Table 46 - Intel/Motorla Address Fall Register (continued)

Address: 140h
Label: inmo_d_gpo
Reset Value: 0000 h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| inmo_d_out8 | 0 | RW | This is the value sent out on the corresponding pin, used <br> in conjuction with the OE bit |
| inmo_d_out9 | 1 | RW | This is the value sent out on the corresponding pin, used <br> in conjuction with the OE bit |
| inmo_d_out10 | 2 | RW | This is the value sent out on the corresponding pin, used <br> in conjuction with the OE bit |
| inmo_d_out11 | 3 | RW | This is the value sent out on the corresponding pin, used <br> in conjuction with the OE bit |
| inmo_d_out12 | 4 | RW | This is the value sent out on the corresponding pin, used <br> in conjuction with the OE bit |
| inmo_d_out13 | 5 | RW | This is the value sent out on the corresponding pin, used <br> in conjuction with the OE bit |
| inmo_d_out14 | 6 | RW | This is the value sent out on the corresponding pin, used <br> in conjuction with the OE bit |
| inmo_d_out15 | 7 | RW | This is the value sent out on the corresponding pin, used <br> in conjuction with the OE bit |
| inmo_d_oe8 | 8 | RW | This is OE bit used to drive the corresponding pin. |
| inmo_d_oe9 | 9 | RW | This is OE bit used to drive the corresponding pin. |
| inmo_d_oe10 | 10 | RW | This is OE bit used to drive the corresponding pin. |
| inmo_d_oe11 | 11 | RW | This is OE bit used to drive the corresponding pin. |
| inmo_d_oe12 | 12 | RW | This is OE bit used to drive the corresponding pin. |

Table 47 - Intel/Motorola Data Out Register

## Address: 140h

Label: inmo_d_gpo
Reset Value: 0000 h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| inmo_d_oe13 | 13 | RW | This is OE bit used to drive the corresponding pin. |
| inmo_d_oe14 | 14 | RW | This is OE bit used to drive the corresponding pin. |
| inmo_d_oe15 | 15 | RW | This is OE bit used to drive the corresponding pin. |

Table 47 - Intel/Motorola Data Out Register (continued)

## Address 142h: <br> Label: inmo_d_gpi0 <br> Reset Value: $\mathbf{0} 000 \mathrm{~h}$

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| inmo_d_in8 | 0 | RO | Current level of the corresponding pin |
| inmo_d_in9 | 1 | RO | Current level of the corresponding pin |
| inmo_d_in10 | 2 | RO | Current level of the corresponding pin |
| inmo_d_in11 | 3 | RO | Current level of the corresponding pin |
| inmo_d_in12 | 4 | RO | Current level of the corresponding pin |
| inmo_d_in13 | 5 | RO | Current level of the corresponding pin |
| inmo_d_in14 | 6 | RO | Current level of the corresponding pin |
| inmo_d_in15 | 7 | RO | Current level of the corresponding pin |
| reserved | $15: 8$ | RO | Reserved. Always read as "0000_0000" |

Table 48 - Intel/Motorola Data In Register

Address 144h:
Label: inmo_d_gpi1
Reset Value: 0000h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| inmo_d_rise8 | 0 | ROL | This bit is set if the corresponding pin goes for '0' to ' 1 ' |
| inmo_d_rise9 | 1 | ROL | This bit is set if the corresponding pin goes for '0' to ' 1 ' |
| inmo_d_rise10 | 2 | ROL | This bit is set if the corresponding pin goes for '0' to ' 1 ' |
| inmo_d_rise11 | 3 | ROL | This bit is set if the corresponding pin goes for '0' to ' 1 ' |
| inmo_d_rise12 | 4 | ROL | This bit is set if the corresponding pin goes for '0' to ' 1 ' |
| inmo_d_rise13 | 5 | ROL | This bit is set if the corresponding pin goes for '0' to ' 1 ' |
| inmo_d_rise14 | 6 | ROL | This bit is set if the corresponding pin goes for '0' to ' 1 ' |

Table 49 - Intel/Motorola Data Rise/Fall Register

Address 144h:
Label: inmo_d_gpi1
Reset Value: 0000h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| inmo_d_rise15 | 7 | ROL | This bit is set if the corresponding pin goes for '0' to ' 1 ' |
| inmo_d_fall8 | 8 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_d_fall9 | 9 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_d_fall10 | 10 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_d_fall11 | 11 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_d_fall12 | 12 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_d_fall13 | 13 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_d_fall14 | 14 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |
| inmo_d_fall15 | 15 | ROL | This bit is set if the corresponding pin goes for '1' to '0' |

Table 49 - Intel/Motorola Data Rise/Fall Register (continued)

### 5.5.2 Main Registers

Address: 200h
Label: control
Reset Value: 0000h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| reserved | $14: 0$ | RO | Reserved. Always read as "0000_0000_0000_000" |
| test_status | 15 | TS | When '1', all the status bits in the register will be set. |

Table 50 - Main Control Register

Address: 202h
Label: status0
Reset Value: 0000h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| cmem_parity_error0 | 0 | ROL | Indicates a parity error on data received from the control <br> memory on the data bits [7:0]. This will only be detected if <br> cmem_parity_conf[0] is ' 0 '. |
| cmem_parity_error1 | 1 | ROL | Indicates a parity error on data received from the control <br> memory on the data bits [15:8]. This will only be detected if <br> cmem_parity_conf[1] is ' 0 '. |
| dmem_parity_error0 | 2 | ROL | Indicates a parity error on data received from the data <br> memory on the data bits [7:0]. This will only be detected if <br> dmem_parity_conf[0] is ' 0 '. |

Table 51 - Main Status Register

| Address: 202h <br> Label: status0 <br> Reset Value: $\mathbf{0 0 0 0 h}$ | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| Label | 3 | ROL | Indicates a parity error on data received from the data <br> memory on the data bits [15:8]. This will only be detected if <br> dmem_parity_conf[1] is '0'. |
| dmem_parity_error1 |  | RW | Reserved. Must be "0000_0000_0000" |
| reserved | $15: 4$ |  |  |

Table 51 - Main Status Register (continued)

Address: 204h
Label: status0 ie
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| cmem_parity_error0_ie | 0 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| cmem_parity_error1_ie | 1 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| dmem_parity_error0_ie | 2 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| dmem_parity_error1_ie | 3 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| reserved | $15: 4$ | RW | Reserved. Must be "0000_0000_0000" |

Table 52 - Main Interrupt Enable Register

| Address: counters <br> Label: 20Eh <br> Reset Value: $\mathbf{0 0 0 0 h}$ | Label | Bit <br> Position | Type |
| :---: | :---: | :---: | :--- |

Table 53 - Main Counter Register

| Address: 220h <br> Label: interrupt_flags <br> Reset Value: 00000 h |  |  | Label |
| :---: | :---: | :---: | :--- |
| cpureg_interrupt_active | 0 | RO |  |
| Position |  |  |  | Type | Description |
| :--- |
| mainreg_interrupt_active |
| When '1', indicates that the interrupt request for this |
| module is active. |

Table 54 - Interrupt Flags Register

Address: 220h
Label: interrupt_flags
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| erroralarm_interrupt_active | 14 | RO | The error alarm occurs after the error alarm timeout <br> period or if the error report structure buffer becomes <br> half full. This is used to ensure that the host can empty <br> the error structure buffer in a timely manner without <br> having to poll its fill continuously. |
| interrupt1_treated | 15 | PUL | Software must write this bit to '1', when it has finished <br> servicing interrupts. |

Table 54 - Interrupt Flags Register (continued)

```
Address: 224h
Label: interrupt1_conf
Reset Value: 0000 \(\bar{h}\)
```

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| min_interrupt1_period | $13: 0$ | RW | Number of us between interrupts (minimum). When 0000h, <br> these is no minimum interval between interrupts. <br> Programming this prevents the host from being flooded with <br> interrupts. Each interrupt will last 1 us. |
| interrupt1_polarity | $15: 14$ | RW | Interrupt polarity and output enable. "00"=active low (open- <br> collector); "01"=active high (open-collector); "10" $=$ drive <br> low; "11" = drive high. Drive low or drive high means that the <br> pin's value will not change regardless of internal interrupts. |

Table 55 - Interrupt1 Configuration Register

| Address: 226h <br> Label: interrupt2_conf <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| reserved | $13: 0$ | RW | Reserved. Must be "0000_0000_0000_00" |
| interrupt2_polarity | $15: 14$ | RW | Interrupt polarity and output enable. "00"=active low (open- <br> collector); "01"=active high (open-collector); "10" = drive low; <br> "11" = drive high. Drive low or drive high means that the pin's <br> value will not change regardless of internal interrupts. |

Table 56 - Interrupt2 Configuration Register

## Address: 228h Label: interrupt1_enable Reset Value: 0000h

| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| :---: | :---: | :---: | :---: |
| cpureg_interrupt1_enable | 0 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| mainreg_interrupt1_enable | 1 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| utoreg_interrupt1_enable | 2 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| txreg_interrupt1_enable | 3 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| rxreg_interrupt1_enable | 4 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| miscreg_interrupt1_enable | 5 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| wheelreg_interrupt1_enable | 6 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| clkreg_interrupt1_enable | 7 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| tdmreg_interrupt1_enable | 8 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| mastreg_interrupt1_enable | 9 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| casalarm_interrupt1_enable | 10 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| tdmalarm_interrupt1_enable | 11 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| clkrecovalarm_interrupt1_enable | 12 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| aal0alarm_interrupt1_enable | 13 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| erroralarm_interrupt1_enable | 14 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[0]. |
| reserved | 15 | RW | Reserved. Must always be "0" |

Table 57 - Interrupt1 Enable Register

## Address: 22Ch <br> Label: interrupt2_enable <br> Reset Value: 0000h

| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| :---: | :---: | :---: | :---: |
| cpureg_interrupt2_enable | 0 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| mainreg_interrupt2_enable | 1 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| utoreg_interrupt2_enable | 2 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| txreg_interrupt2_enable | 3 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| rxreg_interrupt2_enable | 4 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| miscreg_interrupt2_enable | 5 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| wheelreg_interrupt2_enable | 6 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| clkreg_interrupt2_enable | 7 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| tdmreg_interrupt2_enable | 8 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| mastreg_interrupt2_enable | 9 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| casalarm_interrupt2_enable | 10 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| tdmalarm_interrupt2_enable | 11 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| Clkrecovalarm_interrupt2_ena ble | 12 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| aal0alarm_interrupt2_enable | 13 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| erroralarm_interrupt2_enable | 14 | RW | When ' 1 ' and the corresponding interrupt active is ' 1 ', an interrupt will be generated on cpu_int[1]. |
| reserved | 15 | RW | Reserved. Must always be "0" |

Table 58 - Interrupt2 Enable Register

Address: 230h
Label: utopia_clock1
Reset Value: $\overline{0000 h}$

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :---: |
| utopia_porta_clk_oe | 0 | RW | 0 '=tri-state the txa_clk and rxa_clk; '1'=drive the txa_clk and rxa_clk. This should only be enabled if the chip is to drive these clocks, and only after the clock generation has been correctly programmed. |
| utopia_portb_clk_oe | 1 | RW | 0 '=tri-state the txb_clk and rxb_clk; '1'=drive the txb_clk and rxb_clk. This should only be enabled if the chip is to drive these clocks, and only after the clock generation has been correctly programmed. |
| utopia_portc_clk_oe | 2 | RW | This bit must be ' 0 ' as UTOPIA port C clocks must be input only. |
| utopia_txa_clk_select[1:0] | 4:3 | RW | "00" = select clock divisor A; "01" = select clock divisor B; "10" = select clock divisor C; "11" = reserved. There are 3 integer clock divisors used to generate the UTOPIA clocks, and each of the 6 UTOPIA clocks can be selected as any one of the 3 . |
| utopia_txb_clk_select[1:0] | 6:5 | RW | "00" = select clock divisor A; "01" = select clock divisor B; "10" = select clock divisor C; "11" = reserved. There are 3 integer clock divisors used to generate the UTOPIA clocks, and each of the 6 UTOPIA clocks can be selected as any one of the 3 . |
| utopia_txc_clk_select[1:0] | 8:7 | RW | Reserved. |
| utopia_rxa_clk_select[1:0] | 10:9 | RW | "00" = select clock divisor A; "01" = select clock divisor B; "10" = select clock divisor C; "11" = reserved. |
| utopia_rxb_clk_select[1:0] | 12:11 | RW | "00" = select clock divisor A; "01" = select clock divisor B; "10" = select clock divisor C; "11" = reserved. |
| utopia_rxc_clk_select[1:0] | 14:13 | RW | Reserved. |
| reserved | 15 | RW | Reserved. Must always be "0" |

Table 59 - Utopia Clock Register

Address: 232h
Label: utopia_gena Reset Value: 0001h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| utopia_clk_diva[5:0] | $5: 0$ | RW | Integer divisor for input UTOPIA clock. Divides the <br> selected clock source. |

Table 60 - Utopia Clock Generation A Register

| Address: 232h <br> Label: utopia_gena <br> Reset Value: 0001h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| utopia_clk_divisor_load_nowa | 6 | PUL | Written to '11 when the source, divisor and inv. have <br> been correctly programmed. |
| utopia_clk_inva | 7 | RW | Inverts the output of the clock divisor |
| utopia_clk_srca[2:0] | $10: 8$ | RW | "000"=txa_clk_in; "001"=txb_clk_in; "010""txc_clk_in; <br> "011"=rxa_clk_in; "100"rrb_clk_in;"101"=rxc_clk_in; <br> "110"=mclk; "111"=fast_clk. |
| utopia_clk_divisor_reseta | 11 | RW | When '0', the clock divisor module is held in reset. |
| reserved | $15: 12$ | RW | Reserved. Must always be "0000" |

Table 60 - Utopia Clock Generation A Register

## Address: 234h <br> Label: utopia_genb <br> Reset Value: $\mathbf{0 0 0 1 \mathrm { h }}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| utopia_clk_divb[5:0] | $5: 0$ | RW | Integer divisor for input UTOPIA clock. Divides the <br> selected clock source. |
| utopia_clk_divisor_load_nowb | 6 | PUL | Written to '1' when the source, divisor and inv. have been <br> correctly programmed. |
| utopia_clk_invb | 7 | RW | Inverts the output of the clock divisor |
| utopia_clk_srcb[2:0] | $10: 8$ | RW | "000"=txa_clk_n; "001"=txb_clk_in; "010"=txc_clk_in; <br> "011"=rxa_clk_in; "100"=rxb_clk_in;"101"=rxc_clk_in; <br> "110"=mclk; "111"=fast_clk. |
| utopia_clk_divisor_resetb | 11 | RW | When '0', the clock divisor module is held in reset. |
| reserved | $15: 12$ | RW | Reserved. Must always be "0000" |

Table 61 - Utopia Clock Generation B Register

| Address: 236h <br> Label: utopia_genc <br> Reset Value: 0001h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| utopia_clk_divc[5:0] | $5: 0$ | RW | Integer divisor for input UTOPIA clock. Divides the <br> selected clock source. |

Table 62 - Utopia Clock Generation C Register

Address: 236h
Label: utopia_genc
Reset Value: 0001h

| Label | $\begin{array}{c}\text { Bit } \\ \text { Position }\end{array}$ | Type | Description |
| :---: | :---: | :---: | :--- |
| utopia_clk_divisor_load_nowc | 6 | PUL | $\begin{array}{l}\text { Written to '1' when the source, divisor and inv. have been } \\ \text { correctly programmed. }\end{array}$ |
| utopia_clk_invc | 7 | RW | Inverts the output of the clock divisor |
| utopia_clk_srcc[2:0] | $10: 8$ | RW | $\begin{array}{l}\text { "000"=txa_clk_in; "001"=txb_clk_in; "010"=txc_clk_in; } \\ \text { "011"=rxa_clk_in; "100"=rxb_clk_in;"101"=rxc_clk_in; } \\ \text { "110"=mclk; "111"=fast_clk. }\end{array}$ |
| utopia_clk_divisor_resetc | 11 | RW | When '0', the clock divisor module is held in reset. |$]$| reserved | $15: 12$ | RW |
| :---: | :---: | :---: |
| Reserved. Must always be "0000" |  |  |

Table 62 - Utopia Clock Generation C Register

Address: 240h
Label: cmem_parity0
Reset Value: $\overline{0000 h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| cmem_parity_conf[1:0] | $1: 0$ | RW | 0' = parity bits; '1' = user data. In normal chip <br> operation, this field should be set to "00", because <br> the chip does not use the parity bits of the control <br> memory. |
| cmem_parity_generation_add_mask <br> [18:16] | $4: 2$ | RW | Mask of address bits [18:16] to be used to generate <br> parity for the control memory. A '1' in one of these <br> bits indicates that the corresponding address bit will <br> be used to generate parity. |
| reserved | $7: 5$ | RO | Reserved. Must always be "0000" |
| cmem_parity_generation_data_mask | $15: 8$ | RW | Mask of data bits to be used to generate parity for <br> the control memory. A '1' in one of these bits <br> indicates that the data bit will be used to generate <br> parity. |

Table 63 - Control Memory Parity0 Register

Address: 242h
Label: cmem_parity1
Reset Value: $\overline{0000 h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| cmem_parity_generation_add_mask <br> $[15: 0]$ | $15: 0$ | RW | Mask of address bits [15:0] to be used to generate <br> parity for the control memory. $A$ '1' in one of these <br> bits indicates that the corresponding address bit <br> will be used to generate parity. |

Table 64 - Control Memory Parity1 Register

Address: 244h
Label: cmem_conf
Reset Value: $\overline{0} 010 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| cmem_add_lines | $1: 0$ | RW | "11" $=1 \mathrm{MB}$ per chip; "10" $=512 \mathrm{~KB}$ per chip; "01" $=256 \mathrm{~KB}$ per <br> chip; "00" $=128 \mathrm{~KB}$ per chip |
| cmem_mem_type | $3: 2$ | RW | "00" $=$ flowthrough ZBT; "01" $=$ flowthrough SSRAM; "10" $=$ <br> pipelined ZBT; "11" $=$ pipelined SSRAM. |
| cmem_rw_ta | 4 | RW | $0=$ no read/write turn-around cycles; $1=1$ read/write turn-around <br> cycle. |
| reserved | $15: 5$ | RO | Reserved. Always read as "0000_0000_000" |

Table 65 - Control Memory Configuration Register

```
Address: 248h
Label: dmem_parity0
Reset Value: \(\overline{0000 h}\)
```

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| dmem_parity_conf[1:0] | $1: 0$ | RW | 0 ' = parity bits; '1' = user data. |
| dmem_parity_generation_add_mask <br> $[20: 16]$ | $6: 2$ | RW | Mask of address bits [20:16] to be used to <br> generate parity for the data memory. A '1' in one <br> of these bits indicates that the corresponding <br> address bit will be used to generate parity. |
| reserved | 7 | RO | Reserved. Always read as "0" |
| dmem_parity_generation_data_mask | $15: 8$ | RW | Mask of data bits to be used to generate parity <br> for the data memory. A '1' in one of these bits <br> indicates that the corresponding data bit will be <br> used to generate parity. |

Table 66 - Data Memory Parity 0 Register

Address: 24Ah Label: dmem_parity1 Reset Value: $\mathbf{0 0 0 0}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| dmem_parity_generation_add_mask <br> $[15: 0]$ | $15: 0$ | RW | Mask of address bits [15:0] to be used to generate <br> parity for the data memory. A '1' in one of these <br> bits indicates that the corresponding address bit <br> will be used to generate parity. |

Table 67 - Data Memory Parity 1 Register

Address: 24Ch Label: dmem_conf Reset Value: $\overline{0010} \mathbf{h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| dmem_add_lines | $1: 0$ | RW | $" 11 "=1 \mathrm{MB}$ per chip; "10" $=512 \mathrm{~KB}$ per chip; "01" $=256 \mathrm{~KB}$ per chip; <br> $" 00 "=128 \mathrm{~KB}$ per chip |
| dmem_mem_type | $3: 2$ | RW | "00" $=$ flowthrough ZBT; "01" $=$ flowthrough SSRAM; "10" $=$ pipelined <br> ZBT; "11" $=$ pipelined SSRAM. |
| dmem_rw_ta | 4 | RW | $0=$ no read/write turn-around cycles; $1=1$ read/write turn-around <br> cycle. |
| reserved | $15: 5$ | RO | Reserved. Always read as "0000_0000_000" |

Table 68 - Data Memory Configuration Register

### 5.5.3 UTOPIA Registers

Address: 300h
Label: control
Reset Value: 0900h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| rxa_ena | 0 | RW | Enables UTOPIA port RXA |
| rxa_sar | 1 | RW | 0 ' = chip acts as PHY, '1' = chip acts as SAR |
| rxa_width | 2 | RW | 0 ' = 8-bit UTOPIA bus, '1' = 16-bit UTOPIA bus |
| rxb_ena | 3 | RW | Enables UTOPIA port RXB |
| rxb_sar | 4 | RW | 0 ' = chip acts as PHY, '1' = chip acts as SAR |
| rxb_width | 5 | RW | 0 ' = 8-bit UTOPIA bus, '1' = 16-bit UTOPIA bus |
| txa_sar | 6 | RW | 0 ' = chip acts as PHY, '1' = chip acts as SAR |
| txa_width | 7 | RW | 0 ' = 8-bit UTOPIA bus, '1' = 16-bit UTOPIA bus |

Table 69 - Utopia Control Register

## Address: 300h

Label: control
Reset Value: 0900h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| txa_multiphy | 8 | RW | 0' = always drive DAT/PAR/SOC pins; '1' = only drive when selected. Only <br> applicable when chip is in PHY mode. |
| txb_sar | 9 | RW | 0 0' = chip acts as PHY, '1' = chip acts as SAR |
| txb_width | 10 | RW | 0' = 8-bit UTOPIA bus, '1' = 16-bit UTOPIA bus |
| txb_multiphy | 11 | RW | 0 0' = always drive DAT/PAR/SOC pins; '1' = only drive when selected. Only <br> applicable when chip is in PHY mode. |
| add_pin_ena | 12 | RW | If '1', UTOPIA port A is level-2 with addressing. When this bit is set, <br> rxb_width and txb_width must be '0'. |
| reserved | 13 | RW | Reserved. Must always be "0" |
| null_cell_elim | 14 | RW | When '1', all cells with VPI and VCI equal to '0' will be discarded. <br> Otherwise, they will be kept and treated normally. |
| test_status | 15 | TS | When '1', all the status bits in the register will be set. |

Table 69 - Utopia Control Register (continued)

Address: 302h
Label: control1
Reset Value: 0008h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| rxc_ena | 0 | RW | Enables UTOPIA port RXC |
| rxc_sar | 1 | RW | 0 ' = chip acts as PHY, '1' = chip acts as SAR |
| txc_sar | 2 | RW | 0 ' = chip acts as PHY, '1' = chip acts as SAR |
| txc_multiphy | 3 | RW | 0 ' = always drive DAT/PAR/SOC pins; '1' = only drive when <br> selected. Only applicable when chip is in PHY mode. |
| uto_output_enable | 4 | RW | 0' Tri-states the ENA or CLAV pin driven by the chip. '1' drives. <br> This bit should be set to '1' 'fter all PHY/SAR register bits have <br> been programmed, but before all rx_ena bits are set. |
| phy_alarm_pol | $6: 5$ | RW | "00" = PHY alarm disabled, "01" = PHY alarm active-high, "10" = <br> PHY alarm active-low, "11" = reserved. |
| long_lut_entries | 7 | RW | When '0', LUT entries are 4 bytes long. When '1', LUT entries are <br> 8 8 bytes long. Long LUT entries should only be used if header <br> translation is to be performed. |
| rxa_nni_null_elim | 8 | RW | 1 ' consider NNI VPI bits for null cell elimination for port RXA. If '0', <br> the high 4 bits of the header (GFC field) will be ignored for null <br> cell elimination. |

Table 70 - Utopia Control1 Register

Address: 302h
Label: control1
Reset Value: 0008h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| rxb_nni_null_elim | 9 | RW | 1' consider NNI VPI bits for null cell elimination for port RXB. If '0', <br> the high 4 bits of the header (GFC field) will be ignored for null <br> cell elimination. |
| rxc_nni_null_elim | 10 | RW | 1' consider NNI VPI bits for null cell elimination for port RXC. If '0', <br> the high 4 bits of the header (GFC field) will be ignored for null <br> cell elimination. |
| phya_tx_led_conf | 11 | RW | When '0', phya_tx_led pin functions as a LED. When '1', functions <br> as GPIO. |
| phya_rx_led_conf | 12 | RW | When '0', phya_rx_led pin functions as a LED. When '1', functions <br> as GPIO. |
| phyb_tx_led_conf | 13 | RW | When '0', phyb_tx_led pin functions as a LED. When '1', functions <br> as GPIO. |
| phyb_rx_led_conf | 14 | RW | When '0', phyb_rx_led pin functions as a LED. When '1', functions <br> as GPIO. |
| reserved | 15 | RW | Reserved. Must always be "0" |

Table 70 - Utopia Control1 Register (continued)

| Address: 304h <br> Label: status0 <br> Reset Value: 0000h | Bit <br> Labition | Type |  |
| :---: | :---: | :---: | :--- |
| Lx_cell_loss | 0 | ROL | Indicates an overflow in the buffer from UTOPIA to the RX SAR |
| outa_cell_loss | 1 | ROL | Indicates an overflow in the output buffer for port TXA |
| outb_cell_loss | 2 | ROL | Indicates an overflow in the output buffer for port TXB |
| outc_cell_loss | 3 | ROL | Indicates an overflow in the output buffer for port TXC |
| cell_loss_rollover | 4 | CRL | Indicates that the cell_loss_counter register has wrapped |
| tx_arr_rollover | 5 | CRL | Indicates that the tx_sar_cell_arrival register has wrapped |
| rx_dep_rollover | 6 | CRL | Indicates that the rx_sar_cell_departure register has wrapped |
| ia_arr_rollover | 7 | CRL | Indicates that the porta_cell_arrival register has wrapped |
| oa_dep_rollover | 8 | CRL | Indicates that the porta_cell_departure register has wrapped |
| ib_arr_rollover | 9 | CRL | Indicates that the portb_cell_arrival register has wrapped |
| ob_dep_rollover | 10 | CRL | Indicates that the portb_cell_departure register has wrapped |
| ic_arr_rollover | 11 | CRL | Indicates that the portc_cell_arrival register has wrapped |

Table 71 - Utopia Status 0 Register

| Address: 304h <br> Label: status0 <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| oc_dep_rollover | 12 | CRL | Indicates that the portc_cell_departure register has wrapped |
| rxa_parity_error | 13 | ROL | Indicates a parity error on port RXA |
| rxb_parity_error | 14 | ROL | Indicates a parity error on port RXB |
| rxc_parity_error | 15 | ROL | Indicates a parity error on port RXC |

Table 71 - Utopia Status 0 Register (continued)

## Address: 30Ch

Label: status2
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| phy_alarma | 0 | ROLO | This bit is set when phya_alm pin is active. Active polarity is <br> determined by phy_alarm_pol register |
| phy_alarmb | 1 | ROLO | This bit is set when phyb_alm pin is active. Active polarity is <br> determined by phy_alarm_pol register |
| test_a_bit | 2 | ROL | Bit is set when look-up engine gives a pulse on clock recovery VC A. <br> Used for tests. |
| test_b_bit | 3 | ROL | Bit is set when look-up engine gives a pulse on clock recovery VC B. <br> Used for tests. |
| reserved | $15: 4$ | RW | Reserved. Must always be "0000_0000_0000" |

Table 72 - Utopia Status 2 Register

Address: 30Eh
Label: status2 ie
Reset Value: $\mathbf{0} \overline{000}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| phy_alarma_ie | 0 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| phy_alarmb_ie | 1 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| test_a_bit_ie | 2 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| test_b_bit_ie | 3 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| reserved | $15: 4$ | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |

Table 73 - Utopia Interrupt Enable 2 Register

## Address: 310h <br> Label: counters <br> Reset Value: 0000h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :---: |
| emul_mode | 1:0 | EMO | Indicates state of counters and status bits. "00" = normal mode, "01" = reset, " 1 x " = test mode. These bits are only present for tests and should never be used. |
| cell_loss_emul | 2 | EMU | When emul_mode = "11", writing '1' to this bit will increment the cell_loss counter by 1 |
| tx_arr_emul | 3 | EMU | When emul_mode = "11", writing '1' to this bit will increment the tx_arr counter by 1 |
| rx_dep_emul | 4 | EMU | When emul_mode = "11", writing ' 1 ' to this bit will increment the rx_dep counter by 1 |
| ia_arr_emul | 5 | EMU | When emul_mode = "11", writing ' 1 ' to this bit will increment the ia_arr counter by 1 |
| oa_dep_emul | 6 | EMU | When emul_mode = "11", writing ' 1 ' to this bit will increment the oa_dep counter by 1 |
| ib_arr_emul | 7 | EMU | When emul_mode = "11", writing '1' to this bit will increment the ib_arr counter by 1 |
| ob_dep_emul | 8 | EMU | When emul_mode = "11", writing '1' to this bit will increment the ob_dep counter by 1 |
| ic_arr_emul | 9 | EMU | When emul_mode = " 11 ", writing ' 1 ' to this bit will increment the ic_arr counter by 1 |
| oc_dep_emul | 10 | EMU | When emul_mode = "11", writing ' 1 ' to this bit will increment the oc_dep counter by 1 |
| reserved | 15:11 | RW | Reserved. Must always be "0000_0" |

Table 74 - Utopia Counters Register

## Address: 312h <br> Label: cell_loss_counters <br> Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| cell_loss | $15: 0$ | CNT | Counts the number of cells lost due to fifo overflows. This <br> includes the output fifos from the look-up engine to the: RX <br> SAR, ports TXA, TXB, TXC. |

Table 75-Cell Loss Counters Register

Address: 320h
Label: porta_look_up_base
Reset Value: $\mathbf{0 0 0 0} \mathbf{h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| luta_base | $15: 0$ | RW | Bits 19:4 of the address of the look-up table for port A. |

Table 76 - Port A Look Up Table Address Register

Address: 322h
Label: porta_num_vpi_vci_bits
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| num_vpi_vci_bits | $15: 0$ | RW | Indicates the total number of bits that are used to decode the <br> look-up address for port A. $0001 \mathrm{~h}=1$ bit used, FFFFh $=16$ bits <br> used. When subtracted from vci_n, this indicates the number <br> of VPI bits used. |

Table 77 - Port A VPI/VCI Identification Register

Address: 324h
Label: porta_concatenation
Reset Value: 0000 h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vci_na | $4: 0$ | RW | Indicates the number of VCI bits used to <br> decode the look-up address for port A. |
| reserved | $15: 5$ | RW | Reserved. Must always be <br> "0000_000_000" |

Table 78 - Port A Concatenation Register

| Address: 328h <br> Label: porta_vpimatch <br> Reset Value: $\mathbf{0 0 0 0} \mathrm{h}$ |
| :--- |
| Label |
| vpi_matcha |
| Position | Type | Bit | Description |
| :---: | :--- |
| reserved | $11: 0$ |
| RW | For a cell from port A to be considered valid, any bits in its <br> VPI whose corresponding bits in reg 32Ah are '1' must have <br> the value contained in this register. |

Table 79 - Port A VPI Match Register

```
Address: 32Ah
Label: porta_vci_mask
Reset Value: 0000h
```

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vpi_maska | $11: 0$ | RW | For a cell from port A to be considered valid, any bits in its VPI <br> whose corresponding bits in this register are '1' must have the <br> value contained in reg 328h. |
| reserved | $15: 12$ | RW | Reserved. Must always be "0000" |

Table 80 - Port A VCI Mask Register

| Address: 32Ch <br> Label: porta_vci_match <br> Reset Value: $\mathbf{0 0 0} \mathbf{0} \mathrm{h}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| vci_matcha | $15: 0$ | RW | For a cell from port A to be considered valid, any bits in <br> its VCI whose corresponding bits in reg 32Eh are '1' must <br> have the value contained in this register. |

Table 81 - Port A VCI Match Register

Address: 32Eh
Label: porta_vci_mask
Reset Value: $\mathbf{0} 00 \overline{0} \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vci_maska | $15: 0$ | RW | For a cell from port A to be considered valid, any bits in its VCI <br> whose corresponding bits in this register are '1' must have the <br> value contained in reg 32Ch. |

Table 82 - Port A VCI Mask Register

Address: 330h
Label: porta_cell_arrival_high
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| ia_arr[31:16] | $15: 0$ | CNT | Freerunning counter of the number of cells received <br> on port A. |

Table 83 - Port A Cell Arrival Counter High Register

Address: 332h

| Label | Bit <br> Position | Type | Description |
| :--- | :--- | :--- | :--- |
| ia_arr[15:0] | $15: 0$ | CNT | Freerunning counter of the number of cells <br> received on port A. |

Table 84 - Port A Cell Arrival Counter Low Register

Address: 334h
Label: porta_cell_departure_high
Reset Value: $0000 \bar{h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| oa_dep[31:16] | $15: 0$ | CNT | Freerunning counter of the number of cells <br> transmitted on port A. |

Table 85 - Port A Cell Departure Counter High Register

Address: 336h
Label: porta_cell_departure_low
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| oa_dep[15:0] | $15: 0$ | CNT | Freerunning counter of the number of cells transmitted on port A. |

Table 86 - Port A Cell Departure Low Register

| Address: 338h <br> Label: porta_overflow0 <br> Reset Value: $\mathbf{0 0 0 0 h}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type |  |
| ia_rx_cell_max | $4: 0$ | RW | If the cell fill of the RX SAR output FIFO becomes greater than this value, <br> cells from the port A input FIFO will be blocked. Oh = no backpressure |
| ia_oa_cell_max | $9: 5$ | RW | If the cell fill of the Port A output FIFO becomes greater than this value, <br> cells from the port A input FIFO will be blocked. Oh = no backpressure |
| ia_ob_cell_max | $14: 10$ | RW | If the cell fill of the Port B output FIFO becomes greater than this value, <br> cells from the port A input FIFO will be blocked. Oh = no backpressure |
| reserved | 15 | RW | Reserved. Must always be "0" |

Table 87 - Port A Overflow0 Register

```
Address: 33Ah
Label: porta_overflow1
Reset Value: 0000h
```

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| ia_oc_cell_max | $4: 0$ | RW | lf the cell fill of the Port C output FIFO becomes greater than this <br> value, cells from the port A input FIFO will be blocked. Oh = no <br> backpressure |
| reserved | $15: 5$ | RW | Reserved. Must always be "0000_0000_000" |

Table 88 - Port A Overflow1 Register

| Address: 33Ch Label: porta_address Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| porta_add | 4:0 | RW | UTOPIA address to which the chip will repond when programmed as a level-2 PHY. |
| reserved | 15:5 | RW | Reserved. Must always be "0000_0000_000" |

Table 89 - Port A Address Register

Address: 340h
Label: portb_look_up_b
Reset Value: $\overline{\mathbf{0} 000} \overline{\mathrm{~h}}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| lutb_base | $15: 0$ | RW | Bits 19:4 of the address of the look-up table for port B. |

Table 90 - Port B Look Up Table Register

Address: 342h
Label: portb_num_vpi_vci_bits
Reset Value: $\mathbf{0} 000 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| num_vpi_vci_bits | $15: 0$ | RW | Indicates the total number of bits that are used to decode the <br> look-up address for port $B$. 0001h $=1$ bit used, FFFFh $=16$ bits <br> used. When subtracted from vci_n, this indicates the number of <br> VPI bits used. |

Table 91 - Port B VPI/VCI Identification Register

Address: 344h
Label: portb_concatenation
Reset Value: $\overline{0} 000 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vci_nb | $4: 0$ | RW | Indicates the number of VCI bits used to decode the look- <br> up address for port B. |
| reserved | $15: 5$ | RW | Reserved. Must always be "0000_0000_000" |

Table 92 - Port B Concatenation Register

Address: 348h
Label: portb_vpi_match
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vpi_matchb | $11: 0$ | RW | For a cell from port B to be considered valid, any bits in <br> its VPI whose corresponding bits in reg 34Ah are '1' must <br> have the value contained in this register. |
| reserved | $15: 12$ | RW | Reserved. Must always be "0000" |

Table 93 - Port B VPI Match Register

Address: 34Ah
Label: portb_vpi_mask
Reset Value: $\overline{0} 00 \overline{0} \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vpi_maskb | $11: 0$ | RW | For a cell from port B to be considered valid, any bits in its <br> VPI whose corresponding bits in this register are '1' must <br> have the value contained in reg 348h. |
| reserved | $15: 12$ | RW | Reserved. Must always be "0000" |

Table 94 - Port B VPI Mask Register

Address: 34Ch
Label: portb_vci_match
Reset Value: $\overline{0} 00 \overline{0} \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vci_matchb | $15: 0$ | RW | For a cell from port B to be considered valid, any bits in its <br> VCI whose corresponding bits in reg 34Eh are '1' must <br> have the value contained in this register. |

Table 95 - Port B VCI Match Register

## Address: 34Eh <br> Label: portb_vci_mask <br> Reset Value: $\overline{0} 00 \overline{0} \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vci_maskb | $15: 0$ | RW | For a cell from port B to be considered valid, any bits in its VCI whose <br> corresponding bits in this register are '1' must have the value contained <br> in reg 34Ch. |

Table 96 - Port B VCI Mask Register

| Address: 350h <br> Label: portb_cell_arrival_high <br> Reset Value: $\mathbf{0} 000 \overline{\mathrm{~h}}$ |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| ib_arr[31:16] | 15:0 | CNT | Freerunning counter of the number of cells received on port B. |

Table 97 - Port B Cell Arrival Counter High Register

Address: 352h
Label: portb_cell_arrival_low
Reset Value: $\mathbf{0 0 0 0} \overline{\mathrm{h}}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| ib_arr[15:0] | $15: 0$ | CNT | Freerunning counter of the number of cells received on <br> port B. |

Table 98 - Port B Cell Arrival Counter Low Register

Address: 354h
Label: portb_cell_departure_high
Reset Value: ${ }^{0} 000 \overline{0} h$
Table 99 - Port B Cell Departure Counter High Register

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| ob_dep[31:16] | $15: 0$ | CNT | Freerunning counter of the number of cells transmitted on <br> port B. |

Table 99 - Port B Cell Departure Counter High Register

Address: 356h
Label: portb_cell_departure_low
Reset Value: $000 \overline{0} \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| ob_dep[15:0] | $15: 0$ | CNT | Freerunning counter of the number of cells transmitted on <br> port B. |

Table 100 - Port B Cell Departure Counter Low Register

Address: 358h
Label: portb overflow0
Reset Value: $\mathbf{0} 000 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| ib_rx_cell_max | $4: 0$ | RW | If the cell fill of the RX SAR output FIFO becomes greater <br> than this value, cells from the port B input FIFO will be <br> blocked. Oh = no backpressure |
| ib_oa_cell_max | $9: 5$ | RW | If the cell fill of the Port A output FIFO becomes greater <br> than this value, cells from the port B input FIFO will be <br> blocked. Oh = no backpressure |
| ib_ob_cell_max | $14: 10$ | RW | If the cell fill of the Port B output FIFO becomes greater <br> than this value, cells from the port B input FIFO will be <br> blocked. Oh = no backpressure |
| reserved | 15 | RW | Reserved. Must always be "0" |

Table 101-Port B Overflow0 Register

Address: 35Ah
Label: portb_overflow1
Reset Value: $\mathbf{0 0 0 0} \mathbf{h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| ib_oc_cell_max | $4: 0$ | RW | If the cell fill of the Port C output FIFO becomes greater <br> than this value, cells from the port B input FIFO will be <br> blocked. Oh = no backpressure |
| reserved | $15: 5$ | RW | Reserved. Must always be "0000_0000_000" |

Table 102 - Port B Overflow1 Register

## Address: 360h

Label: portc_look_up_base
Reset Value: $\mathbf{0} 000 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| lutc_base | $15: 0$ | RW | Bits 19:4 of the address of the look-up table for port C. |

Table 103 - Port C Look Up Table Register

Address: 362h
Label: portc_num_vpi_vci_bits
Reset Value: $\mathbf{0 0 0 0} \mathbf{h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| num_vpi_vci_bits | $15: 0$ | RW | Indicates the total number of bits that are used to decode the look- <br> up address for port C. 0001h $=1$ 1 bit used, FFFFh $=16$ bits used. <br> When subtracted from vci_n, this indicates the number of VPI bits <br> used. |

Table 104 - Port C VPI/VCI Identification Register

Address: 364h
Label: portc_concatenation
Reset Value: $\mathbf{0 0 0 0} \mathrm{h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vci_nc | $4: 0$ | RW | Indicates the number of VCI bits used to decode the look- <br> up address for port C. |
| reserved | $15: 5$ | RW | Reserved. Must always be "0000_0000_000" |

Table 105 - Port C Concatenation Register

## Address: 368h <br> Label: portc_vpi_match <br> Reset Value: $\mathbf{0} 00 \overline{0} \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vpi_matchc | $11: 0$ | RW | For a cell from port C to be considered valid, any bits in its <br> VPI whose corresponding bits in reg 34Ah are '1' must have <br> the value contained in this register. |
| reserved | $15: 12$ | RW | Reserved. Must always be "0000" |

Table 106 - Port C VPI Match Register

Address: 36Ah Label: portc_vpi_mask
Reset Value: $\mathbf{0} 00 \overline{0} \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vpi_maskc | $11: 0$ | RW | For a cell from port C to be considered valid, any bits in <br> its VPI whose corresponding bits in this register are '1' <br> must have the value contained in reg 348h. |
| reserved | $15: 12$ | RW | Reserved. Must always be "0000" |

Table 107 - Port C VPI Mask Register

Address: 36Ch
Label: portc_vci_match
Reset Value: $\mathbf{0} 000 \overline{\mathrm{O}}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vci_matchc | $15: 0$ | RW | For a cell from port C to be considered valid, any bits in its VCI whose <br> corresponding bits in reg 34Eh are '1' must have the value contained <br> in this register. |

Table 108 - Port C VCI Match Register

Address: 36Eh
Label: portc_vci_mask
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| vci_maskc | $15: 0$ | RW | For a cell from port C to be considered valid, any bits in its VCI whose <br> corresponding bits in this register are '1' must have the value <br> contained in reg 34Ch. |

Table 109 - Port C VCI Match Register
Address: 370h
Label: portc_cell_arrival_high
Reset Value: $\mathbf{0 0 0 0} \mathbf{h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| ic_arr[31:16] | $15: 0$ | CNT | Freerunning counter of the number of cells received on port C. |

Table 110 - Port C Cell Arrival Counter High Register

```
Address: 372h
Label: portc_cell_arrival_low
Reset Value: 0000h
```

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| ic_arr[15:0] | $15: 0$ | CNT | Freerunning counter of the number of cells received on port <br> C. |

Table 111 - Port C Cell Arrival Counter Low Register

Address: 374h
Label: portc_cell_departure_high
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| oc_dep[31:16] | $15: 0$ | CNT | Freerunning counter of the number of cells transmitted <br> on port C. |

Table 112 - Port C Cell Departure Counter High Register

| Address: 376h <br> Label: portc_cell_departure_low <br> Reset Value: $\mathbf{0 0 0 0} \mathbf{h}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| oc_dep[15:0] | $15: 0$ | CNT | Freerunning counter of the number of cells transmitted on <br> port C. |

Table 113 - Port C Cell Departure Counter Low Register

Address: 378h
Label: portc_overflow0
Reset Value: $\mathbf{0 0 0 0} \mathrm{h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| ic_rx_cell_max | $4: 0$ | RW | If the cell fill of the RX SAR output FIFO becomes greater <br> than this value, cells from the port C input FIFO will be <br> blocked. Oh = no backpressure |
| ic_oa_cell_max | $9: 5$ | RW | If the cell fill of the Port A output FIFO becomes greater than <br> this value, cells from the port C input FIFO will be blocked. <br> Oh = no backpressure |
| ic_ob_cell_max | $14: 10$ | RW | If the cell fill of the Port B output FIFO becomes greater than <br> this value, cells from the port C input FIFO will be blocked. <br> Oh = no backpressure |

Table 114 - Port C Overflow0 Register

Address: 378h
Label: portc_overflow0
Reset Value: 0000 h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| reserved | 15 | RW | Reserved. Must always be "0" |

Table 114 - Port C Overflow0 Register (continued)

| Address: 37Ah <br> Label: portc_overflow1 <br> Reset Value: $\mathbf{0 0 0 0 0}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type |  |
| ic_oc_cell_max | $4: 0$ | RW | If the cell fill of the Port C output FIFO becomes greater <br> than this value, cells from the port C input FIFO will be <br> blocked. Oh = no backpressure |
| reserved | $15: 5$ | RW | Reserved. Must always be "0000_0000_000" |

Table 115 - Port C Overflow1 Register

Address: 390h
Label: tx_sar_cell_arrival_high
Reset Value: $\overline{0} 000 \bar{h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tx_arr[31:16] | $15: 0$ | CNT | Bits [31:16] of TX SAR arrival cell counter. Counts the number of <br> cells received by UTOPIA from the TX SAR |

Table 116-TX_SAR Cell Arrival Counter High Register

Address: 392h
Label: tx_sar_cell_arrival_low
Reset Value: 0000h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tx_arr[15:0] | $15: 0$ | CNT | Bits [15:0] of TX SAR arrival cell counter. Counts the <br> number of cells received by UTOPIA from the TX SAR |

Table 117 - TX_SAR Cell Arrival Counter Low Register

```
Address: 394h
Label: rx_sar_cell_departure_high
Reset Value: 0000h
```

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| rx_dep[31:16] | $15: 0$ | CNT | Bits [31:16] of RX SAR departure cell counter. Counts the <br> number of cells sent to the RX SAR from UTOPIA |

Table 118 - RX_SAR Cell Departure Counter High Register

| Address: 396h <br> Label: rx_sar_cell_departure_low <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| rx_dep[15:0] | $15: 0$ | CNT | Bits [15:0] of RX SAR departure cell counter. Counts the <br> number of cells sent to the RX SAR from UTOPIA |

Table 119 - RX_SAR Cell Departure Counter Low Register

Address: 398h Label: tx_sar_overflow Reset Value: $\overline{0000 h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tx_rx_cell_max | $4: 0$ | RW | lf the cell fill of the RX SAR output FIFO becomes greater than this <br> value, cells from the TX SAR input FIFO will be blocked. Oh $=$ no <br> backpressure |
| tx_oa_cell_max | $9: 5$ | RW | lf the cell fill of the Port A output FIFO becomes greater than this value, <br> cells from the TX SAR input FIFO will be blocked. Oh = no backpressure |
| tx_ob_cell_max | $14: 10$ | RW | lf the cell fill of the Port B output FIFO becomes greater than this value, <br> cells from the TX SAR input FIFO will be blocked. Oh $=$ no backpressure |
| reserved | 15 | RW | Reserved. Must always be "0" |

Table 120-TX_SAR Overflow0 Register

Address: 39Ah
Label: tx_sar_overflow
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tx_oc_cell_max | $4: 0$ | RW | If the cell fill of the Port C output FIFO becomes greater than <br> this value, cells from the TX SAR input FIFO will be blocked. Oh <br> = no backpressure |
| reserved | $15: 5$ | RW | Reserved. Must always be "0000_0000_000" |

Table 121 - TX_SAR Overflow1 Register

| Address: 3A0h <br> Label: hec_byte_control <br> Reset Value: 0055 h |
| :--- |
| Label |
| Bit |
| Position | Type $\quad$| Description |
| :--- |
| hec_mask |
| $7: 0$ |
| reserved |
| RW |

Table 122 - HEC Byte Control Register

```
Address: 3A2h
Label: unknown_header_routing
Reset Value: 0000}
```

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| rxa_ncr | $3: 0$ | RW | Normal cell routing for unknow cells received on port A. "xxx1" = Port A, <br> "xx1x" = Port B, "x1xx" = Port C, "1xxx" RX SAR. Cells can be broadcast to <br> multiple destinations by setting multiple bits in this field to '1'. |
| rxb_ncr | $7: 4$ | RW | Normal cell routing for unknow cells received on port B. "xxx1" = Port A, <br> "xx1x" = Port B, "x1xx" = Port C, "1xxx" RX SAR. Cells can be broadcast to <br> multiple destinations by setting multiple bits in this field to '1'. |
| rxc_ncr | $11: 8$ | RW | Normal cell routing for unknow cells received on port C. "xxx1" = Port A, <br> "xx1x" = Port B, "x1xx" = Port C, "1xxx" RX SAR. Cells can be broadcast to <br> multiple destinations by setting multiple bits in this field to '1'. |
| reserved | $15: 12$ | RW | Reserved. Must always be "0000" |

Table 123 - Unknown Header Routing Register

| Address: 3A4h <br> Label: unknown_oam_routing <br> Reset Value: 0000̄h |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | Bit Position | Type | Description |
| rxa_ocr | 3:0 | RW | OAM cell routing for unknow cells received on port A . "xxx1" = Port A, "xx1x" = Port B, "x1xx" = Port C, "1xxx" RX SAR. Cells can be broadcast to multiple destinations by setting multiple bits in this field to ' 1 '. |
| rxb_ocr | 7:4 | RW | OAM cell routing for unknow cells received on port B. "xxx1" = Port A, "xx1x" = Port B, "x1xx" = Port C, "1xxx" RX SAR. Cells can be broadcast to multiple destinations by setting multiple bits in this field to ' 1 '. |
| rxc_ocr | 11:8 | RW | OAM cell routing for unknow cells received on port $C$. "xxx1" = Port $A$, "xx1x" = Port B, "x1xx" = Port C, "1xxx" RX SAR. Cells can be broadcast to multiple destinations by setting multiple bits in this field to ' 1 '. |
| reserved | 15:12 | RW | Reserved. Must always be "0000" |

Table 124 - Unknown OAM Routing Register

| Address: 3C0h <br> Label: gpio_input0 <br> Reset Value: 0000 h |
| :--- |
| Label |
| Bit Position |
| txa_data_input |
| rxa_data_input |
| $15: 0$ |
| RO |

Table 125 - GPIO Input0 Register

Address: 3C2h Label: gpio_input1
Reset Value: 0000 h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| txb_data_input | $7: 0$ | RO | Current level of txb_data pins [15:8] |
| rxb_data_input | $15: 8$ | RO | Current level of rxb_data pins [15:8] |

Table 126 - GPIO Input1 Register

Address: 3C4h
Label: gpio_input2
Reset Value: 0000 h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| phya_alm_input | 0 | RO | Current level of phya_alm pin |
| phyb_alm_input | 1 | RO | Current level of phya_alm pin |
| phya_tx_led_input | 2 | RO | Current level of phya_tx_led pin |
| phya_rx_led_input | 3 | RO | Current level of phya_rx_led pin |
| phyb_tx_led_input | 4 | RO | Current level of phyb_tx_led pin |
| phyb_rx_led_input | 5 | RO | Current level of phyb_rx_led pin |
| reserved | $15: 6$ | RO | Reserved. Always read as "0000_0000_00" |

Table 127-GPIO Input2 Register

## Address: 3C8h <br> Label: txa_data status <br> Reset Value: 0000h

| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| :---: | :---: | :---: | :---: |
| txa_data8_rise | 0 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| txa_data8_fall | 1 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |
| txa_data9_rise | 2 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| txa_data9_fall | 3 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |
| txa_data10_rise | 4 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| txa_data10_fall | 5 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |
| txa_data11_rise | 6 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| txa_data11_fall | 7 | ROL | This bit is set when corresponding pin changes from ' 1 ' to ' 0 ' |
| txa_data12_rise | 8 | ROL | This bit is set when corresponding pin changes from ' 0 ' to ' 1 ' |
| txa_data12_fall | 9 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |
| txa_data13_rise | 10 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| txa_data13_fall | 11 | ROL | This bit is set when corresponding pin changes from ' 1 ' to ' 0 ' |
| txa_data14_rise | 12 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| txa_data14_fall | 13 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |
| txa_data15_rise | 14 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| txa_data15_fall | 15 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |

Table 128-TXA Data Status Register

| Address: 3CAh <br> Label: txa_data_status_ie <br> Reset Value: $\mathbf{0 0 0 0 0}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type |  |
| txa_data8_rise_ie | 0 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data8_fall_ie | 1 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data9_rise_ie | 2 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data9_fall_ie | 3 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data10_rise_ie | 4 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |

Table 129 - TXA Data Interrupt Enable Register

Address: 3CAh
Label: txa_data_status_ie
Reset Value: 0000 h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| txa_data10_fall_ie | 5 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data11_rise_ie | 6 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data11_fall_ie | 7 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data12_rise_ie | 8 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data12_fall_ie | 9 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data13_rise_ie | 10 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data13_fall_ie | 11 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data14_rise_ie | 12 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data14_fall_ie | 13 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txa_data15_rise_ie | 14 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| txata15_fall_ie | 15 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |

Table 129 - TXA Data Interrupt Enable Register (continued)

## Address: 3CCh <br> Label: rxa_data_status <br> Reset Value: 0000h

| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| :---: | :---: | :---: | :---: |
| rxa_data8_rise | 0 | ROL | This bit is set when corresponding pin changes from '0' to '11' |
| rxa_data8_fall | 1 | ROL | This bit is set when corresponding pin changes from ' 1 ' to ' 0 ' |
| rxa_data9_rise | 2 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| rxa_data9_fall | 3 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |
| rxa_data10_rise | 4 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| rxa_data10_fall | 5 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |
| rxa_data11_rise | 6 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| rxa_data11_fall | 7 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |
| rxa_data12_rise | 8 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| rxa_data12_fall | 9 | ROL | This bit is set when corresponding pin changes from ' 1 ' to ' 0 ' |
| rxa_data13_rise | 10 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| rxa_data13_fall | 11 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |
| rxa_data14_rise | 12 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| rxa_data14_fall | 13 | ROL | This bit is set when corresponding pin changes from ' 1 ' to ' 0 ' |
| rxa_data15_rise | 14 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| rxa_data15_fall | 15 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |

Table 130-RXA Data Status Register

Address: 3CEh
Label: rxa_data_status_ie
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| rxa_data8_rise_ie | 0 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data8_fall_ie | 1 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data9_rise_ie | 2 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data9_fall_ie | 3 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data10_rise_ie | 4 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |

Table 131-RXA Data Interrupt Enable Register

Address: 3CEh
Label: rxa_data_status_ie
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| rxa_data10_fall_ie | 5 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data11_rise_ie | 6 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data11_fall_ie | 7 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data12_rise_ie | 8 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data12_fall_ie | 9 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data13_rise_ie | 10 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data13_fall_ie | 11 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data14_rise_ie | 12 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data14_fall_ie | 13 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data15_rise_ie | 14 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rxa_data15_fall_ie | 15 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |

Table 131-RXA Data Interrupt Enable Register (continued)

Address: 3DOh
Label: txb_data_status
Reset Valué: 0000h

| Label | Bit <br> Position | Type |  |
| :---: | :---: | :---: | :--- |
| txb_data8_rise | 0 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| txb_data8_fall | 1 | ROL | This bit is set when corresponding pin changes from '1' to '0' |
| txb_data9_rise | 2 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| txb_data9_fall | 3 | ROL | This bit is set when corresponding pin changes from '1' to '0' |
| txb_data10_rise | 4 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| txb_data10_fall | 5 | ROL | This bit is set when corresponding pin changes from '1' to '0' |
| txb_data11_rise | 6 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| txb_data11_fall | 7 | ROL | This bit is set when corresponding pin changes from '1' to '0' |

Table 132- TXB Data Status Register

## Address: 3DOh

Label: txb_data_status
Reset Valué: 0000h

| Label | Bit <br> Position | Type |  |
| :--- | :---: | :---: | :--- |
| txb_data12_rise | 8 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| txb_data12_fall | 9 | ROL | This bit is set when corresponding pin changes from '1' to '0' |
| txb_data13_rise | 10 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| txb_data13_fall | 11 | ROL | This bit is set when corresponding pin changes from '1' to '0' |
| txb_data14_rise | 12 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| txb_data14_fall | 13 | ROL | This bit is set when corresponding pin changes from '1' to '0' |
| txb_data15_rise | 14 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| txb_data15_fall | 15 | ROL | This bit is set when corresponding pin changes from '1' to '0' |

Table 132-TXB Data Status Register (continued)

| Address: 3D2h <br> Label: txb_data_status_ie <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{aligned} & \text { Bit } \\ & \text { Position } \end{aligned}$ | Type | Description |
| txb_data8_rise_ie | 0 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| txb_data8_fall_ie | 1 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| txb_data9_rise_ie | 2 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| txb_data9_fall_ie | 3 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| txb_data10_rise_ie | 4 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| txb_data10_fall_ie | 5 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| txb_data11_rise_ie | 6 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| txb_data11_fall_ie | 7 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| txb_data12_rise_ie | 8 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| txb_data12_fall_ie | 9 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |

Table 133- TXB Data Interrupt Enable Register

Address: 3D2h
Label: txb_data_status_ie
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :--- | :---: | :---: | :--- |
| txb_data13_rise_ie | 10 | IE | When '1' and the corresponding status bit is '1', an interrupt will be <br> generated. |
| txb_data13_fall_ie | 11 | IE | When '1' and the corresponding status bit is '1', an interrupt will be <br> generated. |
| txb_data14_rise_ie | 12 | IE | When '1' and the corresponding status bit is '1', an interrupt will be <br> generated. |
| txb_data14_fall_ie | 13 | IE | When '1' and the corresponding status bit is '1', an interrupt will be <br> generated. |
| txb_data15_rise_ie | 14 | IE | When '1' and the corresponding status bit is '1', an interrupt will be <br> generated. |
| txb_data15_fall_ie | 15 | IE | When '1' and the corresponding status bit is '1', an interrupt will be <br> generated. |

Table 133-TXB Data Interrupt Enable Register (continued)

## Address: 3D4h Label: rxb_data_status Reset Value: 0000h

| Label | Bit <br> Position | Type |  |
| :---: | :---: | :---: | :--- |
| rxb_data8_rise | 0 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| rxb_data8_fall | 1 | ROL | This bit is set when corresponding pin changes from '1' to '0' |
| rxb_data9_rise | 2 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| rxb_data9_fall | 3 | ROL | This bit is set when corresponding pin changes from '1' to '0' |
| rxb_data10_rise | 4 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| rxb_data10_fall | 5 | ROL | This bit is set when corresponding pin changes from '1' to '0' |
| rxb_data11_rise | 6 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| rxb_data11_fall | 7 | ROL | This bit is set when corresponding pin changes from '1' to '0' |
| rxb_data12_rise | 8 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| rxb_data12_fall | 9 | ROL | This bit is set when corresponding pin changes from '1' to '0' |
| rxb_data13_rise | 10 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| rxb_data13_fall | 11 | ROL | This bit is set when corresponding pin changes from '1' to '0' |
| rxb_data14_rise | 12 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| rxb_data14_fall | 13 | ROL | This bit is set when corresponding pin changes from '1' to '0' |

Table 134-RXB Data Status Register

## Address: 3D4h <br> Label: rxb_data_status <br> Reset Valuè: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| rxb_data15_rise | 14 | ROL | This bit is set when corresponding pin changes from '0' to '1' |
| rxb_data15_fall | 15 | ROL | This bit is set when corresponding pin changes from '1' to '0' |

Table 134-RXB Data Status Register (continued)

| Address: 3D6h <br> Label: rxb_data_status_ie <br> Reset Value: 0000 |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| rxb_data8_rise_ie | 0 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data8_fall_ie | 1 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data9_rise_ie | 2 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data9_fall_ie | 3 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data10_rise_ie | 4 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data10_fall_ie | 5 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data11_rise_ie | 6 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data11_fall_ie | 7 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data12_rise_I | 8 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data12_fall_ie | 9 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data13_rise_ie | 10 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data13_fall_ie | 11 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data14_rise_ie | 12 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data14_fall_ie | 13 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data15_rise_ie | 14 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |
| rxb_data15_fall_ie | 15 | IE | When ' 1 ' and the corresponding status bit is ' 1 ', an interrupt will be generated. |

Table 135-RXB Data Interrupt Enable Register

| Address: 3D8h Label: gpio_status Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| phya_alm_rise | 0 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| phya_alm_fall | 1 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |
| phyb_alm_rise | 2 | ROL | This bit is set when corresponding pin changes from '0' to '11' |
| phyb_alm_fall | 3 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |
| phya_tx_led_rise | 4 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| phya_tx_led_fall | 5 | ROL | This bit is set when corresponding pin changes from ' 1 ' to ' 0 ' |
| phya_rx_led_rise | 6 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| phya_rx_led_fall | 7 | ROL | This bit is set when corresponding pin changes from ' 1 ' to ' 0 ' |
| phyb_tx_led_rise | 8 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| phyb_tx_led_fall | 9 | ROL | This bit is set when corresponding pin changes from ' 1 ' to ' 0 ' |
| phyb_rx_led_rise | 10 | ROL | This bit is set when corresponding pin changes from '0' to ' 1 ' |
| phyb_rx_led_fall | 11 | ROL | This bit is set when corresponding pin changes from ' 1 ' to '0' |
| reserved | 15:12 | ROL | Reserved. Always read as "0000" |

Table 136-GPIO Status Register

Address: 3DAh
Label: gpio_status_ie
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| phya_alm_rise_ie | 0 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| phya_alm_fall_ie | 1 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| phyb_alm_rise_ie | 2 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| phyb_alm_fall_ie | 3 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| phya_tx_led_rise_ie | 4 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| phya_tx_led_fall_ie | 5 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| phya_rx_led_rise_ie | 6 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |

Table 137-GPIO Status Register

| Address: 3DAh <br> Label: gpio_status_ie <br> Reset Value: $\mathbf{0 0 0 0 h}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type |  |
| phya_rx_led_fall_ie | 7 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| phyb_tx_led_rise_ie | 8 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| phyb_tx_led_fall_ie | 9 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| phyb_rx_led_rise_ie | 10 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| phyb_rx_led_fall_ie | 11 | IE | When '1' and the corresponding status bit is '1', an interrupt <br> will be generated. |
| reserved | $15: 12$ | RO | Reserved. Always read as "0000" |

Table 137-GPIO Status Register (continued)

| Address: 3EOh <br> Label: gpio_output0 <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| txa_data_output | $7: 0$ | RW | This is the value sent out on the pins txa_data [15:8], used in <br> conjuction with the OE bit |
| reserved | $15: 8$ | RW | Reserved. Must always be "0000_0000" |

Table 138-GPIO Output0 Register

Address: 3E2h
Label: gpio_output1
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| txb_data_output | $7: 0$ | RW | This is the value sent out on the pins txb_data [15:8], used in <br> conjuction with the OE bit |
| reserved | $15: 8$ | RW | Reserved. Must always be "0000_0000" |

Table 139-GPIO Output1 Register

| Address: 3E4h <br> Label: gpio_output2 <br> Reset Value: 0000 h |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{aligned} & \text { Bit } \\ & \text { Position } \end{aligned}$ | Type | Description |
| phya_tx_led_output | 0 | RW | This is the value sent out on the corresponding pin, used in conjuction with the OE bit |
| phya_rx_led_output | 1 | RW | This is the value sent out on the corresponding pin, used in conjuction with the OE bit |
| phyb_tx_led_output | 2 | RW | This is the value sent out on the corresponding pin, used in conjuction with the OE bit |
| phyb_rx_led_output | 3 | RW | This is the value sent out on the corresponding pin, used in conjuction with the OE bit |
| reserved | 15:4 | RW | Reserved. Must always be "0000_0000_0000" |

Table 140-GPIO Output2 Register

| Address: 3E8h <br> Label: gpio_oe0 <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| txa_data_oe | $7: 0$ | RW | This is OE bit used to drive the txa_data [15:8] pins |
| reserved | $15: 8$ | RW | Reserved. Must always be "0000_0000" |

Table 141 - GPIO Output Enable0 Register

Address: 3EAh
Label: gpio_oe1
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| txb_data_oe | $7: 0$ | RW | This is OE bit used to drive the txb_data [15:8] pins |
| reserved | $15: 8$ | RW | Reserved. Must always be "0000_0000" |

Table 142-GPIO Output Enable1 Register

| Address: 3ECh <br> Label: gpio_oe2 <br> Reset Value: 0000 h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type |  |
| phya_tx_led_oe | 0 | RW | This is OE bit used to drive the corresponding pin |
| phya_rx_led_oe | 1 | RW | This is OE bit used to drive the corresponding pin |
| phyb_tx_led_oe | 2 | RW | This is OE bit used to drive the corresponding pin |
| phyb_rx_led_oe | 3 | RW | This is OE bit used to drive the corresponding pin |
| reserved | $15: 4$ | RW | Reserved. Must always be "0000_0000_0000" |

Table 143-GPIO Output Enable2 Register

### 5.5.4 TDM Registers

Address: 400h
Label: control
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| global_oe | 0 | RW | '0' = ct_d[31:0] forced tri-state; '1' = ct_d[31:0] may be driven. |
| h100_data_loopback | 1 | RW | '0' = no ct_d[31:0] loopback; '1' $=$ ct_d[31:0] loopback. For <br> test only. |
| tdmie_enable | 2 | RW | '0' disables TDM process. Should only be set to '1' when <br> CAM has been programmed |
| stream_mode | $4: 3$ | RW | "00" $=32$ streams, "01" $=16$ streams, "10" $=4$ streams, "11" <br> $=$ reserved. This register is only used for tests: in real <br> operation, it should be left to "00" (32 streams). |
| reserved | $14: 5$ | RW | Reserved. Must always be "0000_0000_00" |
| test_status | 15 | TS | When '1', all the status bits in the register will be set. |

Table 144 - TDM Control Register

Address: 402h
Label: status0
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tdm_out_of_bandwidth0 | 0 | ROL | Indicates that the TDM state machine is out of bandwidth. Fatal <br> chip error, usually due to mclk frequency being too low. |
| cut_vc_detected | 1 | ROL | This bit is set when 255 underruns are detected on a VC whose <br> CAM entry has an enabled "Cut VC status enable" bit |

Table 145 - TDM Status Register

Address: 402h
Label: status0
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| underrun_detected | 2 | ROL | Underrun reported on a voice channel on the ATM link. |
| cas_underrun_detected | 3 | ROL | Underrun reported on CAS bits on the ATM link. |
| rdatamem_overflow | 4 | ROL | Overflow in the TDM RX internal data memory. Fatal chip error. |
| tdmtxpip_overflow | 5 | ROL | Overflow in the TDM TX data memory access cache. Fatal chip <br> error. |
| tdmrxpip_overflow | 6 | ROL | Overflow in the TDM RX data memory access cache. Fatal chip <br> error. |
| reserved | $15: 7$ | ROL | Reserved. Always read as "0000_0000_0" |

Table 145- TDM Status Register (continued)

Address: 404h
Label: status_ie
Reset Value: $\overline{0000 h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tdm_out_of_bandwidth0_ie | 0 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| cut_vc_detected_ie | 1 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| underrun_detected_ie | 2 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| cas_underrun_detected_ie | 3 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| rdatamem_overflow_ie | 4 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| tdmtxpip_overflow_ie | 5 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| tdmrxpip_overflow_ie | 6 | IE | When '1' and the corresponding status bit is '1', an <br> interrupt will be generated. |
| reserved | $15: 7$ | IE | Reserved. Always read as "0000_0000_0" |

Table 146 - TDM Interrupt Enable Register

## Address: 408h

Label: cut_vc_tsst
Reset Value: $\mathbf{0} \mathbf{0} 00 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tsst_number_cut_vc | $11: 0$ | RO | Indicates the TSST on which the last cut VC error occurred. |
| reserved | $15: 12$ | RO | Reserved. Always read as "0000" |

Table 147 - Cut VC TSST Register

Address: 40Ah
Label: underrun_tsst
Reset Value: 0000̄h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tsst_number_underrun | $11: 0$ | RO | Indicates the TSST on which the last underrun error <br> occurred. |
| reserved | $15: 12$ | RO | Reserved. Always read as "0000" |

Table 148 - TSST Underrun Register

Address: 40Ch
Label: cas_underrun_tsst
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tsst_number_cas_underrun | $11: 0$ | RO | Indicates the TSST on which the last CAS underrun error <br> occurred. |
| reserved | $15: 12$ | RO | Reserved. Always read as "0000" |

Table 149 - TSST CAS Underrun Register

Address: 410h
Label: tdmint_reg0
Reset Value: $\overline{0000} \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| h100_samp_clk_delay_flops | $3: 0$ | RW | Number of flops used = 8 + value of this register. "1111" <br> is reserved for selecting falling edge ct_c8 clock. |
| h100_samp_clk_delay_buff | $7: 4$ | RW | "1111" is reserved for selecting rising edge ct_c8 clock |
| h100_oe_clk_delay_flops | $11: 8$ | RW | Number of flops used = 8 + value of this register |
| h100_oe_clk_delay_buff | $15: 12$ | RW | Number of delay buffers used in the delay chain. |

Table 150 - TDM Interrupt 0 Register

Address: 412h Label: tdmint_reg1 Reset Value: $\mathbf{0 0 0 0}$

| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| :---: | :---: | :---: | :---: |
| h100_ts_counter_timeout | 6:0 | RW | Number of mclk cycles in 8 ct_c8 clock cycles - 10\% |
| h100_oe_override_disable | 7 | RW | '1' = do not tri-state the ct_d pin after every timeslot. To respect the H .100 standard, this should be left to ' 0 '. |
| dstream_0_3_freq | 9:8 | RW | ct_d[3:0] stream clock speed. "00" = 2.048 Mhz; "01" = $4.096 \mathrm{MHz} ; " 10 "=8.192 \mathrm{MHz} ; " 11 "=$ reserved. |
| dstream_4_7_freq | 11:10 | RW | ct_d[7:4] stream clock speed. "00" = 2.048 Mhz; "01" = 4. $096 \mathrm{MHz} ; " 10 "=8.192 \mathrm{MHz} ; " 11 "=$ reserved. |
| dstream_8_11_freq | 13:12 | RW | ct_d[11:8] stream clock speed. "00" $=2.048 \mathrm{Mhz} ;$ "01" = 4.096 MHz; "10" = $8.192 \mathrm{MHz} ; " 11 "=$ reserved. |
| dstream_12_15_freq | 15:14 | RW | ct_d[15:12] stream clock speed. "00" = 2.048 Mhz; "01" = 4.096 MHz; "10" = 8.192 MHz; "11" = reserved. |

Table 151 - TDM Interrupt 1 Register

Address: 420h Label: tdmie misc Reset Value: $\mathbf{0} 0$ FFh

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| null_byte | $7: 0$ | RW | Null byte with which the chip will pad if underruns occur and null <br> byte padding is chosen. |
| cas_enable_position | $10: 8$ | RW | Position of the CAS enable bit within the byte on the TDM bus. <br> The CAS nibble will be contained in the nibble in which the <br> enable is not present. For example, if the enable is contained in <br> bit 6, then the nibble will be in bits 3:0. |
| cas_enable_polarity | 11 | RW | Polarity of the CAS enable bit on the TDM bus. |
| reserved | $15: 12$ | RW | Reserved. Must always be "0000" |

Table 152 - TDM Interrupt Enable Misc. Register

| Address: 460h <br> Label: pnt_ro <br> Reset Value: 0000 h |  |  |  |
| :--- | :---: | :---: | :--- |
| Label | Bit <br> Positio <br> $\mathbf{n}$ | Type |  |
| txsar_write_pnt_pcm_monitor <br> [9:0] | $9: 0$ | RO | The current value of the TDM write pointer sent to the TX <br> SAR. Only used for tests. |
| reserved | $15: 10$ | RO | Reserved. Always read as "0000_00" |

Table 153 - TDM Write Pointer 0 Register

| Address: 462h <br> Label: txar_pnt_fire <br> Reset Value:_0000h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Positio <br> n | Type |  |
| txsar_write_pnt_insert[9:0] | $9: 0$ | RW | The test value of the TDM write pointer sent to the TX <br> SAR. Only used for tests. |
| txsar_write_pnt_insert_ena | 10 | RW | When '1', the above pointer will be sent to the TX SAR <br> instead of the valid pointer. Only used for tests. |
| reserved | $15: 11$ | RO | Reserved. Always read as "0000_00" |

Table 154 - TDM Write Pointer 1 Register

| Address: 464h <br> Label: rxsar_pnt_fire <br> Reset Value: $000 \overline{0} h$ |  |  |  |
| :--- | :--- | :--- | :--- |
| Label | Bit <br> Position | Type |  |
| rxsar_write_pnt_insert[14:0] | $14: 0$ | RW | The test value of the TDM read pointer sent to the RX <br> SAR. Only used for tests. |
| rxsar_write_pnt_insert_ena | 15 | RW | When '1', the above pointer will be sent to the RX SAR <br> instead of the valid pointer. Only used for tests. |

Table 155 - TDM Read Pointer Register

### 5.5.5 TX_SAR Registers

Address: 500h
Label: control
Reset Value: 0000h

| Label | Bit <br> Position | Type |  |
| :---: | :---: | :---: | :--- |
| reserved | $14: 0$ | RO | Rescription |
| Rest_status | 15 | TS | When '1', all the status bits in the register will be set. |

Table 156-TX_SAR Control Register

Address: 502h
Label: status
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| global_tx_slip | 0 | ROL | Raised when a bad configuration of the offset field in the TX <br> control structure causes the TX SAR to read data that has <br> not been written yet. |
| txsarpip_overflow | 1 | ROL | Overflow in the TX SAR data memory access cache. Fatal <br> chip error. |
| reserved | $15: 2$ | RO | Reserved. Always read as "0000_0000_0000_00" |

Table 157 - TX_SAR Status Register

| Address: 504h <br> Label: status_ie <br> Reset Value: $\overline{0000 h}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| global_tx_slip_ie | 0 | IE | When '1' and the corresponding status bit is '1', an interrupt will be <br> generated. |
| txsarpip_overflow_ie | 1 | IE | When '1' and the corresponding status bit is '1', an interrupt will be <br> generated. |
| reserved | $15: 2$ | IE | Reserved. Always read as ""0000_0000_0000_00" |

Table 158 - TX_SAR Interrupt Enable Register

Address: 506h
Label: control 1
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| reset_band_per | 0 | PUL | Resets the band_per register field. |
| reserved | $15: 1$ | RO | Reserved. Always read as "0000_0000_0000_000" |

Table 159-TX_SAR Control 1 Register

Address: 508h
Label: data_cell_read
Reset Value: $000 \overline{0} \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| data_read_pnt | $13: 0$ | RO | Chip's read pointer to the AAL0 FIFO. |
| reserved | $15: 14$ | RO | Reserved. Always read as "00" |

Table 160 - TX_SAR Data Read Pointer Register

## Address: 50Ah <br> Label: data_cell_write <br> Reset Value: 0000 h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| data_write_pnt | $13: 0$ | RW | CPU's write pointer to the AAL0 FIFO. |
| reserved | $15: 14$ | RW | Reserved. Must always be "00" |

Table 161 - TX_SAR Data Write Pointer Register

| Address: 50Ch <br> Label: data_celladd <br> Reset Value: $\mathbf{0 0 0 0} \mathrm{h}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| data_add | $13: 0$ | RW | Bits 19:6 of the address of the FIFO in external memory |
| reserved | $15: 14$ | RW | Reserved. Always read as "00" |

Table 162-TX_SAR Data Address Register

Address: 50Eh Label: data_cell_size
Reset Value: 0000 h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| data_size | $13: 0$ | RW | Size of data cell FIFO in 1 cell increments. All zeros $=16 \mathrm{k}$ <br> cells. Minimum 4 cells. |
| reserved | $15: 14$ | RW | Reserved. Must always be "00" |

Table 163 - TX_SAR Data Cell Size Register

Address: 510h
Label: percentage_of_bandwidth
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| band_per | $15: 0$ | RO | Monitor of the maximum number of mclk cycles that it has <br> taken the chip to treat an entire TX SAR frame. If this number <br> is greater than mclk (in Hz) / 8000, some of the frames are <br> overloaded. |

Table 164-Percent of Bandwidth Register

### 5.5.6 Scheduler Registers

Address: 600h
Label: control
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| test_status | 15 | TS | When '1', all the status bits in the register will be set. |

Table 165 - Scheduler Test Status Register

Address: 602h
Label: status
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| out_of_band | 0 | ROL | Indicates that the wheel treatment is more than fr_late frames <br> late. This means that at least some of the frames in the wheels <br> are overloaded. |
| reserved | $15: 1$ | ROL | Reserved. Always read as "0000_0000_0000_000" |

Table 166 - Scheduler Status Register

Address: 604h
Label: status_ie
Reset Value: $\overline{\mathbf{0 0} 00} \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| out_of_band_ie | 0 | IE | When '1' and the corresponding status bit is '1', an interrupt will be <br> generated. |
| reserved | $15: 1$ | IE | Reserved. Always read as "0000_0000_0000_000" |

Table 167 - Scheduler Interrupt Enable Register

Address: 608h
Label: frame_latency
Reset Value: $\mathbf{0} 001 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| fr_late | $3: 0$ | RW | Number of frames by which the chip is allowed to be late. <br> Oh = illegal. |
| reserved | $15: 4$ | RW | Reserved. Must always be "0000_0000_0000" |

Table 168 - Frame Latency Register

Address: 610h
Label: wheel_info_0
Reset Value: $\overline{0} 001 \mathbf{h}^{-}$

| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| :---: | :---: | :---: | :---: |
| wheel0_ena | 0 | RW | Enable for wheel 0. |
| wheel0_inf | 3:1 | RW | Configuration of wheel $0 . " 000 "=$ normal, " 100 " = T1, " $101 "=E 1$, others reserved.I |
| wheel1_ena | 4 | RW | Enable for wheel 1. |
| wheel1_inf | 7:5 | RW | Configuration of wheel 1. " 000 " = normal, " 100 " = T1, " 101 " = E1, others reserved. |
| wheel2_ena | 8 | RW | Enable for wheel 2. |
| wheel2_inf | 11:9 | RW | Configuration of wheel 2. " 000 " = normal, " 100 " = T1, " 101 " = E1, others reserved. |
| wheel3_ena | 12 | RW | Enable for wheel 3. |
| wheel3_inf | 15:13 | RW | Configuration of wheel 3. " 000 " = normal, " 100 " $=\mathrm{T} 1, " 101 "=\mathrm{E} 1$, others reserved. |

Table 169 - Scheduler Configuration \& Enable 0 Register

Address: 612h Label: wheel_inf_1
Reset Value: $\overline{0} 00 \overline{0} \mathrm{~h}$

| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| :---: | :---: | :---: | :---: |
| wheel4_ena | 0 | RW | Enable for wheel 4. |
| wheel4_inf | 3:1 | RW | Configuration of wheel 4. "000" = normal," 100 " = T1, "101" = E1, others reserved. |
| wheel5_ena | 4 | RW | Enable for wheel 5. |
| wheel5_inf | 7:5 | RW | Configuration of wheel 5. "000" = normal, "100" = T1, "101" = E1, others reserved. |
| wheel6_ena | 8 | RW | Enable for wheel 6. |
| wheel6_inf | 11:9 | RW | Configuration of wheel 6. "000" = normal,"100" = T1, "101" = E1, others reserved |
| wheel7_ena | 12 | RW | Enable for wheel 7. |
| wheel7_inf | 15:13 | RW | Configuration of wheel 7. "000" = normal, "100" = T1, "101" = E1, others reserved. |

Table 170 - Scheduler Configuration \& Enable 1 Register

Address: 614h
Label: wheel_inf_2
Reset Value: $\mathbf{0} 000 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| wheel8_ena | 0 | RW | Enable for wheel 8. |
| wheel8_inf | $3: 1$ | RW | Configuration of wheel 8. "000" = normal, "100" $=$ T1, "101" $=$ E1, others <br> reserved. |
| wheel9_ena | 4 | RW | Enable for wheel 9. |
| wheel9_inf | $7: 5$ | RW | Configuration of wheel 9. "000" = normal, "100" $=$ T1, "101" $=$ E1, others <br> reserved. |
| wheelA_ena | 8 | RW | Enable for wheel 10. |
| wheelA_inf | $11: 9$ | RW | Configuration of wheel 10. "000" = normal, "100" $=$ T1, "101" $=$ E1, others <br> reserved. |
| wheelB_ena | 12 | RW | Enable for wheel 11. |
| wheelB_inf | $15: 13$ | RW | Configuration of wheel 11. "000" = normal, "100" $=$ T1, "101" $=$ E1, others <br> reserved. |

Table 171 - Scheduler Configuration \& Enable 2 Register

| Address: 616h <br> Label: wheel_inf_3 <br> Reset Value: $\overline{0} 00 \overline{0} \mathrm{~h}$ |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| wheelC_ena | 0 | RW | Enable for wheel 12. |
| wheelC_inf | 3:1 | RW | Configuration of wheel 12. "000" = normal, "100" = T1, "101" = E1, others reserved. |
| wheelD_ena | 4 | RW | Enable for wheel 13. |
| wheelD_inf | 7:5 | RW | Configuration of wheel 13. " 000 " = normal, " 100 " $=\mathrm{T} 1$, "101" = E1, others reserved. |
| wheelE_ena | 8 | RW | Enable for wheel 14. |
| wheelE_inf | 11:9 | RW | Configuration of wheel 14. "000" = normal, "100" = T1, "101" = E1, others reserved. |
| reserved | 15:12 | RW | Reserved. Must always be "0000" |

Table 172-Scheduler Configuration \& Enable 3 Register

### 5.5.7 RX_SAR Registers

Address: 700h
Label: control
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| always_diagnose | 0 | RW | When '1', an error report structure will be generated for <br> every cell that arrives. Used for tests. |
| reserved | $14: 1$ | RW | Reserved. Must always be "0000_0000_0000_00" |
| test_status | 15 | TS | When '1', all the status bits in the register will be set. |

Table 173-RX_SAR Control Register

Address: 702h
Label: status
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| rxsarpip_overflow | 0 | ROL | Overflow in the RX SAR data memory access cache. <br> Fatal chip error. |
| data_fifo_overflow | 1 | ROL | Overflow in the data cell FIFO in external memory. |
| error_fifo_overflow | 2 | ROL | Overflow in the error report structure FIFO in external <br> memory. |
| reserved | $15: 3$ | ROL | Reserved. Always read as "0000_0000_0000_0" |

Table 174-RX_SAR Status Register

Address: 704h
Label: status ie
Reset Value: $\overline{0} 000 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| rxsarpip_overflow_ie | 0 | IE | When '1' and the corresponding status bit is '1', an interrupt will be <br> generated. |
| data_fifo_overflow_ie | 1 | IE | When '1' and the corresponding status bit is '1', an interrupt will be <br> generated. |
| error_fifo_overflow_ie | 2 | IE | When '1' and the corresponding status bit is '1', an interrupt will be <br> generated. |
| reserved | $15: 3$ | IE | Reserved. Always read as "0000_0000_0000_0" |

Table 175 - RX_SAR Interrupt Enable Register

Address: 00708h
Label: data_cell_read
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| data_read_pnt | $13: 0$ | RW | The CPU's read pointer to the AAL0 cell FIFO. |
| reserved | $15: 14$ | RW | Reserved. Must always be "00" |

Table 176 - RX_SAR Data Read Pointer Register

| Address: 70Ah <br> Label: data_cellwrite <br> Reset Value: $\mathbf{0 0 0 0} \mathrm{h}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| data_write_pnt | $13: 0$ | RO | The chip's write pointer to the AAL0 cell FIFO. |
| reserved | $15: 14$ | RO | Reserved. Always read as "00" |

Table 177 - RX_SAR Data Write Pointer Register

Address: 70Ch
Label: data_cell_add
Reset Value ${ }^{\mathbf{0}} \mathbf{0 0 0} \mathrm{O} \mathrm{h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| data_add | $13: 0$ | RW | Bits 19:6 of the address of the FIFO in external memory |
| reserved | $15: 14$ | RW | Reserved. Must always be "00" |

Table 178 - RX_SAR Data Address Register

Address: 70Eh
Label: data_cell_size
Reset Value: 0000 h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :--- |
| data_size | $13: 0$ | RW | Size of data cell FIFO in 1 cell increments. All zeros $=16 \mathrm{k}$ <br> cells. |
| reserved | $15: 14$ | RW | Reserved. Must always be "00" |

Table 179-RX_SAR Data Cell Size Register

Address: 710h
Label: error_struct_read
Reset Value: $\mathbf{0 0 0 0 h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| error_read_pnt | $15: 0$ | RW | The CPU's read pointer to the error report structure FIFO. |

Table 180 - Error Structure Read Register

Address: 712h
Label: error_struct_write
Reset Value: $\mathbf{0 0 0 0 h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| error_write_pnt | $15: 0$ | RO | The chip's write pointer to the error report structure FIFO. |

Table 181 - Error Structure Write Register

Address: 714h
Label: error_struct_add_high
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| error_add[16] | 0 | RW | Bits 19:3 of the address of the error FIFO in external <br> memory |
| reserved | $15: 1$ | RW | Reserved. Must always be "0000_0000_0000_000" |

Table 182 - Error Structure Address High Register

Address: 716h
Label: error_struct_add_low
Reset Value: $\mathbf{0 0 0 0 h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| error_add[15:0] | $15: 0$ | RW | Bits 19:3 of the address of the error FIFO in external <br> memory |

Table 183 - Error Structure Address Low Register

Address: 718h
Label: error_struct_size
Reset Value: $\mathbf{0 0 0 0 h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| error_size | $15: 0$ | RW | Size of the error structure FIFO (in number of error <br> structures, 8-bytes each). |

Table 184 - Error Structure Size Register

| Address: 720h <br> Label: aal0_timeout_high <br> Reset Value: $\mathbf{0 0 0 0 h}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| aal0_timeout_period[19:16] | $3: 0$ | RW | Time, in us, that an AALO cell can wait in the FIFO <br> before an alarm is generated. |
| reserved | $15: 4$ | RW | Reserved. Must always be "0000_0000_0000" |

## Table 185-AALO Timeout High Register

```
Address: 722h
Label: aal0_timeout_low
Reset Value: 0000h
```

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| aal0_timeout_period[15:0 <br> ] | $15: 0$ | RW | Time, in us, that an AAL0 cell can wait in the FIFO before <br> an alarm is generated. |

Table 186-AALO Timeout Low Register

Address: 724h
Label: error_timeout_high
Reset Value: $\mathbf{0} 000 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| error_timeout_period[19:16] | $3: 0$ | RW | Time, in us, that an error structure can wait in the FIFO <br> before an alarm is generated. |
| reserved | $15: 4$ | RW | Reserved. Must always be "0000_0000_0000" |

Table 187 - Error Timeout High Register

Address: 726h
Label: error_timeout_low
Reset Value: $\mathbf{0 0 0 0 h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| error_timeout_period[15:0] | $15: 0$ | RW | Time, in us, that an error structure can wait in the FIFO <br> before an alarm is generated. |

Table 188 - Error Timeout Low Register

Address: 730h Label: treated_pulses Reset Value: 0 $\mathbf{0} 00 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| aal0_treated_pulse | 0 | PUL | Written to '1' to indicate that AALO cell FIFO has been <br> treated. Another alarm will not be generated until the <br> above timeout has elapsed. |
| error_treated_pulse | 1 | PUL | Written to '1' to indicate that error structure FIFO has <br> been treated. Another alarm will not be generated until <br> the above timeout has elapsed. |
| reserved | $15: 2$ | PUL | Reserved. Always read as "0000_0000_0000_00" |

Table 189 - Treated Pulses Register

### 5.5.8 Clock Registers

| Address: 800h <br> Label: control <br> Reset Value: 0000h | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| Label | $14: 0$ | RO | Reserved. Always Read as "0000_0000_0000_00" |
| reserved | 15 | TS | When '1', all the status bits in the register will be set. |
| test_status |  |  |  |

Table 190-Clock Control Register

Address: 802h
Label: status
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| mclk_count_alarm0 | 0 | ROL | Indicates that bits 31:16 of the mclk_count have reached the <br> value contained in the mclk_count_high_alarm0 field. |
| mclk_count_alarm1 | 1 | ROL | Indicates that bits 31:16 of the mclk_count have reached the <br> value contained in the mclk_count_high_alarm1 field. |
| mclk_count_alarm2 | 2 | ROL | Indicates that bits 31:16 of the mclk_count have reached the <br> value contained in the mclk_count_high_alarm2 field. |
| reserved | $15: 3$ | ROL | Reserved. Always read as "0000_0000_0000_0" |

Table 191 - Clock Status Register

| $\|$$\|$Address: 804h <br> Label: status_ie <br> Reset Value: $\mathbf{0 0 0 0 h}$ <br> Label <br> mclk_count_alarm0_ie <br> molk_count_alarm1_ie <br> Position | Type | Description |  |
| :---: | :---: | :---: | :---: |
| mclk_count_alarm2_ie | 1 | IE | When '1' and the corresponding status bit is '1', and <br> interrupt will be generated. |
| reserved | $15: 3$ | IE | When '1' and the corresponding status bit is '1', and <br> interrupt will be generated. |

Table 192 - Status Interrupt Enable Register

Address: 806h
Label: mclk_count_high_alm0
Reset Value: $\mathbf{0} 000 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| mclk_count_high_alarm0 | $15: 0$ | RW | This register in conjuction with an interrupt enable can be <br> used as a scheduler or as a Periodic Interrupt Controler. |

Table 193 - MCLK Alarm 0 Register

Address: 808h
Label: mclk_count_high
Reset Value: 0000 h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| mclk_count[31:16] | $15: 0$ | RO | Freerunning counter of mclk. |

Table 194 - MCLK Counter High Register

| Address: 80Ah <br> Label: mclk_count_low <br> Reset Value: 0000 h |
| :--- |
| Label |
| mclk_count[15:0] |
| Bit Position |
| 15:0 |
| Type |

Table 195 - MCLK Counter Low Register

Address: 80Ch
Label: mclk_count_high_alm1
Reset Value: $\mathbf{0 0 0 0} \mathbf{h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| mclk_count_high_alarm1 | $15: 0$ | RW | This register in conjuction with an interrupt enable can <br> be used as a scheduler or as a Periodic Interrupt <br> Controller. |

Table 196 - MCLK Alarm 1 Register

```
Address: 80Eh
Label: mclk_count_high_alm2
Reset Value: 0000h
```

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :---: |
| mclk_count_high_alarm2 | $15: 0$ | RW | This register in conjuction with an interrupt enable can be <br> used as a scheduler or as a Periodic Interrupt Controller. |

Table 197 - MCLK Alarm 2 Register

Address: 810h
Label: tx_srts_rego
Reset Value: 0 $\mathbf{0} 00 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tx_srts_enable | 0 | RW | When '1', TX SRTS values will be generated at the rate <br> indicated by the div_p and div_q registers. |
| tx_srts_bus_clk_sel | $2: 1$ | RW | Selects which H.100 Bus clock will be used to generate SRTS. <br> "0x"=ct_c8(which ever is used in slave timing) ; "10" = ct_c8_a; <br> "11"=ct_c8_b. |
| reserved | $7: 3$ | RW | Reserved. Must always be "0000_0" |
| tx_srts_fnxi_input_select | $13: 8$ | RW | Selects which pins must be used as the fnxi to generate the TX <br> SRTS value. See tableTable 28, "Source Selection," on page 85 <br> for a full description. |
| reserved | $15: 14$ | RW | Reserved. Must always be "00" |

Table 198 -TX_SRTS 0 Register

| Address: 812h <br> Label: tx_sts_reg1 <br> Reset Value: $\mathbf{0 0 0 0 h}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| tx_srts_overflow | 0 | ROL | This bit will be set by hardware if the number of TX SRTS <br> values sent by the TX SRTS generation module is <br> greater than the number of TX SRTS values read by the <br> TX SAR. |
| tx_srts_underflow | 1 | ROL | This bit will be set by hardware if the number of TX SRTS <br> values sent by the TX SRTS generation module is <br> smaller than the number of TX SRTS values read by the <br> TX SAR. |
| reserved | $15: 2$ | ROL | Reserved. Always read as "0000_0000_0000_00" |

Table 199 - TX_SRTS 1 Register

Address: 814h
Label: tx_srts_reg2
Reset Value: $\mathbf{0} \mathbf{0} 00 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tx_srts_overflow_ie | 0 | IE | When '1' and the corresponding status bit is '1', and interrupt will <br> be generated. |
| tx_srts_underflow_ie | 1 | IE | When '1' and the corresponding status bit is '1', and interrupt will <br> be generated. |
| reserved | $15: 2$ | IE | Reserved. Always read as "0000_0000_0000_00" |

Table 200 - TX_SRTS 2 Register

Address: 818h
Label: tx_srts_reg4
Reset Value: 00001h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tx_srts_srts8m8c_div_p | $15: 0$ | RW | The 8.192 MHz clock on the H .100 bus (ct_c8) must be divided <br> by a number K in order to match the interval of 8 SRTS carrying <br> cells. For example, a 24 channel AAL1 structured fully filled <br> channel would require a K of (375 / 24) $=15.625 . \mathrm{K} \mathrm{must} \mathrm{then}$ <br> be converted to the values P and Q using the following <br> equation: $\mathrm{K}=\mathrm{P} / \mathrm{Q}$. No rouding errors must be made in this <br> conversion. |

Table 201-TX_SRTS 4 Register

Address: 81Ah Label: tx_srts_reg5
Reset Value: 0001h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tx_srts_srts8m8c_div_q | $15: 0$ | RW | See Table 201, "TX_SRTS 4 Register," on page 165 for a <br> descrption of tx_srts_srts848c_div_p. |

Table 202 - TX_SRTS 5 Register

Address: 820h
Label: adapsrts0_reg0
Reset Value: 0000h

| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| :---: | :---: | :---: | :---: |
| adapsrts0_adaptive_enable | 0 | RW | When ' 1 ', the RX Adaptive block 0 is activated. Received cells on VCs tagged as clock recovery VC 'A' will generate an Adaptive point that will be written to external memory. |
| adapsrts0_rx_srts_enable | 1 | RW | When ' 1 ', the RX SRTS block 0 is activated. Received SRTS values on VCs tagged as clock recovery VC 'A' will be written to external memory. Also, local SRTS value will be generated using of the H. 100 clocks. |
| adapsrts0_ignore_crc | 2 | RW | When '1', this bit forces the CRC in the AAL1 Header to be ignored. CRC Errors are reported no matter the state of this bit. |
| adapsrts0_ignore_parity | 3 | RW | When ' 1 ', this bit forces the Parity bit in the AAL1 Header to be ignored. Parity Errors are reported no matter the state of this bit. |
| adapsrts0_ignore_seq_num | 4 | RW | When '1', this bit forces the Sequence Number $t$ in the AAL1 Header to be ignored. Sequence Number Errors are reported no matter the state of this bit. |
| adapsrts0_pclk_loss | 5 | RW | This bit can be directly routed out on a recov_X pin in order to convey the state (good / bad) of the clock generated by the adapsrts block 0 . |
| adapsrts0_pclk_divisor_load_now | 6 | PUL | When this bit is written to ' 1 ', the pclk_div and pclk_frc are loaded into the digital PLL used to synthesize the pclk. |
| adapsrts0_pclk_divisor_reset | 7 | RW | When ' 0 ', the digital PLL used to synthesize pclk is put in reset state. When ' 1 ', it is not longer in reset. |
| reserved | 15:8 | RW | Reserved. Must always be "0000_0000" |

Table 203 - Adaptive SRTS0 0 Register

## Address: 822h <br> Label: adapsrts0 reg1 <br> Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts0_aal1_crc_error | 0 | ROL | Set when an AAL1-byte CRC error occurs on cells <br> used in the adapsrts0 block. |
| adapsrts0_aal1_bad_parity | 1 | ROL | Set when an AAL1-byte parity error occurs on cells <br> used in the adapsrts0 block. |
| adapsrts0_single_cell_lost | 2 | ROL | Set when a single cell loss error occurs on cells used <br> in the adapsrts0 block. |


| adapsrts0_multi_cell_lost | 3 | ROL | S <br> us |
| :---: | :---: | :---: | :---: |


| adapsrts0_cell_misinserted | 4 | ROL | Set when a cell misinsertion error occurs on cells <br> used in the adapsrts0 block. |
| :---: | :---: | :---: | :---: |
| adapsrts0_timeout_flag | 5 | RO | Set when a the interval between two cells used in the <br> adapsrts0 block is greater than <br> adaprsts0_time_out_period. This bit is automatically <br> cleared by hardware when the timeout condition <br> ceases. |
| adapsrts0_timeout_pulse | 6 | ROL | Set when a the interval between two cells used in the <br> adapsrts0 block is greater than <br> adapsrts0_time_out_period. |
| adapsrts0_rx_srts_remote_overflo | 7 | ROL | Set when two consecutive remote SRTS value were <br> received, and when the second value could not be <br> stored because the interval was too short. |
| adapsrts0_rx_srts_local_overflow | 8 | ROL | Set when two consecutive local SRTS value were <br> received, and when the second value could not be <br> stored because the interval was too short. |
| reserved | $15: 9$ | ROL | Reserved. Always read as "0000_000" |

Table 204 - Adaptive SRTS0 1 Register

Address: 824h
Label: adapsrts0_reg2
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts0_aal1_crc_error_ie | 0 | IE | When '1' and the corresponding status bit is '1', <br> and interrupt will be generated. |
| adapsrts0_aal1_bad_parity_ie | 1 | IE | When '1' and the corresponding status bit is '1', <br> and interrupt will be generated. |
| adapsrts0_single_cell_lost_ie | 2 | IE | When '1' and the corresponding status bit is '1', <br> and interrupt will be generated. |

Table 205 - Adaptive SRTS0 2 Register

Address: 824h
Label: adapsrts0_reg2
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts0_multi_cell_lost_ie | 3 | IE | When '1' and the corresponding status bit is '1', <br> and interrupt will be generated. |
| adapsrts0_cell_misinserted_ie | 4 | IE | When '1' and the corresponding status bit is '1', <br> and interrupt will be generated. |
| reserved | 5 | RO | Reserved. Always read as "0" |
| adapsrts0_timeout_pulse_ie | 6 | IE | When '1' and the corresponding status bit is '1', <br> and interrupt will be generated. |
| adapsrts0_rx_srts_remote_overflow_i | 7 | IE | When '1' and the corresponding status bit is '1', <br> and interrupt will be generated. |
| adapsrts0_rx_srts_local_overflow_ie | 8 | IE | When '1' and the corresponding status bit is '1', <br> and interrupt will be generated. |
| reserved | $15: 9$ | IE | Reserved. Always read as "0000_000" |

Table 205 - Adaptive SRTS0 2 Register (continued)

Address: 826h
Label: adapsrts0_reg3
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts0_ref_input_select | $5: 0$ | RW | In adaptive mode, this field indicates which events the <br> adapsts0 block will consider as a timing reference. <br> Rising edges on the recov_X pins can be used. Cell <br> arrival events on Clock Recovery VC A or B can be <br> used. See Table 28, "Source Selection," on page 85" for <br> more details. |
| reserved | $7: 6$ | RO | Reserved. Always read as "00" |
| adapsrts0_rx_fnxi_input_select | $13: 8$ | RW | In SRTS mode, this selects the fnxi input used in the RX <br> SRTS block. See Table 28, "Source Selection," on <br> page 85" for more details. |
| reserved | $15: 14$ | RO | Reserved. Always read as "00" |

Table 206 - Adaptive SRTS0 3 Register

Address: 828h
Label: adapsrts0_reg4
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts0_time_out_period | $15: 0$ | RW | This value defines the time-out period between two cells <br> in both adaptive and SRTS mode. Unit is 1024 mclk <br> cycles. 0000h will disabled checking. |

Table 207 - Adaptive SRTS0 4 Register

Address: 82Ah
Label: adapsrts0_reg5
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts0_adap_pnt_elim_x | $7: 0$ | RW | In adaptive mode, this value defined how many points <br> will be deleted vs how many points will be kept and <br> written to external memory. It is defined as "keep 1 point <br> out of X". |
| reserved | $15: 8$ | RW | Reserved. Must always be "0000_0000" |

Table 208 - Adaptive SRTS0 5 Register

| Address: 82Ch <br> Label: adapsrts0_reg6 <br> Reset Value: 0001h |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| adapsrts0_srts8m8c_div_p | 15:0 | RW | The clock generated by the digital PLL (pclk at 8 KHz ) must be divided by a number K in order to match the interval of 8 SRTS carrying cells. For example, a 24 channel AAL1 structured fully filled channel would require a K of $(375 / 24)=15.625 . \mathrm{K}$ must then be converted to the values $P$ and $Q$ using the following equation: $K=P /$ Q. No rouding errors must be made in this conversion. |

Table 209 - Adaptive SRTS0 6 Register

Address: 82Eh Label: adapsrts0_reg7 Reset Value: 0001h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts0_srts8m8c_div_q | $15: 0$ | RW | See adapsrts0_srts8m8c_div_p description. |

Table 210 - Adaptive SRTS0 7 Register

Address: 830h
Label: adapsrts0_reg8
Reset Value: 2710h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :---: |
| adapsrts0_pclk_div | $15: 0$ | RW | In order to generate pclk, mclk must be divided by a factor K. This <br> factor is likely to be fractional. The integer part of K is written in the <br> form X / 65536, where X is written in the adapsrts0_pclk_frc register. |

Table 211 - Adaptive SRTS0 8 Register

Address: 832h
Label: adapsrts0_reg9
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts0_pclk_frc | $15: 0$ | RW | For a description see adapsrts0_pclk_div <br> adapsrts0_pclk_div above. |

Table 212 - Adaptive SRTS0 9 Register

Address: 840h
Label: adapsrts1_reg0
Reset Value: 0000̄h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts1_adaptive_enable | 0 | RW | See adapsrts0 registers 820h to 83Eh. |
| adapsrts1_rx_srts_enable | 1 | RW |  |
| adapsrts1_ignore_crc | 2 | RW |  |
| adapsrts1_ignore_parity | 3 | RW |  |
| adapsrts1_ignore_seq_num | 4 | RW |  |
| adapsrts1_pclk_loss | 5 | RW |  |
| adapsrts1_pclk_divisor_load_now | 6 | PUL |  |

Table 213 - Adaptive SRTS1 0 Register

```
Address: 840h
Label: adapsrts1_reg0
Reset Value: 0000̄
```

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts1_pclk_divisor_reset | 7 | RW |  |
| reserved | $15: 8$ | RW | Reserved. Must always be "0000_0000" |

Table 213 - Adaptive SRTS1 0 Register (continued)

Address: 842h
Label: adapsrts1_reg1
Reset Value: 0000̄

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts1_aal1_crc_error | 0 | ROL | See Address: 822h, Table 204, "Adaptive <br> SRTS0 1 Register," on page 167. |
| adapsrts1_aal1_bad_parity | 1 | ROL |  |
| adapsrts1_single_cell_lost | 2 | ROL |  |
| adapsrts1_multi_cell_lost | 3 | ROL |  |
| adapsrts1_cell_misinserted | 4 | ROL |  |
| adapsrts1_timeout_flag | 5 | RO |  |
| adapsrts1_timeout_pulse | 6 | ROL |  |
| adapsrts1_rx_srts_remote_overflow | 7 | ROL |  |
| adapsrts1_rx_srts_local_overflow | 8 | ROL |  |
| reserved | $15: 9$ | ROL | Reserved. Always read as "0000_000" |

Table 214 - Adaptive SRTS1 1Register

Address: 844h
Label: adapsrts1_reg2
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts1_aal1_crc_error_ie | 0 | IE |  |
| adapsrts1_aal1_bad_parity_ie | 1 | IE |  |
| adapsrts1_single_cell_lost_ie | 2 | IE |  |
| adapsrts1_multi_cell_lost_ie | 3 | IE | See "Address: 842h" on page 171. |
| adapsrts1_cell_misinserted_ie | 4 | IE |  |

Table 215 - Adaptive SRTS1 2 Register

Address: 844h
Label: adapsrts1_reg2
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| reserved | 5 | RO |  |
| adapsrts1_timeout_pulse_ie | 6 | IE |  |
| adapsrts1_rx_srts_remote_overflow_ie | 7 | IE |  |
| adapsrts1_rx_srts_local_overflow_ie | 8 | IE |  |
| reserved | $15: 9$ | IE | Reserved. Always read as "0000_000" |

Table 215 - Adaptive SRTS1 2 Register (continued)

Address: 846h
Label: adapsrts1_reg3
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts1_ref_input_select | $5: 0$ | RW | See "Address: 826h" on page 168. |
| reserved | $7: 6$ | RO |  |
| adapsrts1_rx_fnxi_input_select | $13: 8$ | RW |  |
| reserved | $15: 14$ | RO |  |

Table 216 - Adaptive SRTS1 3 Register

Address: 848h
Label: adapsrts1_reg4
Reset Value: 0000h

| Label | Bit Position | Type | Description |
| :---: | :---: | :---: | :---: |
| adapsrts1_time_out_period | $15: 0$ | RW | See "Address: 828h" on page 169. |

Table 217 - Adaptive SRTS1 4 Register

| Address: 84Ah <br> Label: adapsrts1_reg5 <br> Reset Value: 0001h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| adapsrts1_adap_pnt_elim_x | $7: 0$ | RW |  |
| reserved | $15: 8$ | RW | Reserved. Must always be "0000_0000" |

Table 218 - Adaptive SRTS1 5 Register

Address: 84Ch
Label: adapsrts1_reg6
Reset Value: 0000

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :---: |
| adapsrts1_srts8m8c_div_p | $15: 0$ | RW | See "Address: 82Ch" on page 169. |

Table 219 - Adaptive SRTS1 6 Register

Address: 84Eh
Label: adapsrts1_reg7
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :--- | :--- |
| adapsrts1_srts8m8c_div_q | $15: 0$ | RW | See "Address: 82Eh" on page 170. |

Table 220 - Adaptive SRTS1 7 Register

| Address: 850h <br> Label: adapsts1_reg8 <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| adapsrts1_pclk_div | $15: 0$ | RW | See "Address: 830h" on page 170. |

Table 221 - Adaptive SRTS1 8 Register

Address: 852h
Label: adapsrts1_reg9
Reset Value: 0000̄

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts1_pclk_frc | $15: 0$ | RW | See "Address: 832h" on page 170. |

Table 222 - Adaptive SRTS1 9 Register

Address: 860h
Label: pinmux_reg0
Reset Value: 0202h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| pinmux_recov_b_sel | $5: 0$ | RW | This field selects the recov_b output. See "Clockrec <br> Source Select" page for more details. |
| reserved | $7: 6$ | RO | Reserved. Always read as "00". |
| pinmux_recov_a_sel | $13: 8$ | RW | This field selects the recov_a output. See "Clockrec <br> Source Select" page for more details. |
| reserved | $15: 14$ | RO | Reserved. Always read as "00". |

Table 223 - Pin Mux 0 Register

Address: 862h
Label: pinmux_reg1
Reset Value: 0202h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| pinmux_recov_d_sel | $5: 0$ | RW | This field selects the recov_d output. See Table 28- <br> "Source Selection" on page 85 for more details. |
| reserved | $7: 6$ | RO | Reserved. Always read as "00" |
| pinmux_recov_c_sel | $13: 8$ | RW | This field selects the recov_c output. See Table 28 - <br> "Source Selection" on page 85 for more details. |
| reserved | $15: 14$ | RO | Reserved. Always read as "00" |

Table 224 - Pin Mux 1 Register

Address: 864h
Label: pinmux reg2
Reset Value: 0202h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| pinmux_recov_f_sel | $5: 0$ | RW | This field selects the recov_f output. See Table 28- <br> "Source Selection" on page 85 for more details. |
| reserved | $7: 6$ | RO | Reserved. Always read as "00" |
| pinmux_recov_e_sel | $13: 8$ | RW | This field selects the recov_e output. See Table 28- <br> "Source Selection" on page 85 for more details. |
| reserved | $15: 14$ | RO | Reserved. Always read as "00" |

Table 225 - Pin Mux 2 Register

Address: 866h
Label: pinmux_reg3
Reset Value: 0202h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| pinmux_recov_h_sel | $5: 0$ | RW | This field selects the recov_h output. See Table 28- <br> "Source Selection" on page 85 for more details. |
| reserved | $7: 6$ | RO | Reserved. Always read as "00" |
| pinmux_recov_g_sel | $13: 8$ | RW | This field selects the recov_g output. See Table 28- <br> "Source Selection" on page 85 for more details. |
| reserved | $15: 14$ | RO | Reserved. Always read as "00" |

Table 226 - Pin Mux 3 Register

Address: 868h
Label: pinmux_reg4
Reset Value: 0202h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| pinmux_ct_netref2_sel | $5: 0$ | RW | This field selects the ct_netref2 output. See Table 28- <br> "Source Selection" on page 85 for more details. |
| reserved | $7: 6$ | RO | Reserved. Always read as "00" |
| pinmux_ct_netref1_sel | $13: 8$ | RW | This field selects the ct_netref1 output. See Table 28 - <br> "Source Selection" on page 85 for more details. |
| reserved | $15: 14$ | RO | Reserved. Always read as "00" |

Table 227 - Pin Mux 4 Register

Address: 86Ah Label: pinmux_reg5
Reset Value: 0303h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| pinmux_local_netref_16m_sel | $5: 0$ | RW | This field selects the local_16m reference feed into the <br> H.100 master circuit. See Table 28 - "Source Selection" <br> on page 85 for more details. |
| reserved | $7: 6$ | RO | Reserved. Always read as "00" |
| reserved | $13: 8$ | RO | Reserved. Always read as "0000_00" |
| reserved | $15: 14$ | RO | Reserved. Always read as "00" |

Table 228 - Pin Mux 5 Register

Address: 880h
Label: diviclk0 reg0
Reset Value: 0300h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk0_clk_divisor_load_now | 0 | PUL | When this bit is written to '1', the integer clock divisor <br> 0's division value will be loaded into the divisor. |
| diviclk0_clk_divisor_reset | 1 | RW | When this bit is '0', the integer clock divisor 0 is put in <br> reset. Must be '1' for normal operation. |
| diviclk0_even_duty_cycle_selec | 2 | RW | When '1', the duty cycle modifier will be enabled and <br> till generate a 50\% duty cycle output. |
| diviclk0_input_invert_select | 3 | RW | When '1', the input of the integer clock divisor will be <br> inverted before being divided. |
| diviclk0_output_invert_select | 4 | RW | When '1', the output of the integer clock divisor will be <br> inverted before being sent out. |
| reserved | $12: 7$ | RW | Reserved. Must always be "0000_00" <br> diviclk0_input_source_select <br> $13: 8$ <br> RW <br> This field selects the input of the integer clock divisor. <br> See Table 28 - "Source Selection" on page 85 for <br> more details. <br> reserved |
| $15: 14$ | RW | Reserved. Must always be "00" |  |

Table 229 - Integer Clock Divisor0 0 Register

| Address: 882h <br> Label: diviclk0_reg1 <br> Reset Value: $\mathbf{0 0 0 0} \mathbf{0}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type |  |
| diviclk0_ext_loss_pulse | 0 | ROLO | This bit is set when an external signal indicates that the <br> input of the integer clock divisor is invalid (such as a PHY <br> alarm). |
| diviclk0_ext_loss_flag | 1 | RO | Same as diviclk0_ext_loss_pulse, but when the error <br> condition ceases, this bit will clear itself. |
| diviclk0_freq_too_high_pulse | 2 | ROLO | This bit is set when the integer clock divisor input's <br> frequency is off as compared to the expected frequency <br> (received frequency is too high). |
| diviclk0_freq_too_high_flag | 3 | RO | Same as diviclk0_freq_too_high_pulse, but when the error <br> condition ceases, this bit will clear itself. |
| diviclk0_freq_too_low_pulse | 4 | ROLO | This bit is set when the integer clock divisor input's <br> frequency is off as compared to the expected frequency <br> (received frequency is too low). |
| diviclk0_freq_too_low_flag | 5 | RO | Same as diviclk0_freq_too_low_pulse, but when the error <br> condition ceases, this bit will clear itself. |
| reserved | $15: 6$ | RO | Reserved. Always read as "0000_0000_00" |

Table 230 - Integer Clock Divisor0 1 Register

Address: 884h
Label: diviclk0_reg2
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk0_ext_loss_pulse_ie | 0 | IE | When '1' and the corresponding status bit is '1', and <br> interrupt will be generated. |
| reserved | 1 | RO | Reserved. Always read as "00" |
| diviclk0_freq_too_high_pulse_ie | 2 | IE | When '1' and the corresponding status bit is '1', and <br> interrupt will be generated. |
| reserved | 3 | RO | Reserved. Always read as "00" |
| diviclk0_freq_too_low_pulse_ie | 4 | IE | When '1' and the corresponding status bit is '1', and <br> interrupt will be generated. |
| reserved | 5 | RO | Reserved. Always read as "00" |
| reserved | $15: 6$ | RO | Reserved. Always read as "0000_0000_00" |

Table 231 - Integer Clock Divisor0 2 Register

| Address: 886h <br> Label: diviclk0_reg3 <br> Reset Value: 0000 h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| diviclk0_ext_loss_pulse_genloss | 0 | RW | When set, the corresponding status will cause a <br> loss to be signaled for this diviclk. |
| diviclk0_ext_loss_flag_genloss | 1 | RW | When set, the corresponding status will cause a <br> loss to be signaled for this diviclk. |
| diviclk0_freq_too_high_pulse_genloss | 2 | RW | When set, the corresponding status will cause a <br> loss to be signaled for this diviclk. |
| diviclk0_freq_too_high_flag_genloss | 3 | RW | When set, the corresponding status will cause a <br> loss to be signaled for this diviclk. |
| diviclk0_freq_too_low_pulse_genloss | 4 | RW | When set, the corresponding status will cause a <br> loss to be signaled for this diviclk. |
| diviclk0_freq_too_low_flag_genloss | 5 | RW | When set, the corresponding status will cause a <br> loss to be signaled for this diviclk. |
| reserved | $15: 6$ | RW | Reserved. Must always be "0000_0000_00" |

Table 232 - Integer Clock Divisor0 3 Register

## Address: 888h Label: diviclk0_reg4 <br> Reset Value: 0063h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk0_ext_loss_source_select | $5: 0$ | RW | Select's which external pin will signal that the input of <br> the integer clock divisor is good or not. See Table <br> $28-$ "Source Selection" on page 85 for more details. |
| diviclk0_ext_loss_source_polarity | 6 | RW | Input of the integer clock divisor is considered bad <br> when the select ext_loss signal is active (i.e. equal to <br> this bit). 0 = source loss active low; 1 $1=$ source loss <br> active high. |
| diviclk0_output_loss_polarity | 7 | RW | When a loss is detected, the output signal indicating <br> this loss is activated (i.e. the value is this register is <br> sent out). '0' $=$ output loss active low; '1' $=$ output loss <br> active high. |
| reserved | $15: 8$ | RW | Reserved. Must always be "0000_0000" |

Table 233 - Integer Clock Divisor0 4 Register

Address: 88Ah
Label: diviclk0 reg5
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk0_clk_div | $15: 0$ | RW | This is the denominator used to divide the input clock. It <br> ranges from 1 to 65535. |

Table 234 - Integer Clock Divisor0 5 Register

## Address: 890h

Label: diviclk0_reg8
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk0_freqchck_div | $15: 0$ | RW | This value will be used to divide the input clock before check <br> it against the frequency of mclk. The higher the division <br> value, the longer to detect a frequency too high or frequency <br> too low, but the check is more precise. 0000h disables <br> frequency checking. |

Table 235 - Integer Clock Divisor0 8 Register

Address: 892h
Label: diviclk0_reg9
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk0_freqchck_max_mclk_cycles | $15: 0$ | RW | This regsiter defines the maximum number of <br> mclk cycles between two rising edges of the input <br> clock divided by diviclk0_freqchck_div. If a second <br> rising edge has not been detected in <br> diviclk0_freqchck_max_mclk_cycles, the <br> diviclk0_freq_too_low will be detected. |

Table 236 - Integer Clock Divisor0 9 Register

Address: 894h Label: diviclk0_reg10
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk0_freqchck_min_mclk_cycles | $15: 0$ | RW | This regsiter defines the minimum number of mclk <br> cycles between two rising edges of the input clock <br> divided by diviclk0_freqchck_div.lf a second rising <br> edge has been detected in less then <br> diviclk0_freqchck_min_mclk_cycles, the <br> diviclk0_freq_too_high will be detected. |

Table 237 - Integer Clock Divisor0 10 Register

Address: 8A0h
Label: diviclk1_reg0
Reset Value: 0300h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk1_clk_divisor_load_now | 0 | PUL | See regsiters 880h to 89Eh |
| diviclk1_clk_divisor_reset | 1 | RW |  |
| diviclk1_even_duty_cycle_select | 2 | RW |  |
| diviclk1_input_invert_select | 3 | RW |  |
| diviclk1_output_invert_select | 4 | RW |  |
| diviclk1_input_source_select | $13: 8$ | RW |  |
| reserved | $15: 14$ | RW | Reserved. Must always be "00" |

Table 238 - Integer Clock Divisor1 0 Register

Address: 8A2h
Label: diviclk1_reg1
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk1_ext_loss_pulse | 0 | ROLO | See register 882h. |
| diviclk1_ext_loss_flag | 1 | RO |  |
| diviclk1_freq_too_high_pulse | 2 | ROLO |  |
| diviclk1_freq_too_high_flag | 3 | RO |  |
| diviclk1_freq_too_low_pulse | 4 | ROLO |  |
| diviclk1_freq_too_low_flag | 5 | RO |  |
| reserved | $15: 6$ | RO | Reserved. Always read as "0000_0000_00" |

Table 239 - Integer Clock Divisor1 1 Register

| Address: 8A4h Label: diviclk1_reg2 Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| diviclk1_ext_loss_pulse_ie | 0 | IE | See register 884h. |
| reserved | 1 | RO |  |
| diviclk1_freq_too_high_pulse_ie | 2 | IE |  |
| reserved | 3 | RO |  |
| diviclk1_freq_too_low_pulse_ie | 4 | IE |  |
| reserved | 5 | RO |  |
| reserved | 15:6 | RO | Reserved. Always read as "0000_0000_00" |

Table 240 - Integer Clock Divisor1 2 Register

Address: 8A6h
Label: diviclk1_reg3
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk1_ext_loss_pulse_genloss | 0 | RW | See register 886h. |
| diviclk1_ext_loss_flag_genloss | 1 | RW |  |
| diviclk1_freq_too_high_pulse_genloss | 2 | RW |  |
| diviclk1_freq_too_high_flag_genloss | 3 | RW |  |
| diviclk1_freq_too_low_pulse_genloss | 4 | RW |  |
| diviclk1_freq_too_low_flag_genloss | 5 | RW |  |
| reserved | $15: 6$ | RW | Reserved. Must always be "0000_0000_00" |

Table 241 - Integer Clock Divisor1 3 Register

Address: 8A8h
Label: diviclk1_reg4
Reset Value: 0063h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk1_ext_loss_source_select | $5: 0$ | RW | See register 888h. |
| diviclk1_ext_loss_source_polarity | 6 | RW |  |
| diviclk1_output_loss_polarity | 7 | RW |  |
| reserved | $15: 8$ | RW | Reserved. Must always be "0000_0000" |

Table 242-Integer Clock Divisor1 4 Register

| Address: 8AAh <br> Label: diviclk1_reg5 <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| diviclk1_clk_div | $15: 0$ | RW | See register 88Ah. |

Table 243 - Integer Clock Divisor1 5 Register

Address: 8B0h
Label: diviclk1_reg8
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk1_freqchck_div | $15: 0$ | RW | See register 890h. |

Table 244 - Integer Clock Divisor1 8 Register

Address: 8B2h
Label: diviclk1_reg9
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :--- | :---: | :--- | :--- |
| diviclk1_freqchck_max_mclk_cycles | $15: 0$ | RW | See register 892h. |

Table 245 - Integer Clock Divisor1 9 Register

Address: 8B4h
Label: diviclk1_reg10
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :--- | :--- |
| diviclk1_freqchck_min_mclk_cycles | $15: 0$ | RW | See register 894h. |

Table 246 - Integer Clock Divisor1 10 Register

## Address: 8COh <br> Label: diviclk2_reg0 <br> Reset Value: 0300h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk2_clk_divisor_load_now | 0 | PUL | See registers 880h to 89Eh |
| diviclk2_clk_divisor_reset | 1 | RW |  |
| diviclk2_even_duty_cycle_select | 2 | RW |  |
| diviclk2_input_invert_select | 3 | RW |  |
| diviclk2_output_invert_select | 4 | RW |  |
| diviclk2_input_source_select | $13: 8$ | RW |  |
| reserved | $15: 14$ | RW | Reserved. Must always be "00" |

Table 247 - Integer Clock Divisor2 0 Register

Address: 8C2h
Label: diviclk2_reg1
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk2_ext_loss_pulse | 0 | ROLO | See register 8A2h. |
| diviclk2_ext_loss_flag | 1 | RO |  |
| diviclk2_freq_too_high_pulse | 2 | ROLO |  |
| diviclk2_freq_too_high_flag | 3 | RO |  |
| diviclk2_freq_too_low_pulse | 4 | ROLO |  |
| diviclk2_freq_too_low_flag | 5 | RO |  |
| reserved | $15: 6$ | RO | Reserved. Always read as "0000_0000_00" |

Table 248 - Integer Clock Divisor2 1 Register

## Address: 8C4h

Label: diviclk2_reg2
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk2_ext_loss_pulse_ie | 0 | IE | See register 8A4h. |
| reserved | 1 | RO |  |
| diviclk2_freq_too_high_pulse_ie | 2 | IE |  |
| reserved | 3 | RO |  |

Table 249 - Integer Clock Divisor2 2 Register

Address: 8C4h
Label: diviclk2_reg2
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk2_freq_too_low_pulse_ie | 4 | IE |  |
| reserved | 5 | RO |  |
| reserved | $15: 6$ | RO | Reserved. Always read as "0000_0000_00" |

Table 249 - Integer Clock Divisor2 2 Register (continued)

## Address: 8C6h Label: diviclk2_reg3 <br> Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk2_ext_loss_pulse_genloss | 0 | RW | See register 8A6h. |
| diviclk2_ext_loss_flag_genloss | 1 | RW |  |
| diviclk2_freq_too_high_pulse_genloss | 2 | RW |  |
| diviclk2_freq_too_high_flag_genloss | 3 | RW |  |
| diviclk2_freq_too_low_pulse_genloss | 4 | RW |  |
| diviclk2_freq_too_low_flag_genloss | 5 | RW |  |
| reserved | $15: 6$ | RW | Reserved. Must always be "0000_0000_00" |

Table 250 - Integer Clock Divisor2 3 Register

Address: 8C8h
Label: diviclk2 reg4
Reset Value: 0063h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk2_ext_loss_source_select | $5: 0$ | RW | See register 8A8h. |
| diviclk2_ext_loss_source_polarity | 6 | RW |  |
| diviclk2_output_loss_polarity | 7 | RW |  |
| reserved | $15: 8$ | RW | Reserved. Must always be "0000_0000" |

Table 251 - Integer Clock Divisor2 4 Register

## Address: 8CAh

Label: diviclk2_reg5
Reset Value: 0001h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk2_clk_div | $15: 0$ | RW | See register 8AAh. |

Table 252 - Integer Clock Divisor2 5 Register

| Address: 8D0h <br> Label: diviclk2_reg8 <br> Reset Value: 0000h |  |  |  |
| :--- | :---: | :--- | :--- | :--- |
| Label | Bit <br> Position | Type | Description |
| diviclk2_freqchck_div | $15: 0$ | RW | See register 8B0h. |

Table 253 - Integer Clock Divisor2 8 Register

Address: 8D2h
Label: diviclk2_reg9
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk2_freqchck_max_mclk_cycles | $15: 0$ | RW | See register 8B2h. |

Table 254 - Integer Clock Divisor2 9 Register

Address: 8D4h
Label: diviclk2_reg10
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| diviclk2_freqchck_min_mclk_cycles | $15: 0$ | RW | See register 8B4h. |

Table 255 - Integer Clock Divisor2 10 Register

Address: 8EOh
Label: tx_srts_debug
Reset Value: $\mathbf{0} \overline{0} 00 h$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| tx_srts_eight_cell_pulse_rol | 0 | ROL | Only used for tests. |
| tx_srts_fnxi_cnt | $11: 8$ | RO | Only used for tests. |

Table 256-TX SRTS Debug Register

Address: 8EOh
Label: tx_srts_debug
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| reserved | $15: 12$ | RO | Reserved. Always read as "0000" |

Table 256-TX SRTS Debug Register (continued)

| Address: 8E4h <br> Label: adapsrts0_rx_srts_debug <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{aligned} & \text { Bit } \\ & \text { Position } \end{aligned}$ | Type | Description |
| adapsrts0_eight_cell_pulse_rol | 0 | ROL | Only used for tests. |
| adapsrts0_rx_srts_fnxi_cnt | 11:8 | RO | Only used for tests. |
| reserved | 15:12 | RO | Reserved. Always read as "0000" |

Table 257-RX SRTS Debug 0 Register

Address: 8E8h
Label: adapsrts1_rx_srts_debug
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adapsrts1_eight_cell_pulse_rol | 0 | ROL | Only used for tests. |
| adapsrts1_rx_srts_fnxi_cnt | $11: 8$ | RO | Only used for tests. |
| reserved | $15: 12$ | RO | Reserved. Always read as "0000" |

Table 258-RX SRTS Debug 1 Register

| Address: 8ECh <br> Label: adapsrts_aal1_err_debug <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| adapsrts0_valid_cell | 0 | ROL | Only used for tests. |
| adapsrts0_cell_lost | 1 | ROL | Only used for tests. |
| adapsrts1_valid_cell | 2 | ROL | Only used for tests. |
| adapsrts1_cell_lost | 3 | ROL | Only used for tests. |
| reserved | 15:4 | ROL | Reserved. Always read as "0000_0000_0000" |

Table 259-AAL1 Error Debug Register

### 5.5.9 Miscellaneous Registers

Address: 900h
Label: control
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| test_status | 15 | TS | When '1', all the status bits in the register will be set. |

Table 260 - Miscellaneous Control Register

Address: 902h
Label: err_reg0
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adap_srts_remote0_mem_overflow | 0 | ROL | Indicates that the chip did not have time to write an <br> adaptive point structure before the next one was <br> generated. |
| adap_srts_remote1_mem_overflow | 1 | ROL | Indicates that the chip did not have time to write an <br> adaptive point structure before the next one was <br> generated. |
| srts_local0_mem_overflow | 2 | ROL | Indicates that the SRTS value buffer overflowed. |
| srts_local1_mem_overflow | 3 | ROL | Indicates that the SRTS value buffer overflowed. |
| cas_mem_overflow | 4 | ROL | Indicates that the CAS change memory <br> overflowed. |
| silent_tone_error | 5 | ROL | New silent tones were requested before the <br> current ones could be fetched. |
| reserved | $15: 6$ | ROL | Reserved. Always read as "0000_0000_00" |

Table 261 - Miscellaneous Error Register

Address: 908h
Label: silent_tone_reg2
Reset Value: $\overline{0} 000 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| silent_size | $15: 0$ | RW | Size of the silent tone buffers, in bytes minus 1. 0 means <br> 1 byte; FFFFh means 10000h bytes. |

Table 262 - Silent Tone 2 Register

Address: 90Ah
Label: silent_tone_reg3
Reset Value: $\mathbf{0} 000 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| silent_base_add_15_0 | $15: 0$ | RW | The base address of the silent tone buffers (in words). |

Table 263 - Silent Tone 3 Register

Address: 90Ch
Label: silent_tone_reg3
Reset Value: $\mathbf{0} 000 \mathrm{~h}$

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| silent_base_add_18_16 | 2:0 | RW | The base address of the silent tone buffers (in words). |
| reserved | $15: 3$ | RW | Reserved. Must always be "0000_0000_0000_0" |

Table 264 - Silent Tone 4 Register

Address: 90Eh
Label: as0_srts_reg0
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adap_srts0_size | $15: 0$ | RW | Size of the adaptive point/SRTS value buffer minus one. 0 means <br> $1 ;$ FFFFh means 10000h. |

Table 265 - Adaptive Point/SRTS Value 0 Register

```
Address: 910h
Label: as0_srts_reg1
Reset Value: 0000h
```

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adap_srts0_base_add_15_0 | $15: 0$ | RW | The base address of the adaptive point/SRTS value <br> buffer (in points or SRTS values). Note that since <br> adaptive points are 8 words, bits 18:16 of the base <br> address are ignored. |

Table 266 - Adaptive Point/SRTS Base Address Low 0 Register

| Address: 912h <br> Label: as0_srts_reg2 <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| adap_srts0_base_add_18_16 | $2: 0$ | RW | The base address of the adaptive point/SRTS value <br> buffer (in points or SRTS values). Note that since <br> adaptive points are 8 words, bits 18:16 of the base <br> address are ignored. |
| reserved | $15: 3$ | RW | Reserved. Must always be "0000_0000_0000_0" |

Table 267 - Adaptive Point/SRTS Base Address High 0 Register

| Address: 914h <br> Label: as0_srts_reg3 <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| adap_srts_remote0_write_pnt | 15:0 | RO | The chip's write pointer to the adaptive point/SRTS value buffer. |

Table 268 - Adaptive Point/SRTS Write Pointer 0 Register

| Address: 916h <br> Label: as0_srts_reg4 <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| adap_srts_remote0_read_pnt | $15: 0$ | RW | The CPU's read pointer to the adaptive point/SRTS value <br> buffer. |

Table 269 - Adaptive Point/SRTS Read Pointer 0 Register

Address: 918h
Label: as0_srts_reg5
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| srts_local0_write_pnt | $15: 0$ | RO | The chip's write pointer to the local SRTS value buffer. |

Table 270 - Local SRTS Write Pointer 0 Register

Address: 91Ah
Label: as0_srts_reg6
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| srts_local0_read_pnt | $15: 0$ | RW | The CPU's read pointer to the local SRTS value buffer. |

Table 271 - Local SRTS Read Pointer 0 Register

Address: 91Ch
Label: as1_srts_reg0
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adap_srts1_size | $15: 0$ | RW | Size of the adaptive point/SRTS value buffer minus one. 0 means <br> $1 ;$ FFFFh means 10000h. |

Table 272 - Adaptive Point/SRTS Value 1 Register

```
Address: 91Eh
Label: as1 srts reg1
Reset Value: 0000h
```

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| adap_srts1_base_add_15_0 | $15: 0$ | RW | The base address of the adaptive point/SRTS value <br> buffer (in points or SRTS values). Note that since <br> adaptive points are 8 words, bits 18:16 of the base <br> address are ignored. |

Table 273-Adaptive Point/SRTS Base Address Low 1 Register

| Address: 920h <br> Label: as1_srts_reg2 <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| adap_srts1_base_add_18_16 | $2: 0$ | RW | The base address of the adaptive point/SRTS value <br> buffer (in points or SRTS values). Note that since <br> adaptive points are 8 words, bits 18:16 of the base <br> address are ignored. |
| reserved | $15: 3$ | RW | Reserved. Must always be "0000_0000_0000_0" |

Table 274 - Adaptive Point/SRTS Base Address High 1 Register

| Address: 922h <br> Label: as1_srts_reg3 <br> Reset Value: 0000 h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| adap_srts_remote1_write_pnt | $15: 0$ | RO | The chip's write pointer to the adaptive point/SRTS value <br> buffer. |

Table 275 - Adaptive Point/SRTS Write Pointer 1 Register

| Address: 924h <br> Label: as1_sts_reg4 <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| adap_srts_remote1_read_pnt | $15: 0$ | RW | The CPU's read pointer to the adaptive point/SRTS <br> value buffer. |

Table 276 - Adaptive Point/SRTS Read Pointer 1 Register

Address: 926h
Label: as1_srts_reg5
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| srts_local1_write_pnt | $15: 0$ | RO | The chip's write pointer to the local SRTS value buffer. |

Table 277 - Local SRTS Write Pointer 1 Register

Address: 928h
Label: as1_srts_reg6
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| srts_local1_read_pnt | $15: 0$ | RW | The CPU's read pointer to the local SRTS value buffer. |

Table 278 - Local SRTS Read Pointer 1 Register

Address: 92Ah
Label: cas_rego
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| cas_size | $15: 0$ | RW | Size of the CAS change buffer minus 1. 0 means 1; <br> FFFFh means 10000h. |

Table 279 - CAS Change Buffer Size Register

| Address: 92Ch <br> Label: cas_reg1 <br> Reset Value: 0000 h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| cas_base_add_15_0 | $15: 0$ | RW | Base address of the CAS change buffer. |

Table 280 - CAS Change Buffer Base Address Low Register

| Address: 92Eh <br> Label: cas_reg2 <br> Reset Value: 0000 h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| cas_base_add_18_16 | $2: 0$ | RW | Base address of the CAS change buffer. |
| reserved | $15: 3$ | RW | Reserved. Must always be "0000_0000_0000_0" |

Table 281 - CAS Change Buffer Base Address High Register

## Address: 930h

Label: cas_reg3
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| cas_write_pnt | $15: 0$ | RO | The chip's write pointer to the CAS change buffer. |

Table 282-CAS Write Pointer Register

| Address: 932 h <br> Label: cas_reg4 <br> Reset Value: 0000 h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type | Description |
| cas_read_pnt | $15: 0$ | RW | The CPU's read pointer to the CAS change buffer. |

Table 283-CAS Read Pointer Register

| Address: 944h <br> Label: cas_timeout_hig <br> Reset Value: 0000 h |
| :--- |
| Label |
| cas_timeout_period[19:16] |
| Position | Type | Bit | Description |  |
| :---: | :---: | :--- |
| reserved | $15: 4$ | RW |

Table 284 - CAS Timeout High Register

Address: 946h
Label: cas_timeout_low
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| cas_timeout_period[15:0] | $15: 0$ | RW | Time, in us, that a CAS change report can wait in the FIFO <br> before an alarm is generated. |

Table 285 - CAS Timeout Low Register

| Address: 948h <br> Label: treated pulses <br> Reset Value: 0000 h |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type |  |
| cas_treated_pulse | 0 | PUL | Written to '1' to indicate that CAS change FIFO has been <br> treated. Another alarm will not be generated until the above <br> timeout has elapsed. |
| reserved | $15: 1$ | PUL | Reserved. Always read as "0000_0000_0000_000" |

Table 286 - Treated Pulses Register

### 5.5.10 H. 100 Registers

| Address: A00h <br> Label: control0 <br> Reset Value: 0000h |  |  |  |
| :---: | :---: | :---: | :---: |
| Label | $\begin{gathered} \text { Bit } \\ \text { Position } \end{gathered}$ | Type | Description |
| h100_pll_ref_fallback | 0 | RW | ' 0 ' = always use selected external clk (selected by h100_pll_ext_source) to source pll's ref input. '1' = if selected clock fails, fallback on local 16.384 mHz clock. |
| h100_pll_ext_source | 1 | RW | '0' = source ct_c8_a_in. '1' = source ct_c8_b_in |
| h100_pll_local_source | 2 | RW | '0' = source pll's ref input with local 16.384 mHz clock, '1' = source with selected external clk (selected via h100_pll_ext_source). |
| h100_pll_override | 3 | RW | '0' = use embedded PLL to generate outgoing clocks. '1' = bypass PLL and use local 16.384 mHz clock instead. |
| h100_clk_loopback | 4 | RW | '1' = loops the ct_c8 A and B clocks and frames back into the chip. Used for tests |
| h100_sclk_speed | 6:5 | RW | "00" $=2.048 \mathrm{mHz} .010=4.096 \mathrm{mHz} .100=8.192 \mathrm{mHz}$. |
| h100_c8_frame_a_oe | 7 | RW | '0' = tri-states ct_c8_a and ct_frame_a |
| h100_c8_frame_b_oe | 8 | RW | '0' = tri-states ct_c8_b and ct_frame_b |
| h100_comp_oe | 9 | RW | ' 0 ' = tri-states all the compatibility signals |
| h100_pll_fb_input | 11:10 | RW | Selects pll's fb input: "00" = sample A, "10" = sample B, "01" = sample PLL output (div 4), "11" = reserved. |
| h100_frame_sync_sourc | 12 | RW | '0' = sync frame on ct_frame_a_in, '1' = sync frame on ct_frame_b_in |
| h100_tdmint_clk_sel | 13 | RW | specifies which clk to send to tdmint: '0' = use ct_c8_a. '1' = use ct_c8_b. h100_tdmint_clk_fallback can send other clk instead. |
| h100_tdmint_clk_fallbac | 14 | RW | clk to tdmint if selected clk is bad: '0' = always sync on selected clk (h100_tdming_clk_sel). '1' = if selected clk fails, switch over to backup clk. |
| test_status | 15 | TS | When '1', all the status bits in the register will be set. |

Table 287-H. 100 Control 0 Register

Address: A02h
Label: control1
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| h100_force_frame_sync_local | 0 | RW | Only applies if h100_pll_override is '0': '0' = sync frame <br> on selected external frame <br> (h100_frame_sync_source), '1' = sync frame on local <br> 16.384 MHz clk. |
| reserved | $15: 1$ | RW | Reserved. Must always be "0000_0000_0000_000" |

Table 288-H.100 Control 1 Register

Address: A04h
Label: control2
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| h100_min_mclk_ct_c8 | $3: 0$ | RW | Minimum number of mclk cycles between ct_c8 rising edges, <br> typically set using this equation: ((122-35)// <br> mclk_period_ns) - 2. |
| h100_max_mclk_ct_c8 | $8: 4$ | RW | Maximum number of mclk cycles between ct_c8 rising <br> edges, typically set using this equation: ((122 + 35)/ <br> mclk_period_ns) + 2. |
| reserved | $15: 9$ | RW | Reserved. Must always be "0000_000" |

Table 289 - H. 100 Control 2 Register
Address: A08h
Label: flags
Reset Value: 0000h

| Label | Bit <br> Position | Type | Description |
| :---: | :---: | :---: | :--- |
| h100_clk_a_bad_latched | 0 | RO | Indicates that the ct_c8_a period is not within +/-3 35 ns of <br> what is was supposed to be. This is a RO signal, so it <br> means that the error is currently occurring. |
| h100_clk_b_bad_latched | 1 | RO | Indicates that the ct_c8_b period is not within +/-3 35 ns of <br> what is was supposed to be. This is a RO signal, so it <br> means that the error is currently occurring. |
| h100_frame_a_bad_latched | 2 | RO | Indicates that the ct_frame_a is not occurring every 1024 <br> ct_c8_a cycles. This is a RO signal, so it means that the <br> error is currently occurring. |
| h100_frame_b_bad_latched | 3 | RO | Indicates that the ct_frame_b is not occurring every 1024 <br> ct_8_a cycles. This is a RO signal, so it means that the <br> error is currently occurring. |
| reserved | $15: 4$ | RO | Reserved. Always read as "0000_0000_0000" |

Table 290-H. 100 Flags Register

| Address: A0Ah <br> Label: status <br> Reset Value: $\mathbf{0 0 0 0 h}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type |  |
| h100_clk_a_bad_rol | 0 | ROLO | Indicates that the ct_c8_a period is not within +/- 35 ns of what <br> is was supposed to be. This is a ROL signal, so it means that <br> the error has occurred since the last time this bit was cleared. |
| h100_clk_b_bad_rol | 1 | ROLO | Indicates that the ct_c8_b period is not within +/- 35 ns of what <br> is was supposed to be. This is a ROL signal, so it means that <br> the error has occurred since the last time this bit was cleared. |
| h100_frame_a_bad_rol | 2 | ROLO | Indicates that the ct_frame_a is not occurring every 1024 <br> ct_c8_a cycles. This is a ROL signal, so it means that the error <br> has occurred since the last time this bit was cleared. |
| h100_frame_b_bad_rol | 3 | ROLO | Indicates that the ct frame_b is not occurring every 1024 <br> ct_8_a cycles. This is a ROL signal, so it means that the error <br> has occurred since the last time this bit was cleared. |
| reserved | $15: 4$ | ROLO | Reserved. Always read as "0000_0000_0000" |

Table 291-H. 100 Status Register

| Address: A0Ch <br> Label: status_ie <br> Reset Value: $\mathbf{0 0 0 0 h}$ |  |  |  |
| :---: | :---: | :---: | :--- |
| Label | Bit <br> Position | Type |  |
| h100_clk_a_bad_rol | 0 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| h100_clk_b_bad_rol | 1 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| h100_frame_a_bad_rol | 2 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| h100_frame_b_bad_rol | 3 | IE | When '1' and the corresponding status bit is '1', an interrupt will <br> be generated. |
| reserved | $15: 4$ | IE | Reserved. Always read as "0000_0000_0000' |

Table 292-H. 100 Interrupt Enable Register

### 6.0 Statistics

### 6.1 TDM statistics

Underrun counter: 16 -bit counter that counts the number of underruns detected by TDM interface.

### 6.2 TX SAR statistics

Percentage of bandwidth utilisation: a register (0510h) which indicates how many mclk cycles were required to treat the last frame.

Transmitted Cell Counter: 32-bit counter that counts the number of cells transmitted on a particular VC. Each VC has its own counter in its structure.

### 6.3 RX SAR statistics

Error reporting structures allow software-based counters for the following errors:

- P-byte absent error
- P-byte framing error
- P-byte range error
- P-byte parity error
- Overrun error
- Underrun error
- AAL1 CRC error
- AAL1 parity error
- Single cell loss
- Multiple cell loss
- Cell misinsertion

Received Cell Counter: 32-bit counter that counts the number of cells received on a particular VC. Each VC has its own counter in its structure.

### 6.4 UTOPIA statistics

Transmitted Cell Counters: 32-bit counters that counts the number of cells transmitted on a particular UTOPIA port. Three counters are available for each ports.

Received Cell Counters: 32-bit counters that counts the number of cells received on a particular UTOPIA port. Each port has a dedicated counter.

Cell loss counter: 16-bit counters that counts the number of cells lost in the UTOPIA module.

### 7.0 Programming the fast_clk PLL

The frequency received on mclk_src pin is used by the MT90503's PLL to generate a much higher frequency (fast_clk). It is then divided down to the output mem_clk frequency.

The $X, Y$ and $Z$ divider can be programmed to be any value as defined in Table 293 and Table 294 on page 201. The MT90503 can support mclk_src with a frequency ranging from 30 MHz to 80 MHz . Only frequencies between 50 MHz and 53.3 MHz are not supported by the PLL. The $X$ and Y divisor indicate what values can be programmed in the pll_conf registers 128h. Table 294, "Z Divisor Table," on page 201 indicates the range of output mem_clk that can be achieved. Note that the output mem_clk cannot be programmed to be above 80 MHz , or below 40 MHz .

The fast_clk PLL drives the output mem_clk pins. These pins provide both TTL and PECL interfaces for the output mem_clk. For both types, the output pins for the mem_clk is always driven. However, when the output pins are not being used, the register bits that control the toggling of these two pins should be disabled to reduce power consumption.

The user must configure the MT90503 to select the desired input mem_clk type, i.e., either PECL or TTL. The input mem_clk serves as the main clock (mclk) for the MT90503 and must be present for the MT90503 to function. It is absolutely necessary for the input mem_clk to be present and one of the inputs to be selected. The output mem_clk, however, are convenience for the user and do not have to be connected. These outputs eliminate the need for a second, high-speed oscillator to drive the input mem_clk.

The clock that is connected to the mem_clk inputs on the MT90503, whether it is the TTL or PECL, must be in phase with the clock connected to SSRAM used with the chip. The maximum skew allowed is $\pm 0.5 \mathrm{~ns}$.


Figure 50 - mem_clk Output and fast_clk Generation Circuits


Figure 51 - mem_clk Input and mclk Generation Circuit

| Div $\mathbf{X}$ | Div $\mathbf{Y}$ | mclk_src (MHz) | fast_clk (MHz) |
| :---: | :---: | :---: | :---: |
| - | - | 0 to 30 | - |
| 1 | 6 | 30 to 33.33 | 180 to 200 |
| 1 | 5 | 33.33 to 40 | 166.66 to 200 |
| 1 | 4 | 40 to 50 | 160 to 200 |
| - | - | 50 to 53.33 | - |
| 1 | 3 | 53.33 to 66.66 | 160 to 200 |
| 2 | 5 | 66.66 to 80 | 166.66 to 200 |
| 1 | 2 | 80 | 160 to 200 |

Table 293-Register 0128h Frequency Values

| Div Z | output mem_clk |
| :---: | :---: |
| 2 | 80 |
| 3 | 53.3 to 66.6 |
| 4 | 40 to 50 |

Table 294-Z Divisor Table

### 8.0 Electrical Specifications

### 8.1 DC Characteristics

## Absolute Maximum ratings

|  | Parameter | Symbol | Min | Max | Units |
| :---: | :--- | :---: | :---: | :---: | :---: |
| 1 | Supply Voltage - 3.3 Volt Rail | $\mathrm{V}_{\mathrm{DD}}$ | -0.3 | 3.9 | V |
| 2 | Voltage on 3.3V Input pins | $\mathrm{V}_{\mathrm{I}}$ | -1.0 | 3.6 | V |
| 3 | Continuous current at digital inputs | $\mathrm{I}_{\mathrm{I}}$ |  | 4.0 | mA |
| 4 | Continuous current at digital outputs | $\mathrm{I}_{\mathrm{O}}$ |  | 5.3 | mA |
| 5 | Storage Temperature | $\mathrm{T}_{\mathrm{s}}$ | -40.0 | +85.0 | ${ }^{\circ} \mathrm{C}$ |

* Exceeding these figures may cause permanent damage. Functional operation under these conditions is not guaranteed.

Voltage measurements are with respect to ground ( $V_{S S}$ ) unless otherwise stated. Long-term exposure to absolute maximum ratings may affect device reliability, and permanent damage may occur if operate exceeding the rating. The device should be operated under recommended operating condition.

## Recommended Operating Conditions

|  | Characteristics | Symbol | Min | Type | Max | Units | Test Conditions |
| :---: | :--- | :---: | :---: | :---: | :---: | :---: | :--- |
| 1 | Operating Temperature | $\mathrm{T}_{\mathrm{OP}}$ | -40.0 | 25.0 | +85.0 | ${ }^{\circ} \mathrm{C}$ | 2048 Channels <br> with heat sink <br> Note 1. |
| 2 | Operating Temperature | $\mathrm{T}_{\mathrm{OP}}$ | -40.0 | 25.0 | +70.0 | ${ }^{\circ} \mathrm{C}$ | 2048 Channels <br> with no heat sink |
| 3 | Operating Temperature | $\mathrm{T}_{\mathrm{OP}}$ | -40.0 | 25.0 | +85.0 | ${ }^{\circ} \mathrm{C}$ | 1024 Channels <br> with no heat sink |
| 4 | Supply Voltage, 3.3 Volt Rail | $\mathrm{V}_{\mathrm{DD}}$ | 3.0 | 3.3 | 3.6 | V |  |
| 5 | Input Voltage -3.3 V inputs | $\mathrm{V}_{\mathrm{I}}$ | $\mathrm{Vss}-0.5$ | 3.3 | $\mathrm{VDD}+0.3$ | V |  |

Note 1: Suitable heat sinks: Part Number 66435, Avvid Thermalloy and Part Number HS2141, Intricast, or other similar heat sinks. Typical figures are at $25^{\circ} \mathrm{C}$ and are for design aid only; not guaranteed and not subject to production testing.

Voltage measurements are with respect to ground $\left(\mathrm{V}_{\mathrm{SS}}\right)$ unless otherwise stated.

## DC Characteristics

|  | Characteristics | Symbol | Min | Typ | Max | Unit s | Test Conditions |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | Supply Current - 3.3 V supply | $\mathrm{I}_{\mathrm{DD}}$ |  | 720.0 |  | mA | 50.0 MHz, Nominal output loads, 1024 Channels, 16 streams active. |
| 2 | Supply Current - 3.3 V supply | IDDN |  | 970.0 |  | mA | 80.0 MHz, Nominal output loads, 2048 Channels |
| 3 | Device Power Dissipation | PDD1 |  | 2.38 |  | W | 50.0 MHz, Nominal output loads, 1024 Channels, 16 streams active. |
| 4 | Device Power Dissipation | $P_{\text {DD2 }}$ |  | 3.2 |  | W | 80.0 MHz, Nominal output loads, 2048 Channels |
| 5 | Input High Voltage | $\mathrm{V}_{\mathrm{IH}}$ | 2.0 |  | $\begin{gathered} \text { VDD+0 } \\ .3 \end{gathered}$ | V |  |
| 6 | Input Low Voltage | $\mathrm{V}_{\text {IL }}$ | Vss-0.5 |  | 0.8 | V |  |
| 7 | Switching Threshold | $V_{\text {TC }}$ |  | 1.4 | 2.0 | V |  |
| 8 | Schmitt Trigger Positive Threshold | $\mathrm{V}_{\text {t+ }}$ |  | 1.7 | 2.0 | V |  |
| 9 | Schmitt Trigger Neg. Threshold | $V_{t-}$ |  | 0.8 | 1.0 | V |  |
| 10 | Input Leakage Current | 1 | -10.0 |  | 10.0 | $\mu \mathrm{A}$ | $\mathrm{V}_{\text {IN }}=\mathrm{V}_{\mathrm{DDx}}$ or $\mathrm{V}_{\text {Ss }}$ |
| 11 | Input Pin Capacitance | C |  | 2.5 |  | pF |  |
| 12 | Output Pin Capacitance | $\mathrm{C}_{0}$ |  | 2.0 |  | pF |  |
| 13 | Output High Impedance Leakage | $\mathrm{I}_{\mathrm{OZ}}$ | -10.0 | +/-1.0 | 10.0 | $\mu \mathrm{A}$ | $\mathrm{V}_{\mathrm{O}}=\mathrm{V}_{\mathrm{SS}}$ or $\mathrm{V}_{\mathrm{DD}}$ |
| 14 | Output HIGH Voltage | $\mathrm{V}_{\mathrm{OH}}$ | 2.4 |  | VDD | V | $\mathrm{I}_{\mathrm{OH}}=$ rated current |
| 15 | Output LOW Voltage | $\mathrm{V}_{\text {OL }}$ |  | 0.2 | 0.4 | V | $\mathrm{I}_{\text {OL }}$ = rated current |
| 16 | 3.3V output HIGH current (4 mA buffer) | $\mathrm{I}_{\mathrm{OH}}$ |  |  | 4.0 | mA | $\mathrm{V}_{\mathrm{OH}}=2.4 \mathrm{~V}$ |
| 17 | 3.3 V output LOW current (4 mA buffer) | $\mathrm{IOL}^{\text {O }}$ |  |  | 4.0 | mA | $\mathrm{V}_{\mathrm{OL}}=0.4 \mathrm{~V}$ |
| 18 | 3.3 V output HIGH current (8 mA buffer) | $\mathrm{I}_{\mathrm{OH}}$ |  |  | 8.0 | mA | $\mathrm{V}_{\mathrm{OH}}=2.4 \mathrm{~V}$ |
| 19 | 3.3 V output LOW current (8 mA buffer) | $\mathrm{IOL}^{\text {L }}$ |  |  | 8.0 | mA | $\mathrm{V}_{\mathrm{OL}}=0.4 \mathrm{~V}$ |
| 20 | 3.3 V output HIGH current (12 mA buffer) | $\mathrm{I}_{\mathrm{OH}}$ |  |  | 12.0 | mA | $\mathrm{V}_{\mathrm{OH}}=2.4 \mathrm{~V}$ |
| 21 | 3.3 V output LOW current (12 mA buffer) | IOL |  |  | 12.0 | mA | $\mathrm{V}_{\mathrm{OL}}=0.4 \mathrm{~V}$ |
| 22 | Junction-to-Ambient Thermal Resistance | $\theta_{\mathrm{J}-\mathrm{A}}$ |  | 14.225 |  | ${ }^{\circ} \mathrm{C} / \mathrm{W}$ | 0 cfm air flow (natural convection airflow only) |

b. $\mathrm{T}_{\mathrm{OP}}=0^{\circ} \mathrm{C}$ to $70^{\circ} \mathrm{C} ; \mathrm{V}_{\mathrm{DD}}=3.3 \mathrm{~V} \pm 5 \%$

Voltage measurements are with respect to ground ( $\mathrm{V}_{\mathrm{SS}}$ ) unless otherwise stated.

### 8.1.1 Precautions During Power Sequencing

Latch-up is not a concern during power sequencing. The only requirement for sequencing 3.3 V and 5 V supplies during power up is that the MT90503 be either held in reset until the rails are stable or have its global_tri_state pin held low (tristate). However, to minimise over-voltage stress during system start-up, the 3.3 V supply applied to the MT 90503 should be brought to a level of at least $\mathrm{V}_{\mathrm{DD}}=3.0 \mathrm{~V}$ before a signal line is driven to a level greater than or equal to 3.3 V . This practice can be implemented either by ensuring that the 3.3 V power turns on simultaneously with or before the system 5 V supply turns on, or by ensuring that all 5 V signals are held to a logic LOW state during the time that $\mathrm{V}_{\mathrm{DD}}<3.0 \mathrm{~V}$. This condition is also met also if the MT90503 is held in reset until $\mathrm{V}_{\mathrm{DD}}$ reaches 3.0 V .

Regardless of the method chosen to limit over-voltage stress during power up, exposure must be limited to no more than +6.5 V input voltage $\left(\mathrm{V}_{\mathrm{IN}}\right)$. The global_tri_state pin of the MT90503 can be asserted low on powerup to prevent bus contention.

### 8.1.2 Precautions During Power Failure

Latch-up is not a concern in power failure mode. Although extended exposure of the MT90503 to 5 V signals during 3.3 V supply power failure is not recommended, there are no restrictions as long as $\mathrm{V}_{\mathrm{IN}}$ does not exceed the absolute maximum rating of 6.5 V . To minimise over-voltage stress during a 3.3 V power supply failure, the designer should either link the power supplies to prevent this condition or ensure that all 5 V signals connected to the MT90503 are held in a logic LOW state until the 5 V supply is deactivated.

### 8.1.3 Pull-ups

Pull-ups from the 5 V rail to 3.3 V (5 V tolerant) outputs of the MT90503 can cause reverse leakage currents into those 3.3 V outputs when they are active HIGH. (No significant reverse current is present during the high impedance state.) If the application can put the MT90503 in a state where MCLK is stopped, and a large number of 3.3 V output buffers are held in a static HIGH state, current can flow from the 5 V rail to the 3.3 V rail. If this MCLK-stopped state can not be avoided, the user should determine if the total MT90503 reverse current will have a negative impact on the system 3.3 V power supply. Alternatively, the global_tri_state pin of the MT90503 can be asserted low to put all outputs in the high impedance state.

### 8.2 H. 110 Diode Clamp Rail

As the MT90503 has a diode clamp to the 5 V rail, the diode clamp must be no more than 0.7 V below $\mathrm{V}_{\mathrm{DD}}$ when the pin is not tristated. This can be accomplished by asserting the global_tri_state pin low or by keeping the MT90503 in reset until all rails are stable.

### 8.3 AC Characteristics

All pins are tested with 50 pf worst case loading and 15 pf best case loading unless otherwise specified.

| Clock Name | Minimum <br> Frequency | Maximum <br> Frequency | Required For <br> Device Operation | Minimum <br> Duty Cycle | Maximum <br> Duty Cycle |
| :---: | :---: | :---: | :---: | :---: | :---: |
| mem_clk_i | 40 MHz | 80 MHz | Yes | $40 \%$ | $60 \%$ |
| mclk_src | 30 MHz | 80 MHz | Yes | $40 \%$ | $60 \%$ |
| rxa_clk | 1 MHz | 50 MHz | No | $40 \%$ | $60 \%$ |
| rxb_clk | 1 MHz | 50 MHz | No | $40 \%$ | $60 \%$ |

Table 295-Clock Networks

| Clock Name | Minimum <br> Frequency | Maximum <br> Frequency | Required For <br> Device Operation | Minimum <br> Duty Cycle | Maximum <br> Duty Cycle |
| :---: | :---: | :---: | :---: | :---: | :---: |
| rxc_clk | 1 MHz | 50 MHz | No | $40 \%$ | $60 \%$ |
| txa_clk | 1 MHz | 50 MHz | No | $40 \%$ | $60 \%$ |
| txb_clk | 1 MHz | 50 MHz | No | $40 \%$ | $60 \%$ |
| txc_clk | 1 MHz | 50 MHz | No | $40 \%$ | $60 \%$ |
| pll_clk | 8.192 MHz | 65.536 MHz | No | $40 \%$ | $60 \%$ |
| ct_c8_a/b | 8.192 MHz | 8.192 MHz | No | $40 \%$ | $60 \%$ |
| recov_a | 0 MHz | 80 MHz | No | $40 \%$ | $60 \%$ |
| recov_b | 0 MHz | 80 MHz | No | $40 \%$ | $60 \%$ |
| recov_c | 0 MHz | 80 MHz | No | $40 \%$ | $60 \%$ |
| recov_d | 0 MHz | 80 MHz | No | $40 \%$ | $60 \%$ |
| recov_e | 0 MHz | 80 MHz | No | $40 \%$ | $60 \%$ |
| recov_f | 0 MHz | 80 MHz | No | $40 \%$ | $60 \%$ |
| recov_g | 0 MHz | 80 MHz | No | $40 \%$ | $60 \%$ |
| recov_h | 0 MHz | 80 MHz |  |  |  |

Table 295 - Clock Networks (continued)

| Characteristic | Sym | Min | Typ | Max | Unit <br> $\mathbf{s}$ | Test Conditions |
| :--- | :---: | :---: | :---: | :---: | :---: | :--- |
| mem_clk_i Frequency | $\mathrm{t}_{\mathrm{MF}}$ | 40.0 | 80.0 | 80.0 | MHz | 30ppm clock recommended for <br> TDM PLLs |
| mem_clk_i Pulse Width <br> (HIGH / LOW) | $\mathrm{t}_{\mathrm{MH} / \mathrm{L}}$ | 5.0 | 6.25 | 7.5 | ns | For 80MHz operation. |

Table 296 - MCLK - Master Clock Input Parameters

### 9.0 Interface Timing

### 9.1 CPU Interface Timing



Note:
$\overline{\text { write_active }}=(\overline{\text { inmo_cs }}=0)$ AND $(\overline{\text { inmo_wr }}=0)$

Figure 52 - Non-multiplexed CPU Write Access - Intel Mode


Figure 53 - Non-multiplexed CPU Read Access - Intel Mode

| Symbol | Description Write Access | Min | Typ | Max | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: |
| t1 | write_access_active falling to inmo_d/inmo_a_das/ inmo a valid |  |  | 2*mclk_src - 4 | ns |
| t2 | inmo_rdy_ndtack rising to write_access_active rising | 0 |  |  | ns |
| t3 | inmo_rdy_ndtack rising to inmo_d/inmo_a_das/ inmo_a invalid | 0 |  |  | ns |
| t4 | write_access_active falling to inmo_rdy_ndtack falling | 0 |  | 12 | ns |
| t5 | Write Access Time |  |  | 740 | ns |
| t6 | write_access_active rising to inmo_rdy_ndtack tristate | 0 |  | 10 | ns |
| Symbol | Description Read Access | Min | Typ | Max | Unit |
| t1 | read_access_active falling to inmo_a_das/inmo_a valid |  |  | 2*mclk_src - 4 | ns |
| t2 | inmo_rdy_ndtack rising to read_access_active rising | 0 |  |  | ns |
| t3 | inmo_rdy_ndtack rising to inmo_a_das/inmo_a invalid | 0 |  |  | ns |
| t4 | read_access_active falling to inmo_rdy_ndtack falling | 0 |  | 12 | ns |
| t5 | Read Access Time |  |  | See Table 301 | ns |
| t6 | read_access_active rising to inmo_rdy_ndtack tristate | 0 |  | 10 | ns |
| t7 | inmo_d valid to inmo_rdy_ndtack rising | mclk_src - 4 |  |  | ns |
| t8 | read_access_active falling to inmo_d driving | $\begin{gathered} 3^{*} \mathrm{mclk} \text { _src } \\ -4 \end{gathered}$ |  |  | ns |

Table 297 - Non-multiplexed CPU Interface Intel Mode


Note:
$\overline{\text { write_active }}=(\overline{\overline{\text { inmo_cs }}=0)}$ AND $(\overline{\overline{\text { inmo_wr }}}=0)$

Figure 54 - Multiplexed CPU Write Access - Intel Mode


Figure 55 - Multiplexed CPU Read Access - Intel Mode


Figure 56 - Non-Multiplexed CPU Interface Write Access - Motorola Mode

| Symbol | Description Write Access | Min | Typ | Max | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: |
| t1 | write_access_active falling to inmo_ale falling |  |  | 2*mclk_src-4 | ns |
| t2 | write_access_active falling to inmo_d valid (writes) |  |  | 2*mclk_src-4 | ns |
| t3 | inmo_rdy_ndtack rising to write_access_active rising | 0 |  |  | ns |
| t4 | inmo_rdy_ndtack rising to inmo_ale rising \& inmo_d invalid | 0 |  |  | ns |
| t5 | Write Access Time |  |  | 740 | ns |
| t6 | write_access_active rising to inmo_rdy_ndtack tristate | 0 |  | 10 | ns |
| t7 | write_access_active falling to inmo_rdy_ndtack falling | 0 |  | 12 | ns |
| t8 | inmo_ale high pulse width | 5 |  |  | ns |
| t9 | inmo_d valid to inmo_ale falling | 5 |  |  | ns |
| t10 | inmo_ale falling to inmo_d invalid | 0 |  |  | ns |
| Symbol | Description Read Access | Min | Typ | Max | Unit |
| t1 | read_access_active falling to inmo_ale falling |  |  | 2*mclk_src - 4 | ns |
| t2 | inmo_rdy_ndtack rising to read_access_active rising | 0 |  |  | ns |
| t3 | inmo_rdy_ndtack rising to inmo_ale rising | 0 |  |  | ns |
| t4 | read_access_active falling to inmo_rdy_ndtack falling | 0 |  | 12 | ns |
| t5 | Read Access Time |  |  | See Table 301 | ns |
| t6 | read_access_active rising to inmo_rdy_ndtack tristate | 0 |  | 10 | ns |
| t7 | inmo_ale high pulse width | 5 |  |  | ns |
| t8 | read_access_active falling to inmo_d driving | 3*mclk_src - 4 |  |  | ns |
| t9 | inmo_d valid to inmo_rdy_ndtack falling | mclk_src - 4 |  |  | ns |
| t10 | inmo_d valid to inmo_ale falling | 5 |  |  | ns |
| t11 | inmo_ale falling to inmo_d invalid | 0 |  |  | ns |

Table 298 - Multiplexed CPU Interface Intel Mode


Figure 57 - Non-multiplexed CPU Interface Read Access - Motorola Mode

| Symbol | Description Write Access | Min | Typ | Max | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: |
| t1 | Address \& Data Setup -- $\overline{\text { inmo_cs }}$ and $\overline{\text { inmo_rd_ds }}$ asserted to inmo_a[14:0] and inmo_d[15:0] and inmo_a_das valid |  |  | 2*mclk_src-4 | ns |
| t3 | Address \& Data Hold -- inmo_rdy_ndtack low to inmo_a[14:0[ and inmo_d[15:0] and inmo_a_das invalid | 0 |  |  | ns |
| t4 | Inmo_rdy_ndtack high -- inmo_cs and inmo_rd_ds asserted to inmo_rdy_ndtack driving one | 0 |  | 12 | ns |
| t5 | Inmo_rdy_ndtack delay -- $\overline{\text { inmo_cs }}$ and inmo_rd_ds asserted to inmo_rdy_ndtack driving zero |  |  | 740 | ns |
| t6 | Inmo_rdy_ndtack hold -- $\overline{\text { inmo_cs }}$ and $\overline{\text { inmo_rd_ds }}$ de-asserted to inmo_rdy_ndtack drivingb one | 0 |  | 10 | ns |
| t7 | Inmo_rdy_ndtack high impedance -inmo_rdy_ndtack driving one to inmo_rdy_ndtack high impedance | 2 |  | 8 | ns |
| t11 | R/W setup -- R/W asserted to inmo_cs asserted | 0 |  |  | ns |
| t12 | R/W hold -- inmo_cs deasserted to R/W deasserted | 0 |  |  | ns |
| Note: $\mathrm{t} 1, \mathrm{t} 4$, and t 5 are dependent upon the last of $\overline{\mathrm{inmo}} \mathrm{cs}$ and $\overline{\mathrm{inmo} \text { rd_ds }}$ to be asserted. t 6 is dependent on the first of inmo_cs and inmo_rd_ds to be de-asserted. |  |  |  |  |  |
| Symbol | Description Read Access | Min | Typ | Max | Unit |
| t1 | Address Setup -- inmo_cs and inmo_rd_ds asserted to inmo_a[14:0[ and inmo_a_das valid |  |  | 2*mclk_src - 4 | ns |
| t3 | Address Hold -- inmo_rdy_ndtack low to inmo_a[14:0] and inmo_a_das invalid | 0 |  |  | ns |

Table 299 - Non-multiplexed CPU Interface Motorola Mode

| Symbol | Description Write Access | Min | Typ | Max | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: |
| t4 | Inmo_rdy_ndtack high -- inmo_cs and inmo_rd_ds asserted to inmo_rdy_ndtack driving one | 0 |  | 12 | ns |
| t5 | Inmo_rdy_ndtack delay -- $\overline{\mathrm{inmo}}$ _cs and $\overline{\mathrm{inmo} \text { _rd_ds }}$ asserted to inmo_rdy_ndtack asserted |  |  | See Table 301 | ns |
| t6 | Inmo_rdy_ndtack hold -- inmo_cs or inmo_rd_ds deasserted to inmo_rdy_ndtack driving one | 0 |  | 10 | ns |
| t7 | Inmo_rdy_ndtack high impedance -inmo_rdy_ndtack driving one to inmo_rdy_ndtack high-impedance | 2 |  | 8 | ns |
| t9 | Data to inmo_rdy_ndtack delay -- inmo_d[15:0] valid to inmo_rdy_ndtack asserted | mclk_src-4 |  |  | ns |
| t10 | Data output hold -- inmo_cs or inmo_rd_ds deasserted to inmo_d[15:0] invalid | 0 |  | 10 | ns |
| t11 | R/W setup -- R/W asserted to inmo_cs asserted | 0 |  |  | ns |
| t12 | R/W hold -- inmo_cs deasserted to R/W deasserted | 0 |  |  | ns |
| Note: t 1 , t 4 , and t 5 are dependent upon the last of $\overline{\mathrm{inmo}} \overline{\mathrm{cs}}$ and $\overline{\mathrm{inmo} \text { _rd_ds }}$ to be asserted. t , and t 10 are dependent on the first of inmo_cs and inmo_rd_ds to be-de-asserted. |  |  |  |  |  |

Table 299 - Non-multiplexed CPU Interface Motorola Mode (continued)


Notes:
$\overline{\text { write_active }}=(\overline{\text { inmo_cs }}=0)$ AND $(\overline{\text { inmo_ds }}=0)$ AND $($ inmo_r/ $\bar{w}=0)$

Figure 58 - Multiplexed CPU Interface Write Access - Motorola Mode


Figure 59 - Multiplexed CPU Interface Read Access - Motorola Mode

| Symbol | Description Write Access | Min | Typ | Max | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: |
| t1 | write_access_active falling to inmo_ale falling |  |  | 2*mclk_src - 4 | ns |
| t2 | write_access_active falling to inmo_d valid (writes) |  |  | 2*mclk_src - 4 | ns |
| t3 | inmo_rdy_ndtack falling to write_access_active rising | 0 |  |  | ns |
| t4 | inmo_rdy_ndtack falling to inmo_ale rising \& inmo_d invalid | 0 |  |  | ns |
| t5 | Write Access Time |  |  | 740 | ns |
| t6 | write_access_active rising to inmo_rdy_ndtack rising | 0 |  | 10 | ns |
| t7 | inmo_rdy_ndtack rising to inmo_rdy_ndtack tristate | 2 |  | 8 | ns |
| t8 | inmo_ale high pulse width | 5 |  |  | ns |
| t9 | write_access_active falling to inmo_rdy_ndtack driven high | 0 |  | 12 | ns |
| t10 | inmo_d valid to inmo_ale falling | 5 |  |  | ns |
| t11 | inmo_ale falling to inmo_d invalid | 0 |  |  | ns |
| Symbol | Description Read Access | Min | Typ | Max | Unit |
| t1 | read_access_active falling to inmo_ale falling |  |  | 2*mclk_src - 4 | ns |
| t2 | inmo_rdy_ndtack falling to read_access_active rising | 0 |  |  | ns |
| t3 | inmo_rdy_ndtack falling to inmo_ale rising | 0 |  |  | ns |
| t4 | read_access_active falling to inmo_rdy_ndtack driving high | 0 |  | 12 | ns |
| t5 | Read Access Time |  |  | See Table 301 | ns |
| t6 | read_access_active rising to inmo_rdy_ndtack rising | 0 |  | 10 | ns |
| t7 | inmo_rdy_ndtack rising to inmo_rdy_ndtack tristate | 2 |  | 8 | ns |
| t8 | read_access_active falling to inmo_d driving | 3*mclk_src - 4 |  |  | ns |
| t9 | inmo_d valid to inmo_rdy_ndtack falling | mclk_src-4 |  |  | ns |
| t10 | read_access_active rising to inmo_d tri-state | 0 |  | 10 | ns |
| t11 | inmo_ale high pulse width | 5 |  |  | ns |
| t12 | inmo_d valid to inmo_ale falling | 5 |  |  | ns |
| t13 | inmo_ale falling to inmo_d invalid | 0 |  |  | ns |

Table 300 - Multiplexed CPU Interface Motorola Mode

| Symbol | Description | Burst Length | Max | Unit |
| :---: | :---: | :---: | :---: | :---: |
| t5 | register and internal <br> memory access | 1 word | 740 | ns |
| t 5 | SSRAM | 1 word | 1.07 | $\mu \mathrm{~s}$ |
| t 5 | SSRAM | 8 words | 1.44 | $\mu \mathrm{~s}$ |
| t 5 | SSRAM | 128 words | 8.78 | $\mu \mathrm{~s}$ |

Table 301-t5 Read Access Time

### 9.2 UTOPIA Interface Timing



Figure 60 - UTOPIA Timing

|  | Characteristics | Symbol | Min | Max | Units |
| :---: | :--- | :---: | :---: | :---: | :---: |
| 1 | Input setup time | t 1 | 4 |  | ns |
| 2 | Input hold time | t 2 | 1 |  | ns |
| 3 | Clock to data valid | t 3 |  | 12 | ns |
| 4 | Clock to data change | t 4 | 2 |  | ns |
| 5 | Clock rising to signal driven | t 5 | 1 |  | ns |
| 6 | Clock rising to signal tri-state | t 6 | 1 | 20 | ns |

Table 302 - UTOPIA Bus Timing

### 9.3 External Memory Timing



Figure 61 - Flowthrough ZBT External Memory Timing - Write Access


Figure 62 - Flowthrough ZBT External Memory Timing - Read Access


Figure 63 - Flowthrough SSRAM External Memory Timing - Write Access


Figure 64 - Flowthrough SSRAM External Memory Timing - Read Access


Figure 65 - Late-write External Memory Timing - Write Access


Figure 66 - Late-write External Memory Timing - Read Access


Figure 67 - Pipelined ZBT External Memory Timing - Write Access


Figure 68 - Pipelined ZBT External Memory Timing - Read Access


Figure 69 - Pipelined External Memory Timing - Write Access


Figure 70 - Pipelined External Memory Timing - Read Access

|  | Characteristics | Symbol | Min | Typ | Max | Units | Test Conditions |
| :---: | :--- | :---: | :---: | :---: | :---: | :---: | :--- |
| 1 | Input setup time | t 1 | 2 |  |  | ns |  |
| 2 | Input hold time | t 2 | 0 |  |  | ns |  |
| 3 | Clock to data valid | t 3 |  |  | 8.30 | ns | Primetime tested (load = 50 pf) |
| 4 | Clock to data change | t 4 | 2 |  |  | ns | Primetime tested (load = 50 pf) |
| 3 s | Clock to data valid | t 3 |  |  | 7 | ns | Spice tested with 2 memory chips |
| 4 s | Clock to data change | t 4 | 2 |  |  | ns | Spice tested with 2 memory chips |
| 5 | Clock rising to signal driven | t 5 | 2 |  |  | ns |  |
| 6 | Clock rising to signal tri-state | t 6 |  |  | 10 | ns |  |

Table 303-Memory Interface Timing

## $9.4 \quad$ H.100/H. 110 Interface Timing



Figure 71-H. 100 Input, Output, and Frame Sampling
(100 Message Channel Clock

Figure 72-H. 100 Message Channel Clock, Transmission Delay, and Reception Delay


Figure 73-H. 100 Clock Skew (when chip is Master)

| Symbol | Description | Min | Typical | Max | Unit | Notes |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| t1 | ct_c8 rise to ct_d valid |  |  | 69 | ns |  |
| t2 | ct_c8 rise to ct_d invalid | 102 |  |  | ns |  |
| t3 | ct_d valid to ct_c8 fall | 3 |  |  | ns |  |
| t4 | ct_c8 fall to ct_d invalid | 1 |  |  | ns |  |
| t5 | ct_d valid to ct_c8 rise | 5 |  |  | ns |  |
| t6 | ct_c8 rise to ct_d invalid | 0 |  |  | ns |  |
| t10 | ct_c8 rise to ct_d tri-state |  |  | 122 | ns | 200 pf |
| t11 | ct_c8 rise to ct_d invalid | 102 |  |  | ns | 200 pf |
| t12 | ct_c8 rise to ct_d invalid | 2 |  |  | ns | 200 pf |
| t13 | ct_c8 rise to ct_d driven | 2 |  |  | ns | 200 pf |
| t14 | ct_c8 rise to ct_d valid |  |  | 22 | ns | 200 pf |
| t20 | ct_frame valid to ct_c8 rise | 5 |  |  | ns |  |
| t21 | ct_c8 rise to ct_frame invalid | 5 |  |  | ns |  |
| t30 | ct_c8 rise to mc_clock rise |  |  | 15 | ns |  |
| t31 | ct_c8 fall to mc_clock fall |  |  | 15 | ns |  |
| t40 | mc_tx fall to ct_mc low | 3 |  | 15 | ns | 200 pf |
| t41 | mc_tx rise to ct_mc tri-state | 3 |  | 15 | ns | 200 pf |
| t50 | ct_mc fall to mc_rx fall | 3 |  | 15 | ns |  |
| t60 | ct_c8_a, ct_c8_b, ct_frame_a, ct_frame_b, c2, c4, sclk, sclkx2, frcomp, c16+, c16- maximum skew when generated by the chip |  |  | 5 | ns | 200 pf |

Table 304-H.100/H. 110 Interface Timing

### 9.5 H.100/H.110 Clocking Signals

The MT90503's H.100/H. 110 interface generates all of the following signals with the specified timing.


Figure 74-H.100/H. 110 Clocking Signals


Figure 75 - TDM Bus Timing - Compatibility Clock Generation


Figure 76 - TDM Data Bus Timings

| Symbol | Parameter | Min | Typ | Max | Units |
| :---: | :--- | :---: | :---: | :---: | :---: |
|  | Clock edge rate (All Clocks) | 0.25 |  | 2 | V/ns |
| Tc8p | Clock CT_C8 Period | $122.066-\mathrm{F}$ |  | $122.074+\mathrm{F}$ | ns |
| Tc8h | Clock CT_C8 High Time | $49-\mathrm{F}$ |  | $73+\mathrm{F}$ | ns |
| Tc8l | Clock CT_C8 Low Time | $49-\mathrm{F}$ |  | $73+\mathrm{F}$ | ns |
| Tsamp | Data Sample Point |  | 90 |  | ns |
| Tdoz | Data Output to HiZ Time | -20 |  | 0 | ns |
| Tzdo | Data HiZ to Output Time | 0 |  | 22 | ns |
| Tdod | Data Output Delay Time | 0 |  | 22 | ns |
| Tdv | Data Valid Time | 0 |  | 69 | ns |
| Tdiv | Data Invalid Time | 102 |  | 112 | ns |
| Tfp | /CT_FRAME Width | 90 | 122 | 180 | ns |
| Tfs | ICT_FRAME Setup Time | 45 |  | 90 | ns |
| Tfh | ICT_FRAME Hold Time | 45 |  | 90 | ns |
| F | Phase Correction | 0 |  | 10 | ns |

Table 305-H.100/H. 110 Clocking Signals

### 10.0 Glossary of Terms

AALO: ATM Adaptation Layer $\mathbf{0}$. AALO is a straight packaging of 48 bytes of data within an ATM cell. AALO can be used to treat either data cells (managed by CPU) or CBR cells (managed by TX/RX SAR).

AAL1: ATM Adaptation Layer 1. AAL1 is used to transport constant bit rate (CBR) data on ATM. The main features of AAL1 are a cell sequence number that allows the detection of lost cells and a p-byte that allows reconverging of multi-channel VCs.

AAL5: ATM Adaptation Layer 5. The main feature of AAL5 is a 32-bit CRC at the end of the cell that allows the detection and correction of errors in the data payload. In this design, AAL5 cells are used uniquely to treat CBR information.

ATM: Asynchronous Transfer Mode. ATM is a networking standard based on 53-byte cells and is capable of carrying voice, data and video information simultaneously.

CAS: Channel Associated Signalling. Signalling bits used to indicate the state of the channel.
CBR: Constant Bit Rate. Cells in CBR format are sent out at a regular rate. CBR is applicable to voice channels.

CDV: Cell Delay Variation. When cells arrive on a UTOPIA port, they arrive with a certain delay with respect to when they were sent. CDV is a measure of how much that delay varies on a VC.

CLP: Cell Loss Priority. A 1-bit field in the ATM cell header that corresponds to the loss priority of a cell; cells with CLP = 1 can be discarded in a congestion situation.

CNT: Counter. Events in the MT90503 will cause the counter to increment.

CRC: Cyclic Redundancy Check. The CRC is a method of error detection and correction that is applied to a certain field of data. CRC is an efficient method of error detection because the odds of erroneously detecting a correct payload are low.

DS1: Digital-Signal Level 1. DS1 is an electrical interface for digital transmissions that contains 24 64-Kbps channels. The physical interface defined to carry DS1 channels is T1.

E1: E1 is the European equivalent of T1. They are similar with the main difference being E1 runs at 2.048 MHz instead of 1.544 Mbps , carrying 3064 kbps channels.

ESF: Extended Super-Frame. ESF is a T1 format that defines multiframes as consisting of 24 frames, each one of which contains 1 byte per channel.

FASTCAS: FASTCAS is not an acronym. It is capitalised in this document because it is a reserved word. FASTCAS means that multiframe integrity is not respected between the TDM and ATM buses. The TDM data is processed as soon as it is received, while CAS is sent when it is available.

FIFO: First In, First Out. A FIFO memory is one in which the first byte to have been written into the memory is the first one to be read from the read port.

GFC: Generic Flow Control. The GFC field is kept in the 4 highest bits of an ATM cell's header and is used for local functions (not carried end-to-end). The default value is "0000", meaning that GFC protocol is not enforced.

## GPI: General Purpose Input

## GPI/O: General Purpose Input or Output

H.100/H.110: A TDM bus standard developed by ECTF to provide backward compatibility to existing TDM buses with more bandwidth and potential for development.

HEC: Header Error Check. Using the fifth octet in the ATM cell header, ATM equipment may check for an error and correct the contents of the header. A CRC algorithm allows for single-error correction and multiple-error detection.

IE: Interrupt Enable. This is a register bit that enables a status event to generate an interrupt. This bit is always active-high.

## ISR: Interrupt Service Routine

JTAG: Joint Test Action Group.
LUT: Look-Up Table. In the UTOPIA module of the MT90503, the LUT is used to associate the data from received cells with the proper TDM output channels. The LUT is contained in the external memory.

MFS: Multi-Frame Support. The MT90503 is capable of supporting the multi-frame standards of E1 and T1.
MVIP: Multi-Vendor Integration Protocol. MVIP is a standard for transmitting data on a TDM bus.

NNI: Network-Node Interface. NNI ATM cells do not have a GFC nibble, instead having an extra nibble of VPI.
OAM: Operations Administration and Maintenance. MSB within the PTI field of the ATM cell header which indicates if the ATM cell carries management information such as fault indications.

OC-3: Optical Carrier level-3. A Sonet channel that carries a bandwidth of 155.55 Mbps .

OC-12: Optical Carrier level-12. A Sonet channel that carries a bandwidth of 622.08 Mbps.
PC: Process Control bit. This is a register bit type that is written to ' 1 ' to initiate a hardware process. When the process completes, the hardware clears the bit.

PCM: Pulse Code Modulation. PCM is the basic method of encoding an analog voice signal into digital form.

PHY: PHYsical layer. The bottom layer of the ATM Reference Model, it provides ATM cell transmission over the physical interfaces that interconnect the various ATM devices.

PLL: Phase Lock Loop. A phase lock loop is a component that generates an output clock by synchronising itself to an input clock. PLLs are often used to multiply the frequency of clocks.

PTI: Payload Type Identifier. The PTI field is a 3-bit header field that encodes various cell management information. Bit 2 (MSB) indicates OAM information or user data, bit 1 is Explicit Forward Congestion Control Indication (whether the cell may have been delayed by network congestion -- never examined by the MT90503) and bit 0 (LSB) indicates that a CBR-AAL5 cell is the final cell in a frame.

PUL: PULse bit. This is a register bit that is written to ' 1 ' to indicate an event to the hardware. This bit is always read at ' 0 '.

RAM: Random Access Memory. RAM is the main memory in the computer. It is called "random" because any random address can be accessed in an equal amount of time.

RO: Read Only. This serves to define registers that cannot be written to by the CPU.
ROL: Read Only Latch. This defines status bits. Status bits cannot be written to ' 1 ' by the CPU; however, once the status bit is set, the CPU can clear it by writing a ' 0 ' over it.

RW: Read Write. This type of register bit will be readable and writeable by the CPU.

SAR: Segmentation And Reassembly. Method of partitioning, at the source, frames into ATM cells and reassembling, at the destination, these cells back into information frames; lower sublayer of the AAL which inserts data from the information frames into cells and then adds the required header, trailer, and/or padding bytes to create 48 -byte payloads to be transmitted to the ATM layer.

## SCSA: Signal Computing System Architecture

SRTS: Synchronous Residual Time Stamp. SRTS is a clock recovery technique, which transmits timing information over the network to allow the source clock to be reconstructed at the other end. SRTS is sent in a 4 -bit value transmitted over 8 AAL1 cells.

T1: The North American telephone industry standard for interconnecting digital communication systems. A T1 channel consists of 24 voice channels, each at 64 kbps giving a total data rate of 1.544 Mbps .

TDM: Time Division Multiplexing. TDM buses carry voice data divided according to frames. In a single $125 \mu \mathrm{~s}$ frame, the TDM bus will have carried one byte from each channel it contains.

TSST: Time Slot/STream. A TSST represents a single voice channel on a TDM bus. For example, Time Slot 0, Stream 0 would be a TSST.

UNI: User-Network Interface. UNI ATM cells contain a GFC nibble in the header that is in place of four additional VPI bits in NNI ATM cells.

UTOPIA: Universal Test and Operations PHY Interface for ATM. A PHY-level interface to provide connectivity between ATM components.

VC: Virtual Circuit. VCs define a point-to-point connection between two nodes in a network. A single ATM cell carries data that belongs to a single VC.

VCC: Virtual Channel Connection.

VCI: Virtual Channel Identifier. This is the label given to an ATM VC to identify it and determine its destination. The VCI is a 16 -bit number that is included in the header of an ATM cell.

VPI: Virtual Path Identifier. A virtual path determines the way an ATM cell should be routed. The VPI is an 8-bit (in UNI) or 12-bit (in NNI) number that is included in the header of an ATM cell.

VTOA: Voice and Telephony Over ATM. A set of standards that enable voice and telephony services over ATM.

### 11.0 Physical Specification



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